We give a purely combinatorial proof of the Glaisher-Crofton identity which is derived from the analysis of discrete structures generated by the iterated action of the second derivative. The argument illustrates the utility of symbolic and generating function methodology of modern enumerative combinatorics. The paper is meant for nonspecialists as a gentle introduction to the field of graphical calculus and its applications in computational problems.

1. Introduction

Many computational problems involve action of complex expressions in derivatives on functions. A typical example is the exponential of a Hamiltonian acting on some initial condition which is a formal solution to the evolution equation. Applications of the latter range from classical heat and diffusion theory, financial mathematics, and economy to quantum field theory, hence practical interest in operational formulas enabling explicit evaluation of such expressions. Methods used to this effect usually involve operator and special function techniques, integral transforms, umbral calculus methods, etc. See comprehensive review of the subject [1, 2]. In this paper we develop another approach based on modern combinatorial methods of the analysis and enumeration of structures via generating functions [3–5].

The most known operational identities that involve the exponential of the first derivative are formulas for the shift and dilation operators

\[
\exp \left( \lambda \frac{d}{dx} \right) F(x) = F(x + \lambda)
\]

\[
\exp \left( \lambda x \frac{d}{dx} \right) F(x) = F(e^\lambda x)
\]

where \( F(x) \) is an arbitrary function. (Here, we leave subtle problems of convergence aside and consider \( F(x) \) as a formal power series in one variable \( x \).) They are a special case of the general closed-form operational expression

\[
\exp \left[ \lambda \left( q(x) \frac{d}{dx} + v(x) \right) \right] F(x) = g(\lambda, x) \cdot F(T(\lambda, x)),
\]

where functions \( T(\lambda, x) \) and \( g(\lambda, x) \) are specified by the following equations

\[
\frac{\partial T(\lambda, x)}{\partial \lambda} = q(T(\lambda, x)), \quad T(0, x) = x,
\]

\[
\frac{\partial g(\lambda, x)}{\partial \lambda} = v(T(\lambda, x)) \cdot g(\lambda, x), \quad g(0, x) = 1.
\]

See [1, 2, 6] for the proof based on operator techniques and [7, Sect. 6] for a recently developed combinatorial approach. Note that formulas (1) and (2) are valid for any function \( F(x) \). However, this is a very unique situation which holds only for the exponential of an expression linear in the first derivative. For the second derivative the closed-form formulas are not known and the best that can be done is evaluation on specific functions. There are only a few explicit
examples which include the formula for the exponential generating function of Hermite polynomials [8] and the Glaisher-Crofton identity [9–11]

\[ \exp \left( -\frac{1}{2} \frac{d^2}{dx^2} \right) \exp (xt) = \sum_{n=0}^\infty He_n(x) \frac{t^n}{n!} \]  

(5)

and the Glaisher-Crofton identity [9–11]

\[ \exp \left( \alpha \frac{d^2}{dx^2} \right) \exp (-x^2) = \frac{1}{\sqrt{1+4\alpha}} \exp \left( -\frac{x^2}{1+4\alpha} \right). \]  

(6)

Formulas of this type are usually derived using integral representations in the complex domain. (For example, to derive (6) one may use the integral representation \( \exp(-x^2) = \int_{-\infty}^{\infty} \exp(-x^2+2ix\xi)d\xi; \) see [10, 11].) However, in this paper we demonstrate that it is also possible to prove these identities on the basic algebraic level by analysing combinatorial structures generated by the iterated action of the second derivative. In this approach functions are treated as generating functions enumerating simple combinatorial objects (like sets of subsets, cycles, sequences, etc.) and consequently expressions in the derivatives transform objects in the initial class into richer structures whose generating functions can be quickly identified with the methods of symbolic combinatorics [3, 4]. We remark that methodology exposed in this note treads in the footsteps of combinatorial approach to algebraic identities developed by D. Foata in [12, 13].

Our goal in this paper is to develop and promote general combinatorial methodology for solving computational problems which, in many cases, provides better insight into algebraic and analytic manipulations. We illustrate this approach by explaining combinatorial meaning of the exponential of the second derivative and use this interpretation to derive equations (5) and (6).

2. Combinatorics of Derivatives

Action of derivatives on a function can be seen as a transformation of combinatorial structures. This viewpoint comes from interpreting the function as a generating function enumerating objects in some combinatorial class. In the following we formalise this intuition by describing the relevant constructions and develop a broader picture which includes higher derivatives and their exponentials. This framework will be illustrated by a simple combinatorial proof and interpretation of Taylor’s formula in Section 2.1 and equations (5) and (6) in Sections 2.2 and 3.

2.1. Generating Functions, First Derivative, and Taylor’s Formula. Let us consider a combinatorial class \( \mathcal{F} \) which is defined as a denumerable collection of objects built of atoms represented by \( X \) according to some well specified procedure. A typical combinatorial problem consists in enumeration of objects in \( \mathcal{F} \) according to the size \( |\cdot| : \mathcal{F} \rightarrow \mathbb{N} \) which is usually the number of atoms. In other words, one seeks the sequence \( f_n = \# \{ \phi \in \mathcal{F} : |\phi| = n \} \) which counts the number of objects comprised of exactly \( n \) atoms. This sequence can be encoded in a generating function

\[ F(x) = \sum_{\phi \in \mathcal{F}} x^{|\phi|} = \sum_{n=0}^\infty f_n x^n, \]  

(7)

which is a convenient tool for enumeration of complex structures via the so-called transfer rules. The latter translate combinatorial constructions into algebraic manipulations of the corresponding generating functions (see [3–5] for a comprehensive treatment of the subject and the appendix for a quick extract of a few transfer rules used in this paper). In the following we will briefly review combinatorics of the first derivative and recall a simple combinatorial interpretation of Taylor’s formula.

Here we will be concerned with the derivative operation \( yD_x \) acting on some well-defined class \( \mathcal{F} \) which consists in “selecting in all possible ways a single atom of type \( X \) in each element of \( \mathcal{F} \) and replacing it with an atom of a new type \( Y \).” In other words, one may think of the new class \( yD_x \mathcal{F} \) as formed of all structures taken from \( \mathcal{F} \) in which one of the atoms \( X \) gets “repainted” into a new colour \( Y \). Since each structure in \( \mathcal{F} \) built of \( n \) atoms of type \( X \) gives rise to \( n \) new ones with \( (n-1) \) atoms \( X \) and a single \( Y \), then the generating function enumerating objects in the new class \( yD_x \mathcal{F} \) is given by

\[ G(x, y) = \sum_{n, l \geq 0} g_{n,l} x^n y^l = \sum_{n \geq 1} \sum_{l=0}^{n} f_{n} x^{n-1} y, \]  

(8)

where \( g_{n,l} \) counts objects according to the number of \( X \)'s and \( Y \)'s, respectively. This is substantiated in the standard transfer rule:

\[ Selecting \ singleton:\]  

\[ \mathcal{G} = yD_x \mathcal{F} \Rightarrow \]  

\[ G(x, y) = y \frac{d}{dx} F(x). \]  

(9)

Now, let us consider the \( k \)-th derivative \( (1/k!)(yD_x)^k \) acting on \( \mathcal{F} \). Combinatorially it means “select in all possible ways an unordered collection of \( k \) atoms of type \( X \) and replace (repaint) them by atoms of type \( Y \).” Clearly, for each structure of size \( n \) in \( \mathcal{F} \) we have \( \binom{n}{k} \) possible choices, and hence the generating function of the new class \( (1/k!)(yD_x)^k \mathcal{F} \) evaluates to

\[ G(x, y) = \sum_{n, l \geq 0} g_{n,l} x^n y^l = \sum_{n \geq k} \sum_{l=0}^{n} \binom{n}{k} f_{n} x^{n-k} y^l. \]  

(10)

In consequence we get the following transfer rule:

\[ Selecting \ k-subset:\]  

\[ \mathcal{G} = \frac{1}{k!} (yD_x)^k \mathcal{F} \Rightarrow \]  

\[ G(x, y) = \frac{1}{k!} y^k \frac{d^k}{dx^k} F(x), \]  

(11)
which gives the combinatorial interpretation of the $k$-th derivative on the level of combinatorial structures.

This brings an interesting perspective on the derivative of a function which we will develop throughout the paper. Namely, one may think of a function $F(x)$ as a generating function of some combinatorial class $\mathcal{F}$. Then differentiation yields a new generating function which enumerates objects in the new class comprised of structures taken from $\mathcal{F}$ in which some of the atoms $\mathcal{X}$ were replaced with $\mathcal{Y}$'s (how many are replaced depends on the order of the derivative). Hence, the derivative of a function $F(x)$ can be understood as a well-defined combinatorial transformation of the associated combinatorial class $\mathcal{F}$ in a sense that on the level of generating functions it corresponds to simple differentiation (cf. (9) and (11)).

For illustration of this viewpoint let us recall the usual Taylor's formula

$$\sum_{k=0}^{\infty} \frac{y^k}{k!} f^{(k)}(x) = F(x + y). \quad (12)$$

Surprisingly, it admits a transparent combinatorial interpretation (see [3, Note III.31] or [7, Note 3]). To see this we observe that the left-hand side is the sum of derivatives $exp(yD_x)$ applied to some function $F(x)$ which can be considered as the generating function of some class of objects built of atoms $\mathcal{X}$. Then from our previous discussion the exponential $exp(yD_x)$ corresponds to

"selecting in all possible ways an arbitrary number of $\mathcal{X}$'s and replacing them by $\mathcal{Y}$'s."

(since the sum contains derivatives of arbitrary order we may choose subsets of arbitrary cardinality). On the other hand, this is the same as substituting each atom $\mathcal{X}$ either with atom $\mathcal{Y}$ (which makes no real effect) or with atom $\mathcal{Y}$ (which means the replacement). Hence we have the following combinatorial equivalence

$$\exp(yD_x)\mathcal{F} = \mathcal{F} \circ (\mathcal{X} + \mathcal{Y}), \quad (13)$$

which on the level of generating functions, by virtue of (11) and the transfer rule for substitutions (see the appendix, (A.4)), directly translates into (12). Hence from the combinatorial point of view Taylor's formula is a simple manifestation of the following transfer rule (cf. (13)):

Selecting arbitrary subset:

$$\mathcal{G} = \exp(yD_x)\mathcal{F} \implies \quad (14)$$

$$G(x, y) = F(x + y),$$

which applies to any combinatorial class $\mathcal{F}$ and its generating function $F(x)$. This is a typical example of combinatorial methodology which draws on the fact that in many cases the same combinatorial structure allows different specifications.

2.2. Second Derivative and Hermite Polynomials. Combinatorial interpretation of the second derivative can be developed along similar lines. From the above we know that the 2-nd derivative $(1/2)(yD_x)^2$ acting on $\mathcal{F}$ consists in

"selecting in all possible ways an unordered pair of $\mathcal{X}$ atoms and replacing the chosen $\mathcal{X}$'s by atoms of type $\mathcal{Y}$."

We will call such (unordered) pair a doubleton. Clearly, for an object composed of $n$ atoms this can be done in $n(n-1)/2$ ways, which agrees with the algebraic identity

$$(1/2)(yD_x)^2 x^n = (n(n-1)/2)x^{n-2}y^2.$$ 

More generally, by iterating $k$ times one picks out a sequence of doubletons in the original structure. Hence, we define the following construction:

Selecting $k$-subset of doubletons:

$$\mathcal{G} = \frac{1}{k!}\left(\frac{1}{2}(yD_x)^2\right)^k \mathcal{F} \implies \quad (15)$$

which consists in

"selecting in all possible ways a set of $k$ unordered pairs (doubletons) of $\mathcal{X}$ atoms and replacing them by $\mathcal{Y}$ atoms."

Note that, as in (11), we deem the order in the sequence irrelevant by introducing factor $1/k!$ in front of the iterated derivative (hence the “set” and not the “sequence” in the description). For a quick check of this specification we observe that the coefficient on the right-hand side of the identity

$$\frac{1}{k!}\left(\frac{1}{2}(yD_x)^2\right)^k x^n$$

$$= \frac{1}{k!} \frac{n(n-1)(n-2)(n-3)}{2} \cdots (n-2k+2)(n-2k+1) x^{n-2k} y^{2k} \quad (16)$$

coincides with the number of possible ways of choosing a set of $k$ unordered pairs from the set of $n$ objects, and hence by linearity we establish correctness of the description and transfer rule (15).

Now, we are in position to give interpretation of the exponential of second derivative:

Selecting arbitrary subset of doubletons:

$$\mathcal{G} = \exp\left(\frac{1}{2}(yD_x)^2\right)\mathcal{F} \equiv \sum_{k=0}^{\infty} \frac{1}{k!}\left(\frac{1}{2}(yD_x)^2\right)^k \mathcal{F}, \quad (17)$$

whose combinatorial meaning comes down to

"selecting in all possible ways an arbitrary subset of (unordered) pairs of $\mathcal{X}$ atoms and replacing each chosen $\mathcal{X}$ by atom of type $\mathcal{Y}$."

This is the sum of the derivative operations of the type (15) which on the level of generating functions is the sum of the derivatives. Unfortunately it does not come close to any
of the derivative operator terms of enumeration of structures. Our aim is to understand these formulas in (or weight) which, if needed, can be specified to the required structures proceeds also with respect to this additional label allows getting rid of negative or noninteger factors entering trick to introduce additional labels (or weights) which often only (cf. coefficients in (18)). (It is a typical combinatorial interpretation exp \( \exp(\frac{1}{2}(yD_x)^2) \)) of the exponential of the second derivative on the gaussian. is the Glaisher-Crofton formula (6) which evaluates action of the exponential and differentiating the combinatorial model of Hermite polynomials (cf. [4, 12, 13]). For the purpose at hand we will leave it is the exponential generating function of Hermite polynomials (cf. (5)). For the purpose at hand we will leave case of the action on monomials and provide a link with will illustrate our combinatorial methodology on a simpler before we proceed to this result, discussed in Section 3, we will illustrate our combinatorial methodology on a simpler case of the action on monomials and provide a link with combinatorial model of Hermite polynomials (cf. [4, 12, 13]).

Let us start with the explicit expression which is obtained from expanding the exponential and differentiating the monomial, i.e.,

\[
\exp\left(\frac{1}{2}(yD_x)^2\right) x^n \equiv \sum_{k=0}^{\lfloor n/2 \rfloor} \frac{n!}{2^k k! (n-2k)!} x^{n-2k} y^{2k}.
\]

(18)

For \( y = i \) it specialises to the Hermite polynomial \( H_n(x) \). More generally, we may also write

\[
\exp\left(\frac{1}{2}(yD_x)^2\right) \exp(xt) = \exp\left(\frac{1}{2}(yt)^2\right) \cdot \exp(xt)
\]

(19)

which stems from the fact that \( \exp(xt) \) is an eigenvector of the derivative operator \( D_x \) to eigenvalue \( t \). Again, for \( y = i \) it is the exponential generating function of Hermite polynomials (cf. (5)). For the purpose at hand we will leave variable \( y \) unspecified so as to deal with positive integers only (cf. coefficients in (18)). (It is a typical combinatorial trick to introduce additional labels (or weights) which often allows getting rid of negative or noninteger factors entering multiplicatively in the expressions. Then enumeration of structures proceeds also with respect to this additional label (or weight) which, if needed, can be specified to the required value at the end.) Our aim is to understand these formulas in terms of enumeration of structures.

In order to use combinatorial description of (17) we interpret \( \exp(xt) \) as the exponential generating function of the labelled class of sets \( \text{Set}(X\mathcal{F}) \) (see [3, Sect. II] for a precise definition and discussion of labelled classes and their relation with exponential generating functions). It is comprised of sets whose atoms \( \mathcal{F} \) carry integer labels 1, 2, 3, . . . , and additionally to each \( \mathcal{F} \) an atom (or weight) of type \( X \) is attached; see Figure 1 on the left. We have the following translation rule (see (A.5)):

\[
\mathcal{F} = \text{Set}(X\mathcal{F}) \implies F(x,t) = \sum_{n \geq 0} x^n t^n \frac{n!}{n!} = \exp(xt).
\]

(20)

(Clearly, for given \( n \) one can built one such set and its weight is \( x^n \).) Now, following the combinatorial description (17) action of \( \exp((1/2)(yD_x)^2) \) on an individual set in \( \text{Set}(X\mathcal{F}) \) consists in selecting in all possible ways a subset of doubletons. This amounts to splitting of each original set into products of two subsets: one comprising singletons and the other doubletons. Additionally, these subsets differ in that the atoms in the singletons (untouched by the derivatives) carry the weight \( X \), while each atom forming the doubleton (arising from nontrivial action of the second derivative) carries the label \( Y \). Yet another way of seeing the resulting class of objects is to understand them simply as a set of singletons \( X\mathcal{F} \) and doubletons \((1/2)(Y\mathcal{F})^2 \). See Figure 1 for illustration. Formally, one writes the following sequence of combinatorial equivalences:

\[
\exp\left(\frac{1}{2}(yD_x)^2\right) \text{Set}(X\mathcal{F}) = \text{Set}(X\mathcal{F}) \times \text{Set}\left(\frac{1}{2}(Y\mathcal{F})^2\right) = \text{Set}\left(X\mathcal{F} + \frac{1}{2}(Y\mathcal{F})^2\right),
\]

(21)

which on the level of generating functions readily transform (cf. the appendix) into a sequence of algebraic equalities providing a combinatorial proof of (19).

As a consequence of this discussion we get a simple combinatorial model of Hermite polynomials. Namely, coefficients of \( He_n(x) = \sum_{k=0}^n h_{n,k} x^k \) count the number of possible ways of

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**Figure 1:** On the left, there is an instance of a structure in class \( \mathcal{F} = \text{Set}(X\mathcal{F}) \), i.e., a set of labelled atoms \( \mathcal{F} \) which are additionally weighted by an atom of type \( X \). In the middle, example of a structure arising from the action of \( \exp((1/2)(yD_x)^2) \) on \( \mathcal{F} \) which is obtained by selecting a set of doubletons of \( X \) atoms and replacing them by atoms of type \( Y \). Selected pairs are depicted by wavy lines. On the right, decomposition of such a structure into a product of two sets comprising, respectively, singletons \( X\mathcal{F} \) (untouched by the derivatives) and doubletons \((1/2)(Y\mathcal{F})^2 \) (selected by the derivatives).
“selecting a subset of \(k\) doubletons (each weighted by \(-1\)) out of a set of \(n\) distinguishable objects.”

It is exactly the coefficient of \(x^{n-2k}\) in (18) for \(y = i\). Another way to see it directly from our combinatorial description of the exponential of the second derivative is to interpret \(x^n\) as the generating function of a single \(n\) element set; i.e., we have \(\mathcal{T} = \text{Set}_n(\mathcal{X}) \implies F(x) = x^n\). Then combinatorial model in terms of the choice of doubletons is a simple consequence of the specification of (17). We note that this model is a rephrasing of the interpretation of Hermite polynomials in terms of weighted involutions \([4, \text{Sect. 2.3.}]\). Moreover, it can be straightforwardly extended to provide a combinatorial interpretation of a larger class of multivariate Hermite-Kampé de Féret polynomials \([1, 2]\).

3. Proof of the Glaisher-Crofton Identity

Here we prove the identity (6) by a purely combinatorial argument by analysing structures generated by the second derivative discussed above. We will proceed in a step by step manner explaining the details of combinatorial constructions and structures that appear along the way. Although most of them are standard in combinatorial community we take a rather explicit and methodological route that may be of help for an unaccustomed reader. In the following we adopt the standard notation from the book [3] (see also the appendix).

Our goal is to calculate the explicit form of the expression (cf. the left hand side of (6) with \(\alpha = y^2/2\) and \(t = -1\))

\[
\exp \left( \frac{1}{2} (yD_x)^2 \right) \exp (x^2t) = \sum_{k=0}^{\infty} \frac{1}{k!} (yD_x)^{2k} \cdot \frac{(x^2t)^n}{n!} = R(x, y, t).
\]

Following our combinatorial strategy we will treat \(\exp(x^2t)\) as an exponential generating function enumerating some combinatorial objects. Let us define them as a labelled class of sets \(\text{Set}(\mathcal{X}^2 \mathcal{T})\). This class is comprised of sets of labelled atoms \(\mathcal{T}\) (i.e., each atom carries the integer label) which are weighted with two atoms of type \(\mathcal{X}\). We will depict them as sets of doubletons of unlabelled atoms \(\mathcal{X}\) such that each doubleton carries the labelled marker \(\mathcal{T}\). See Figure 2 on the left for illustration. Clearly, we have the following transfer rule (cf. (A.5)):

\[
\mathcal{C} = \text{Set}(\mathcal{X}^2 \mathcal{T}) \implies C(x,t) = \sum_{n \geq 0} x^{2n} t^n/n! = \exp (x^2t).
\]

We know from the discussion of (17) that on the combinatorial level exponential \(\exp((1/2)(yD_x)^2)\) consists in selecting in all possible ways unordered pairs of \(\mathcal{X}'s\) (not necessarily attached to the same \(\mathcal{T}\)) and replacing each \(\mathcal{X}\) in the chosen pairs by \(\mathcal{Y}\). Figure 2 in the middle illustrates a generic structure arising in this procedure. If we denote the resulting class of structures by \(\mathcal{R}\), then we may write

\[
\mathcal{R} = \exp \left( \frac{1}{2} (yD_x)^2 \right) \mathcal{C} \implies \\
\exp \left( \frac{1}{2} (yD_x)^2 \right) \exp (x^2t) = R(x, y, t),
\]

where \(R(x, y, t) = \sum_{k,l \geq 0} R_{k,l} x^k y^l (t^n/n!)\) is the exponential generating function enumerating structures in \(\mathcal{R}\). Note that this gives a precise combinatorial meaning to (22).

Since we have reduced our goal to finding the exponential generating function \(R(x, y, z)\) one needs to come up with a systematic specification of structures in \(\mathcal{R}\). For this purpose
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Figure 3: A generic closed chain in class $\mathcal{A}$ embedded in the plane can be seen as a cycle built out of two kinds of blocks, each of type $\mathcal{Y}^2\mathcal{T}$, arising from different choices of $\mathcal{X}$'s attached to the same $\mathcal{T}$ by the derivative on either the left or right (hence possible crossings of the wavy lines).

Figure 4: If embedded in the line, each open chain in $\mathcal{B}$ forms a sequence whose inner part (in the grey box) is built out of blocks of type $\mathcal{Y}\mathcal{T}\mathcal{Y}$ with two blocks of type $\mathcal{X}\mathcal{Y}\mathcal{T}$ attached at the ends. Here again, each block can have two configurations due to different choices made by the derivatives (possible crossings of the wavy lines).

Let us observe that doubletons, initially detached from one another in $\mathcal{C}$, are now tied together to form either open or closed chains. Then, we may group them together splitting each structure in $\mathcal{R}$ into a product of two sets: one containing only closed and the second open chains. This entails the following combinatorial decomposition and its translation to the exponential generating functions (cf. (A.3) and (A.5)):

$$\mathcal{R} = \text{Set}(\mathcal{A}) \times \text{Set}(\mathcal{B}) \implies R(x, y, t) = \exp(A(x, y, t)) \cdot \exp(B(x, y, t)),$$

where $\mathcal{A}$ is the class of open chains and $\mathcal{B}$ is the class of closed chains; see Figure 2 on the right. Hence the problem comes down to finding both exponential generating functions $A(x, y, t)$ and $B(x, y, t)$ enumerating, respectively, objects of types $\mathcal{A}$ and $\mathcal{B}$.

Let us start with the class of closed chains $\mathcal{A}$. Its elements embedded in the plane can be seen as cycles whose building blocks have a finer structure of type $\mathcal{Y}^2\mathcal{T}$ which occurs in two possible configurations arising from two possible choices of $\mathcal{X}$ by the derivative in the same doubleton; see Figure 3 for pictorial explanation. Therefore the whole class is specified as follows:

$$\mathcal{A} = \text{Cyc}(2\mathcal{Y}^2\mathcal{T}),$$

and by the standard transfer rule for labelled classes, (A.6), we get

$$A(x, y, t) = \log \frac{1}{1 - 2y^2t}.$$ (27)

The class of open chains $\mathcal{B}$ can be described in a similar manner. First we observe that each such chain can be embedded in the line in two possible ways. Then it forms a sequence which can be decomposed into the inner part which is a sequence of blocks of type $\mathcal{Y}\mathcal{T}\mathcal{Y}$ with two additional blocks of types $\mathcal{X}\mathcal{T}\mathcal{X}$ and $\mathcal{Y}\mathcal{T}\mathcal{X}$ attached at the ends (left and right, respectively). Here as well, each block occurs in two possible configurations arising from two possible choices of $\mathcal{X}$ by the derivative in the same doubleton. See Figure 4 for illustration. This gives the following combinatorial specification:

$$\mathcal{B} = \mathcal{X}\mathcal{T}\mathcal{X} + \frac{1}{2} (2\mathcal{X}\mathcal{Y}\mathcal{T}) \times \text{Seq}(2\mathcal{Y}\mathcal{T}\mathcal{Y}) \times (2\mathcal{Y}\mathcal{T}\mathcal{X}).$$ (28)

Note that the coefficient $1/2$ stems from the double counting due to embedding in a line and the term $\mathcal{X}\mathcal{T}\mathcal{X}$ makes up for a single structure left out by the above description. Having specified $\mathcal{B}$ we obtain exponential generating function by means of the standard transfer rules, (A.7), which yield
\[ B(x, y, t) = x^2 t + \frac{1}{2} 2xyt \cdot \frac{1}{1 - 2y^2t} \cdot 2xyt \]
\[ = \frac{x^2 t}{1 - 2y^2t} \] (29)

Now, by substituting (27) and (29) to (25), we get
\[ R(x, y, t) = \frac{1}{1 - 2y^2t} \exp \left( \frac{x^2 t}{1 - 2y^2t} \right). \] (30)

This completes the proof of the Glaisher-Crofton identity, (6), which is readily obtained from (22). In conclusion, let us remark that we benefit from the proof by a deeper combinatorial insight into the nature of both factors on the r.h.s. of (6). The latter can be interpreted as exponential generating functions enumerating sets of closed and open chains, respectively, formed by derivatives acting on the gaussian.

4. Discussion and Outlook

Many computational problems require keeping track and skilful rearrangement of terms involved in algebraic expressions. It often comes down to the analysis of their structural properties and counting terms grouped with respect to some relevant characteristics. This is a natural domain of application for modern combinatorics which has developed a large array of tools for systematic treatment of such problems. In this paper, we have considered a few examples where it can be effectively used for evaluation of the action of the exponential in the derivatives on a function. Fundamental in this approach is treatment of the function as a generating function enumerating some simple combinatorial objects. This shift in perspective allows interpreting the derivatives (and their exponentials) as combinatorial constructors which produce a new class of objects which often can be enumerated with combinatorial flair. We have illustrated this approach by showing simple combinatorial proofs and interpretation of Taylor’s formula, connecting exponential in second derivative with a model of Hermite polynomials, and deriving the Glaisher-Crofton identity. It is worth emphasising that the proofs are purely combinatorial and thus do not require any arguments involving integral representations or analyticity. We note that our exposition builds on the seminal papers of D. Foata [12–14] who gave combinatorial insight into the related Mehler formula. In this paper, the focus is shifted towards graphical calculus based on symbolic methods in combinatorial insight into the nature of both factors on the r.h.s. of (6). The latter can be interpreted as enumerating combinatorial structures built of atoms of several kinds.

Finally, let us remark that combinatorial approach to derivatives also provides an interesting insights into operator identities. One example is a systematic treatment of the normal ordering problem [7, 20, 21]. Clearly, majority of operator identities admit combinatorial interpretation as they typically arise from algebraic manipulation of discrete structures [4, 22, 23]. As such it opens the whole field of application for combinatorial approach. In this paper we have illustrated this point only on a few examples which can be seen as instances of a broad class of operator identities amenable to combinatorial methodology (cf. Sack identity, Baker-Campbell-Hausdorff formula, Rodrigues-type formulas, Crofton identities, etc. [1, 2]).

Appendix

A. Combinatorial Constructions

Our primary reference for combinatorial analysis is the standard book Analytic Combinatorics by Ph. Flajolet and R. Sedgewick [3]. Here, we briefly recall basic terminology and a few standard translation rules for labelled constructions used in this paper.

Suppose we are given a combinatorial class \( C \) which consists of a denumerable collection of objects built of the labelled atoms \( \mathcal{F} \) (see [3, Ch. II] for precise definition of the labelled class). Usually size of an object is the number of atoms it is built of, and a typical problem is to count the number of structures of a given size. In other words, one seeks the sequence \( c_n = \# C_n \), where \( C_n = \{ \phi \in C : |\phi| = n \} \) which is conveniently encoded in the exponential generating function (e.g.f.)

\[ C(t) = \sum_{n \geq 0} c_n \frac{t^n}{n!} = \sum_{\phi \in \mathcal{F}} t^{|\phi|}. \] (A.1)

Let us remark that the reason for the use of exponential generating functions, rather than ordinary generating functions (o.g.f.), is simplicity of transfer rules in the domain of labelled classes. (Ordinary generating functions are typically used for enumeration of unlabelled structures, cf. [3, Ch. I].)

The point of combinatorial analysis of structures is construction of complex classes from simpler ones. The initial building blocks include the atomic class \( \mathcal{F} \), which comprises a single element of size 1 and has e.g.f. \( C(t) = t \), and...
the neutral class $\mathcal{C}$, which consists of a single element of size 0 and has e.g.f. $C(t) = 1$. Then complex structures are built by well defined set of theoretical constructions which provide a precise specification of the class. Remarkably, these constructions can be translated into algebraic equations for the corresponding generating functions which solve the enumeration problem. Below, we give a short list of such constructions and translation rules that we exploit in this paper.

The most basic one is the disjoint union, henceforth denoted by “+”, which corresponds to

$$C = \mathcal{A} + \mathcal{B} \implies C(t) = A(t) + B(t).$$

(A.2)

Another one is the labelled product, denoted by “×”, which forms a cartesian product of objects and relabels the atoms in order-consistent manner. We have the following translation rule:

$$C = \mathcal{A} \times \mathcal{B} \implies C(t) = A(t) \cdot B(t).$$

(A.3)

If objects of one structure $\mathcal{B}$ are substituted into atoms of another structure $\mathcal{A}$ and relabelled in the order-consistent way, then the e.g.f. of such constructed class is given by (assuming $\mathcal{B}_0 = \emptyset$, i.e., $B(0) = 0$)

$$C = \mathcal{A} \circ \mathcal{B} \implies C(t) = A(B(t)).$$

(A.4)

It is then possible to form the class of all (labelled) sequences, sets, and cycles (respectively denoted by Set, Seq, and Cyc) built from objects in $\mathcal{A}$. The corresponding generating functions are given by the following dictionary (assuming $\mathcal{A}_0 = \emptyset$, i.e., $A(0) = 0$):

$$C = \text{Set}(\mathcal{A}) \implies C(t) = \exp(A(t)),$$

(A.5)

$$C = \text{Cyc}(\mathcal{A}) \implies C(t) = \log\frac{1}{1-A(t)},$$

(A.6)

$$C = \text{Seq}(\mathcal{A}) \implies C(t) = \frac{1}{1-A(t)}.$$  

(A.7)

This is a nonexhaustive selection of possible constructions which is used in the present paper. For a comprehensive survey of the methods of combinatorial enumeration via generating functions, we refer to the classic books on this subject [3–5].

**Conflicts of Interest**

The authors declare that there are no conflicts of interest regarding the publication of this paper.

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**References**


