

Research Article

Cooperative Networking towards Maritime Cyber Physical Systems

Tingting Yang,¹ Hailong Feng,¹ Chengming Yang,² Zhonghua Sun,¹ Jiadong Yang,¹ Fan Sun,¹ Ruilong Deng,³ and Zhou Su⁴

¹Navigation College, Dalian Maritime University, 1 Linghai Road, Dalian, Liaoning 116026, China

²School of Naval Architecture, Ocean & Civil Engineering, Shanghai Jiao Tong University, Shanghai 200240, China

³Department of Electrical and Computer Engineering, University of Alberta, Edmonton, AB, Canada T6G 1H9

⁴School of Mechatronic Engineering and Automation, Shanghai University, Shanghai 200444, China

Correspondence should be addressed to Tingting Yang; yangtingting820523@163.com and Zhou Su; zhousu@ieee.org

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An innovative paradigm named Cooperative Cognitive Maritime Cyber Physical Systems (CCMCPs) is developed to achieve high-speed and low-cost communication services. The analysis of the available white space at sea, as well as the framework, is presented. Specifically, a bilevel game with two stages of PUs-to-SUs (primary users to secondary users) and SUs-to-SUs is proposed, to address the resource allocation issue of Decode-and-Forward (DF) relay mode with maximal-ratio combining (MRC) receiving mode in destination. Stackelberg game with priority is employed between PU and SUs, while a symmetrical system model is considered among SUs-to-SUs. The game theoretic procedure that converges to Nash equilibrium based on the utility and payoff function is illustrated. Simulation results demonstrate that our proposed strategy could effectively increase the throughput as well as the payoffs of the system.

1. Introduction

Maritime Cyber Physical Systems (CPSs) target the tight incorporation of wireless communications, control, and computing technologies into the navigation transportation system, which possesses the typical characteristics of CPSs revolutionizing the navigation pattern to be safer and more efficient through real-time embedded systems for distributed sensing, computation, and control between cyber and physical systems [1]. As we can see, the intelligent transportation systems in maritime field have proliferative interest in vis-a-vis maritime CPSs to achieve a low-cost alternative for current maritime satellite system for data transmission. Some advanced wireless technologies have been previously adopted to build maritime wideband networks. For example, the Worldwide Interoperability for Microwave Access (WiMAX) technology is exploited for wireless-broadband-access for Seaport (WISEPORT) project in Singapore [2]. The Fourth Generation Long-Term Evolution (4G LTE) broadband service has been developed in US Navy ships since 2011

[3]. In our previous work [4, 5], a novel complementary scheme, for example, delay tolerant networks (DTNs) which exploits store-carry-and-forward routing scheme [6], is utilized in maritime scenario. However, the above scheme should apply the limited authorized spectrum resource for maritime communication, which is an important restriction for maritime communication. The cognitive technology has been already utilized to construct maritime mesh/ad hoc networks [7], to alleviate the insufficient spectrum provision for maritime purpose indeed. At the same time, a maritime cognitive mesh/ad hoc network could achieve the connection among neighboring vessels, sea farm, oil/gas platform, and marine beacons and buoys to sense the maritime information, exchange it, and send it to the land-based network through wireless communication between vessel-to-vessel and vessel to-infrastructure. In the maritime CPSs, all types of physical systems involving intelligent vessels should be equipped with seamlessly integrated embedded computing systems and in-vessel networking systems, which have the convergent functions of computation,

communication, and control. The maritime CPSs involve interactions between vessel controllers, communication networks, and physical world. The behaviors of the physical world such as the conditions of cargo and the vessel are dynamic and continuous changing with time while the process of communication and calculation in maritime CPSs is discrete. Endowed with control, monitoring, and data gathering functions, the maritime CPSs dramatically enhance the controllability, adaptability, efficiency, functionality, reliability, safety, and usability of maritime service. And thereby they promote a great deal of magnetic utilizations in terms of monitoring (e.g., transferring supervision videos gathering internal or external of vessels to shore-based authority being uploaded after collecting from bridge, engine room, or other critical places), safety (e.g., maritime safety information dissemination), infotainment (e.g., mobile office and multimedia data download and upload, especially for cruise industry), and cargo online management (e.g., real-time cargo status notification and handling management).

However, the remarkable fundamental characteristics of maritime environment are long transferring distance and irregularly worsening channel due to the obstruction, that is, sea clutters. Therefore, the communication distance range is still unsatisfactory. Fortunately, the cooperative technology fits the maritime communication environment very much, which could achieve diversity gain through collaboration between entities depending on relay transmission principle [8]. It is employed to overcome the disadvantage of channel fading, strengthen the reliability of delivery, augment the coverage, and utilize the spectrums definitely. In this paper, a state-of-the-art prototype, that is, Cooperative Cognitive Maritime CPSs (CCMCPSs), is developed on the basis of wireless mesh/ad hoc networks to provide high-speed and low-cost communications for maritime paradigm. It is envisioned that the CPSs which incorporate information communications technology and sensing-enabled vessels could impose new opportunities, application, and agendas as well as challenges for maritime community. However, the resource allocation and scheduling issue in Cooperative Cognitive Maritime CPSs are challenging, which is still an open issue.

In this paper, we will describe this new system in detail, including the framework, available white space on the sea, regulation requirement, and standards in maritime society. Heterogeneous network framework of CPSs is proposed which refers to the resource allocation issues between PUs-to-SUs (primary users to secondary users) and SUs-to-SUs, respectively. Specifically, a bilevel game with two stages of PUs-to-SUs and SUs-to-SUs is proposed to address the resource allocation issue of Decode-and-Forward (DF) relay mode with maximal-ratio combining (MRC) receiving mode in destination. Stackelberg game with priority is employed between PU and SUs, while a symmetrical system model is considered among SUs-to-SUs. Specifically, the contribution of this paper is threefold.

- (1) An innovative paradigm named Cooperative Cognitive Maritime Cyber Physical Systems is developed to achieve high-speed and low-cost communication.

services, based on the analysis of available white space at sea.

- (2) A bilevel game with two stages of PUs-to-SUs and SUs-to-SUs is proposed to address the resource allocation issue of DF relay mode with MRC receiving mode in destination.
- (3) Stackelberg game with priority is employed between PU and SUs, while a symmetrical system model is considered among SUs-to-SUs. And the game theoretic procedure that converges to Nash equilibrium based on the utility and payoff function is illustrated.

The remainder of this paper is organized as follows. Section 2 reviews the related work. System model and white space analysis are illustrated in Section 3. Section 4 describes the Stackelberg game based resource allocation between PU and SUs. Section 5 demonstrated the game theory of symmetrical cooperative three-node model among SUs. In Section 6, simulation results validate the performance of our scheme. Finally, conclusions and references end this paper. Since many symbols are used in this paper, some important notation definitions are tabulated in Notations.

2. Related Work

The research related to CPSs has drawn much attention recent years. Kim and Kumar gave a nearly exhaustive survey that covers every aspect of CPS research [1]. The achievements in many research disciplines have been reviewed, including the history of CPSs, some advanced theoretical foundations and technologies of networked control systems, hybrid systems, wireless sensor networks, the evolution of technology, and some interest issues captured by CPSs. This survey also laid a great foundation for the later research. The work in [9] studied the multicast routing design for decentralized sensors and controllers in CPSs. The application of CPSs in intelligent transportation system is considered as a new research hotspot in the future. In [10], the authors extensively studied the CPSs in intelligent transportation system, and artificial systems, computational experiments, and parallel execution methodology are introduced based on data-driven model. For railway CPSs, Zhang proposed an aspect-oriented approach to modeling this system, whose velocity, flow, and density are dynamic and continuous changing with time during the process of communication and calculation [11]. For vehicular CPSs, defined as a combination of traditional state-based discrete control model and continuous models based on physical environments, the author provided adequate requirements specification for this specific system in [12]. In the literature, although there are few related research in maritime scenarios, the existing literatures such as [13] also provide us a myriad of hints. To our knowledge, this is the first work to study maritime CPSs.

We combine cooperative and cognitive technology to construct an innovative maritime communication network paradigm in this paper, that is, a Cooperative Cognitive Maritime wireless mesh/ad hoc network, to efficiently explore White Space (WS) on the sea and address the limited

transmission range and distorted signal transmission simultaneously. Although the research on maritime scenario is still in the early stage, the counterparts on land-based network could provide a solid foundation for our work [14–22]. In [23], the authors studied the broadcast scheduling issue with minimum latency and redundancy analysis for cognitive radio networks, which could significantly improve the existing performance of latency and redundancy for CRNs. In [24], the distributed data collection problem for asynchronous CRNs is studied considering the proper carrier-sensing range and the fairness issue. In [25], the sensing-throughput trade-off issue and the impact of different system parameters on the optimal sensing time for cooperative spectrum sensing in CRNs are investigated, based on the soft decision schemes. They also compared the performance in different relay modes. In cognitive system, not only SUs could cooperate with PUs, but also could cooperate with each other. Generally speaking, in the research of cooperative system, it is always assumed that relays would like to help source to transmit information. However, a challenge in practical maritime communication system is how to effectively build up a cooperative system in limited transmitter power level on the sea. On the other hand, nodes are independent selfish individuals and they are not voluntary to help other nodes. So, the user must pay fee to buy resources to attain corresponding help. Then, the other user calculates payoff of itself to decide whether to join the cooperative communication or not. It is inevitable to consume resources such as power and rate in helping other nodes to transmit; rational nodes are not obligatory to involve in this cooperate activity. Game theory is a good mathematical method to research effective cooperative condition and requirement. In [26], the authors investigated multichannel assignment in wireless sensor networks utilizing the game theoretic approach. In [27], the authors proposed adjustable price and rate algorithm based on cooperative stimulating method in ad hoc networks. One classic literature about channel allocation in cognitive radio networks applying Stackelberg game [28] achieved a minimum required SIR through setting a new potential utility function. However, it utilized the iteration scheme to get virtual price as cost of accessing channel. In [29], a joint pricing and power allocation strategy for dynamic spectrum access networks with stackelberg game model is proposed. However, it is about the resource allocation study between one pair of PU and SU. In [30], cooperative spectrum access of primary and secondary users, as well as MAC protocol for multichannel cognitive radio networks, is developed. In [31], the authors discussed a nonsymmetrical cooperative three-node model. But, in practice, the two users are likely the potential relays, which is better accordant with fairness of network. To address the resource allocation issue of this cooperative cognitive mesh/ad hoc maritime network, Stackelberg game with priority is employed between PU and SUs, while a symmetrical system model is considered among SUs-to-SUs in this paper.

3. System Model and WS Analysis

Cyber systems are embedded in all types of physical systems (we refer to them as PUs, SUs, and BS in this paper). Such

networks comprise sensors and on-board units installed in the vessels (PUs and SUs) as well as shore-side BS. The data collected from the sensors on the vessels (PUs) could be shown on the bridge for navigators, delivered by the SUs or BS, and finally received by the maritime authorities, ship owners, or shipping companies, and then a feedback is given to the controlling service to the vessels. This software-intensive system demonstrates very complex maritime CPSs with intricate interplay between the physical world and the cyber world. This section presents the details of the system model as well as the analysis of WS on the sea.

3.1. System Model. We consider the scenario that an innovative paradigm CCMCPSSs could be formed by exploring the connection between neighboring ships, sea farm, oil/gas platform, maritime security/safe monitors, and marine beacons and buoys. Depending on the node mobility, the network close to the shore could be the combination of the following two classes of network: (1) immobile marine infrastructure mesh network; (2) mobile ship-to-ship/shore mesh/ad hoc network. The nodes are equipped with a mesh/ad hoc module that is capable of implementing cognitive radio functions. The cognitive nodes that sense or collect data regularly could be the relays to deliver traffic. And then, the immobile or mobile nodes could connect to the land-based network via shore-based base station easily. In this cognitive system, PUs could be the radio devices installed near the coastal region such as authorities on land, or licensed vessels. Thereby, SUs could be the devices on the sea, such as the unlicensed vessels neighboring, sea farm, oil/gas platform, and marine beacons and buoys. SUs could cooperate with each other and form a cluster to cooperate with PUs to further strive for transmission opportunities. Therefore, the unlicensed users could utilize the unused frequency spectrum resources opportunistically. And sensing or monitoring data by the mesh/ad hoc nodes could be transmitted to the land-based administrative agencies through wireline/wireless network on land by dynamically adopting the unused frequency channels by PUs. Network topology is depicted in Figure 1.

As illustrated in Figure 2, PU would like to send data to base station (BS). When the direct channel quality is not very well or the PU has large capacity data to transmit, the shore infrastructure-based PU could cooperate with ad hoc nodes SUs as relays which adopt Decode-and-Forward (DF) mode to relay the message for PU. And the destination BS employs maximal-ratio combining (MRC) reception. We only consider the scenario of single channel; that is, one PU corresponds to the shore-based BS on one channel in time slot T . There are totally n SUs participating in the relaying process. The Time Division Multiple Access (TDMA) based media access control (MAC) protocol is employed. The portion of α is utilized for cooperative transmission, including the PU-to-SUs in the first half period $(\alpha/2)T$ and SUs-to-BS to relay the data from PU in the second half period. The remaining period $(1 - \alpha)T$ is exploited for SUs to transmit their own data to the corresponding secondary receivers as the compensation for relaying. The channel coefficients are independent and distributed with complex Rayleigh variables, receiving complex additive white Gaussian noises

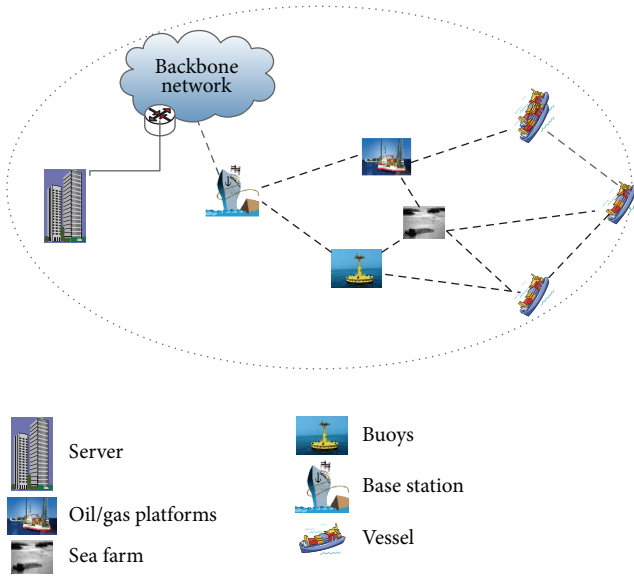


FIGURE 1: System model.

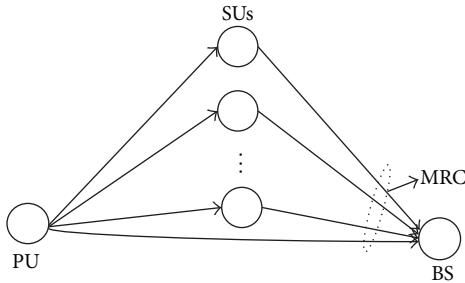


FIGURE 2: Cooperative model between PU and SUs.

(AWGN) with zero mean and the same variance N_0 . The channel coefficients between PU and BS, PU and SU_j , and SU_j and receivers accordingly are, respectively, $h^{p,b}$, $h_s^{s,b}$, and h_s^j . The total bandwidth allocated to PU is W . The power of PU is P_p . For simplicity, SU_j consumes the power P_s^j both for relaying transmission and their own transmission. The power of both the PU and SUs is under the constraint of P_{\max} .

3.2. The Analysis of WS on the Sea. There is abundant TVWS and cellular WS resources at sea nowadays, especially on the open ocean. The present maritime communication system, that is, GMDSS (Global Maritime Distress and Safety System) [32] partly occupies the frequency bands from 300 kHz to 300 MHz but does not utilize them efficiently. "Maritime WS" is nominated for the unused bands for maritime communication. Furthermore, the fundamental of maritime WS usage should be as follows. (1) The PUs such as maritime authorities or licensed vessels should be strictly protected avoiding harmful interference [33]. (2) The priority of communication at sea, that is, distress, urgency, safety, and routine, should be guaranteed. (3) The CPSs cognitive nodes at sea should have the capability of sensing the "Maritime WS" spectrum.

4. Game Theory Based Resource Allocation between PU and SUs

In this section, the resource allocation between PU and SUs will be discussed. Due to the selfishness of each entity, the optimization aims for PU and SUs are distinctive, where PU targets maximizing its throughput and SUs desire to obtain the bandwidth to deliver via cooperation. Stackelberg game is utilized here to model the cooperation process to achieve Nash equilibrium.

4.1. Utility of PU and SUs. We define the utility for both PU and SUs as the transmission rate via DF cooperative transmission, when MRC receiving mode is adopted at BS. The utility of PU is presented as follows:

$$U_p = \frac{\alpha W}{2} \log_2 \left[1 + \frac{P_p}{N_0} |h^{p,b}|^2 + \sum_{j=1}^n \frac{P_s^j}{N_0} |h_{s_j}^{s,b}|^2 \right]. \quad (1)$$

The utility of SU is described as follows:

$$U_s(\alpha) = (1 - \alpha) W \log_2 \left(1 + \sum_{j=1}^n \frac{P_s^j}{N_0} |h_s^j|^2 \right) - c \left(1 - \frac{\alpha}{2} \right) \sum_{j=1}^n P_s^j, \quad (2)$$

where c ($0 < c < 1$) is the weight of consumed energy in the overall utility [34]. Then, the game theoretic process is to determine the best access time portion α to SUs, as well as the consumed power P_s^j for SU_j .

4.2. Stackelberg Game Analysis. Stackelberg game is a sequential game with priority, that is, existing leader and follower, which could be estimated by *Backward induction method*. It consists of two stages. In the first stage, we determine the best choice of the follower (SUs), by making a hypothesis that the scheme of leader (PU) is prefixed. Then, in the second stage, the optimal strategy of PU is determined formally, on the basis of the optimal schemes of SUs. We give a clear explanation again that the payment resources for the two parties to join this Stackelberg game are access time portion α of PU and transmitting power P_s^j of SUs.

4.2.1. SUs Payment Selection Game

Definition 1. Denote $P_s^{*j}(\alpha)$ to be the optimal payment selection of SUs, that is, optimal transmitting power of SUs, such that the utility could reach to the maximum value when $P_s^{*j}(\alpha)$ is chosen, for any given α ($0 < \alpha < 1$), $U_s(P_s^{*j}(\alpha), \alpha) \geq U_s(P_s^j(\alpha), \alpha)$.

Theorem 2. The optimal payment selection $P_s^{*j}(\alpha)$ of SUs, that is, optimal transmitting power ($P_s^{*j}(\alpha)$) = $\arg \max U_p(\alpha)$, is

given by the following equation, when the PU chooses a certain time allocation coefficient α for cooperation:

$$P_s^{*j}(\alpha) = \min \left\{ \left\{ \frac{(1-\alpha)W}{cn(1-\alpha/2)\ln(2)} - \frac{N_0}{n \sum_{j=1}^n |h_s^j|^2} \right\}, P_{\max} \right\}. \quad (3)$$

Proof. As described above, we first suppose that the PU allocates α to SUs for cooperation. Then, the SUs payment selection game is to maximize the utility by choosing the optimal transmitting power P_s^j , which is formulated as follows:

$$\begin{aligned} \max_{P_s^j} \quad & U_s(\alpha) \\ & = (1-\alpha)W \log_2 \left(1 + \sum_{j=1}^n \frac{P_s^j}{N_0} |h_s^j|^2 \right) \\ & \quad - c \left(1 - \frac{\alpha}{2} \right) \sum_{j=1}^n P_s^j \end{aligned} \quad (4)$$

$$\text{s.t.} \quad 0 \leq P_s^j \leq P_{\max}.$$

Taking the first-order partial derivative of the above SU utility function with respect to $\sum_{j=1}^n P_s^j$, we have

$$\begin{aligned} \frac{\partial U_s}{\partial \sum_{j=1}^n P_s^j} &= \frac{n(1-\alpha) \sum_{j=1}^n |h_s^j|^2}{\left(1 + \sum_{j=1}^n P_s^j |h_s^j|^2 / N_0 \right) N_0 \ln 2} \\ &\quad - nc \left(1 - \frac{\alpha}{2} \right). \end{aligned} \quad (5)$$

Let $\partial U_s / \partial \sum_{j=1}^n P_s^j = 0$, and the optimal transmitting power $\sum_{j=1}^n P_s^j$ could be achieved as follows:

$$\sum_{j=1}^n P_s^{*j}(\alpha) = \frac{(1-\alpha)W}{c(1-\alpha/2)\ln(2)} - \frac{N_0}{\sum_{j=1}^n |h_s^j|^2}. \quad (6)$$

For simplicity, we assume all the transmitting power P_s^{*j} employs the average value of $\sum_{j=1}^n P_s^{*j}(\alpha)$. Then, we get

$$\begin{aligned} P_s^{*j}(\alpha) &= \sum_{j=1}^n P_s^{*j}(\alpha) \cdot \frac{1}{n} \\ &= \frac{1}{n} \left\{ \frac{(1-\alpha)W}{c(1-\alpha/2)\ln(2)} - \frac{N_0}{\sum_{j=1}^n |h_s^j|^2} \right\}. \end{aligned} \quad (7)$$

Considering the power constraint, the SU payment selection scheme P_s^{*j} is presented as

$$\begin{aligned} P_s^{*j}(\alpha) &= \min \left\{ \left\{ \frac{(1-\alpha)W}{cn(1-\alpha/2)\ln(2)} - \frac{N_0}{n \sum_{j=1}^n |h_s^j|^2} \right\}, P_{\max} \right\}. \end{aligned} \quad (8)$$

□

The first-order derivative of the above expression with respect to α is $- \alpha W / cn(\alpha - 2)^2 \ln(2)$, which is definitely negative, so the SU's utility has the optimal value.

4.2.2. Maximizing PU's Utility. We maximize the PU's utility on the basis of having the knowledge of SUs' optimal payment choice. Then, the optimization issue could be formulated as

$$\begin{aligned} \max_{\alpha} \quad & U_p \\ & = \frac{\alpha W}{2} \log_2 \left[1 + \frac{P_p}{N_0} |h^{p,b}|^2 + \sum_{j=1}^n \frac{P_s^j}{N_0} |h_s^{s,b}|^2 \right] \end{aligned} \quad (9)$$

$$\text{s.t.} \quad 0 < \alpha < 1.$$

Definition 3. Denote α^* to be the optimal payment selection of PU, that is, optimal time allocation coefficient of PU, such that the utility of PU could reach to the maximum value when α^* is chosen for cooperation.

Theorem 4. The optimal payment selection α^* of time allocation coefficient for PU to cooperate, that is, $(\alpha^*) = \arg \max U_p(\alpha)$, is given by the following equation:

$$\alpha^* = \frac{(6B + \ln(A+1) + A^2 \ln(A+1) + 4B^2 + 2A \ln(A+1) + 6AB)}{(8B + 4B^2 + 8AB)}. \quad (10)$$

Here, $(P_p/N_0)|h^{p,b}|^2 = A$, $(W/cn \ln 2 N_0) \sum_{j=1}^n |h_s^{s,b}|^2 = B$.

Proof. In order to solve this optimization problem, we apply the optimization results of P_s^{*j} obtained in formulation

(7) into the utility function of PU. In the second item of formulation (7), the denominator $n \sum_{j=1}^n |h_s^j|^2$ is much more larger than the numerator N_0 , so the second item could be ignored for convenient calculation. Then, we substitute $P_s^{*j}(\alpha) = (1 - \alpha)W / cn(1 - \alpha/2) \ln(2)$ into the utility function of the PU and get

$$U_p(\alpha) = \frac{\alpha W}{2} \log_2 \left[1 + \frac{P_p}{N_0} |h^{p,b}|^2 + \frac{1}{N_0} \sum_{j=1}^n \frac{(1 - \alpha) W}{cn \ln 2 (1 - \alpha/2)} |h_{sj}^{s,b}|^2 \right]. \quad (11)$$

In order to calculate conveniently, we set $(P_p / N_0) |h^{p,b}|^2 = A$, $(W / cn \ln 2 N_0) \sum_{j=1}^n |h_{sj}^{s,b}|^2 = B$. Then, for $U_p(\alpha)$, we do the Taylor series expansion with respect to α ,

$$\begin{aligned} U_p(\alpha) &= \frac{(w \ln(A + 1))}{(2 \ln 2)} \\ &+ \frac{(w(l(A + 1) - (2B)/(A + 1)) \times (\alpha - 1))}{(2 \ln 2)} \\ &- \frac{(w((2B^2)/(A + 1)^2 + 4B/(A + 1))(\alpha - 1)^2)}{2 \ln 2}. \end{aligned} \quad (12)$$

Taking the first-order derivative of $U_p(\alpha)$ regarding α and letting $\partial U_p / \partial \alpha = 0$ to obtain the optimal value of α , we have

$$\alpha^* = \frac{(6B + \ln(A + 1) + A^2 \ln(A + 1) + 4B^2 + 2A \ln(A + 1) + 6AB)}{(8B + 4B^2 + 8AB)}. \quad (13)$$

The optimal value α^* is obviously pertaining to (0, 1). \square

5. Game Theory Based Resource Allocation between SUs

In Section 4, we focus on the game between PU and SUs in cognitive mesh/ad hoc network. Actually, the scheme is that SUs cooperate with each other to thereby form a cluster firstly, and then the cluster cooperates with PU further to attain the transmission opportunities as the benefits. However, the SUs are independent and selfish such that they are not voluntary to help others with no profit. Thereby they should pay fees to buy resource to obtain transmission assistance. And the users could count for the payoff to determine whether to cooperate with each other or not. A symmetrical system model is exploited, and a price game based on payoff function is developed. Then, the game theoretic procedure converging to Nash equilibrium is described further.

5.1. Game Theory Model. In the cooperative symmetry, with triple nodes, source node s , relay node r , and destination node d are players. Strategies space is stimulating prices μ and ν , which are the price paid to the partner, while R_{rs} and R_{sr} are defined as the transmitting speed for the counterpart. Payoff function is defined as the difference of utility function and price function; that is, $\text{Payoff} = U - P$. Denote U as utility function and denote P as price function.

(i) *Utility Function.* We utilize a common utility function described as the received data when consuming one-unit energy [35]. The unit throughput is defined as 1/bit. R_s and ber_{sd} are denoted as transmitting speed and bit error rate (BER) of source node. R_r and ber_{rd} are, respectively, denoted as the transmitting speed and BER of relaying node. T is transmitting duration.

Then, the utility of noncooperative source node is

$$U_{s,\text{non}} = 1 \cdot \text{Throughput} = R_s \cdot T \cdot (1 - \text{ber}_{sd}). \quad (14)$$

The utility of cooperative source node is

$$U_{s,\text{coop}} = 1 \cdot \text{Throughput}_{\text{local}} + 1 \cdot \text{Throughput}_{\text{relay}}. \quad (15)$$

The utility of noncooperative relay node is

$$U_{r,\text{non}} = 1 \cdot \text{Throughput} = R_r \cdot T \cdot (1 - \text{ber}_{rd}). \quad (16)$$

The utility of cooperative relay node is

$$\begin{aligned} U_{r,\text{coop}} &= 1 \cdot \text{Throughput} \\ &= (R_r - R_{rs}) \cdot T \cdot (1 - \text{ber}_{rd}) + R_{sr} T \\ &\quad \cdot (1 - \text{ber}_{sr}). \end{aligned} \quad (17)$$

(ii) *Price Function.* The service price is denoted as $\lambda_i = \lambda_0(R_i / R_{\max})$, where λ_0 is the criterion price. R_i and R_{\max} are indicated as transmitting rate and maximum rate, respectively. Denote λ_s as the unit price of source and λ_r as the unit price of relay; λ_{rs} is shown as the unit price of relay node in helping source node to send data; λ_{sr} is the unit price of source in helping relay to transmit data; μ and ν are the stimulating price of source and relay, respectively; λ_{r-rs} is the unit price of relay node with cooperation; λ_{s-sr} is the unit price of source node with cooperation.

The price function of noncooperative source node is

$$P_{s,\text{non}} = \lambda_s \cdot R_s \cdot T \cdot (1 - \text{ber}_{sd}). \quad (18)$$

The price function of cooperative source node is

$$\begin{aligned} P_{s,\text{coop}} &= \lambda_{s-sr} \cdot (R_s - R_{sr}) \cdot T \cdot (1 - \text{ber}_{sd}) \\ &+ (\mu + \lambda_{rs}) \cdot R_{rs} \cdot T \cdot (1 - \text{ber}_{rs}) - \nu \cdot R_{sr} \cdot T \\ &\quad \cdot (1 - \text{ber}_{sr}). \end{aligned} \quad (19)$$

The price function of noncooperative relay node is

$$P_{r_non} = \lambda_r \cdot R_r \cdot T \cdot (1 - \text{ber}_{rd}). \quad (20)$$

The price function of cooperative relay node is

$$\begin{aligned} P_{r_coop} &= \lambda_{r-sr} \cdot (R_r - R_{rs}) \cdot T \cdot (1 - \text{ber}_{rd}) \\ &+ (\nu + \lambda_{sr}) \cdot R_{sr} \cdot T \cdot (1 - \text{ber}_{sr}) - \mu \cdot R_{rs} \\ &\cdot T \cdot (1 - \text{ber}_{rs}). \end{aligned} \quad (21)$$

(iii) *Payoff Function.* Payoff function is defined as the difference between utility function and price function; that is, $\text{Payoff} = U - P$ (U is utility function and P is price function).

The payoff function of noncooperative source node is

$$\text{payoff}_{s_non} = (1 - \lambda_s) \cdot R_s \cdot T \cdot (1 - \text{ber}_{sd}). \quad (22)$$

The payoff function of cooperative source node is

$$\begin{aligned} \text{payoff}_{s_coop} &= (1 - \lambda_{s-sr}) \cdot (R_s - R_{sr}) \cdot T \cdot (1 - \text{ber}_{sd}) \\ &+ \nu \cdot R_{sr} \cdot T \cdot (1 - \text{ber}_{sr}) \\ &+ (1 - \mu - \lambda_{rs}) \cdot R_{rs} \cdot T \cdot (1 - \text{ber}_{rs}). \end{aligned} \quad (23)$$

The payoff function of noncooperative relay node is

$$\text{payoff}_{r_non} = (1 - \lambda_r) \cdot R_r \cdot T \cdot (1 - \text{ber}_{rd}). \quad (24)$$

The payoff function of cooperative relay node is

$$\begin{aligned} \text{payoff}_{r_coop} &= (1 - \lambda_{r-sr}) \cdot (R_r - R_{rs}) \cdot T \cdot (1 - \text{ber}_{rd}) \\ &+ \mu \cdot R_{rs} \cdot T \cdot (1 - \text{ber}_{rs}) \\ &+ (1 - \nu - \lambda_{sr}) \cdot R_{sr} \cdot T \cdot (1 - \text{ber}_{sr}). \end{aligned} \quad (25)$$

5.2. Game Theory Analysis. The game theoretic procedure converging to Nash equilibrium by the iteration algorithm is described in this section. Specifically, the value boundary of u and v is induced. The value boundary of relaying rate R_{rs} and R_{sr} is also obtained.

5.2.1. Value Boundary Analysis

Lemma 5. Using game theory converging into Nash equilibrium, the stimulating price of source and relay u and v , as well as R_{rs} and R_{sr} , is obtained.

Proof. (i) *Value Boundary of u and v .* With cooperation, according to the rules that $\text{Payoff}_{s_coop} - \text{Payoff}_{s_non} \geq 0$ and $\text{Payoff}_{r_coop} - \text{Payoff}_{r_non} \geq 0$, we have

$$\begin{aligned} u &\leq 1 - \lambda_s + \frac{\nu R_{rs} (1 - \text{ber}_{sr}) + \lambda_s \cdot R_s (1 - \text{ber}_{sd})}{R_{rs} (1 - \text{ber}_{rs})} \\ &- \frac{R_{sr} (1 - \text{ber}_{sd}) - \lambda_{s-sr} (R_s - R_{sr}) (1 - \text{ber}_{sd})}{R_{rs} (1 - \text{ber}_{rs})}, \end{aligned} \quad (26)$$

$$\begin{aligned} u &\geq \frac{[R_{rs} + \lambda_{r-sr} (R_r - R_{rs})] (1 - \text{ber}_{rd})}{R_{sr} (1 - \text{ber}_{sr})} \\ &+ \frac{(\nu + \lambda_{sr} - 1) R_{sr} (1 - \text{ber}_{sr})}{R_{sr} (1 - \text{ber}_{sr})}, \end{aligned} \quad (27)$$

$$\begin{aligned} v &\leq \frac{[\lambda_r - R_{rs} - \lambda_{r-sr} (R_r - R_{rs})] (1 - \text{ber}_{rd})}{R_{sr} (1 - \text{ber}_{sr})} \\ &+ \frac{\mu R_{rs} (1 - \text{ber}_{rs})}{R_{sr} (1 - \text{ber}_{sr})} 1 - \lambda_{sr}, \end{aligned} \quad (28)$$

$$\begin{aligned} v &\geq [R_{sr} + \lambda_{s-sr} (R_s - R_{sr}) - \lambda_s R_s] (1 - \text{ber}_{sd}) \\ &- (1 - u - \lambda_s) R_{rs} (1 - \text{ber}_{rs}). \end{aligned} \quad (29)$$

Jointly considering the formulas (26)–(29), μ_{\min} and ν_{\min} could be determined simultaneously.

(ii) *Value Boundary of R_{rs} and R_{sr} .* payoff_{r_coop} is a quadratic function with respect to R_{rs} . When u is determined, the peak value of R_{rs}^* could be decided via taking first-order partial derivative of R_{rs} as follows:

$$\begin{aligned} \frac{\partial \text{payoff}_{r_coop}}{\partial R_{rs}} &= -T (1 - \text{ber}_{rd}) \\ &+ \frac{2\lambda_0 (R_r - R_{rs}) T (1 - \text{ber}_{rd})}{R_{\max}} \\ &+ uT (1 - \text{ber}_{rs}). \end{aligned} \quad (30)$$

When $\partial \text{payoff}_{r_coop} / \partial R_{rs} = 0$, we have

$$R_{rs}^* = R_r - R_{rs \max} \left[\frac{1}{2\lambda_0} - \frac{\mu (1 - r_{rs})}{2\lambda_0 (1 - \text{ber}_{rd})} \right]. \quad (31)$$

Similarly, payoff_{s_coop} is a quadratic function with respect to R_{sr} . When v is determined, the peak value of R_{sr}^* could be obtained via taking the first-order partial derivative of R_{sr} as follows:

$$\begin{aligned} \frac{\partial \text{payoff}_{s_coop}}{\partial R_{sr}} &= (R_s - R_{sr}) T (1 - \text{ber}_{sd}) \\ &- \frac{\lambda_0}{R_{\max}} (R_s - R_{sr})^2 T (1 - \text{ber}_{sd}) \\ &+ (1 - \mu - \lambda_s) R_s T (1 - \text{ber}_{rs}) \\ &+ \nu R_{sr} T (1 - \text{ber}_{sr}). \end{aligned} \quad (32)$$

When $\partial \text{payoff}_{s,\text{coop}} / \partial R_{sr} = 0$, we have

$$R_{sr}^* = R_s - R_{sr} \max \left[\frac{1}{2\lambda_0} - \frac{\nu(1-r_{sr})}{2\lambda_0(1-\text{ber}_{sd})} \right]. \quad (33)$$

As we see, R_{rs}^* is just relating to μ but not ν ; R_{sr}^* is just relating to ν but not μ . \square

5.3. Game Theoretic Procedure. This issue is a two-party game with priority decision. Based on the above derivation, $\text{payoff}_{s,\text{coop}}$ is a monotone decreasing function with respect to μ and monotone increasing function with respect to ν . $\text{payoff}_{r,\text{coop}}$ is a monotone decreasing function with respect to ν and monotone increasing function with respect to μ . Therefore, we could conclude that all the two parties hope the counterpart pays for higher simulating price and vice versa. With cooperation, they would like to choose μ_{\min} and ν_{\min} . Then, the cooperation transmitting rate paid for the counterpart could be calculated by formulation (31)–(33). If they are lower than the maximum transmitting rate, the game theoretic procedure will be finished and the cooperation process will be initiated. If they do not meet the requirement, the above procedure will be revalued and repeated

$$\begin{aligned} \text{payoff}_{\text{total-coop}} &= \text{payoff}_s + \text{payoff}_r + \text{payoff}_{\text{BS}} \\ &= \text{payoff}_s + \text{payoff}_r + P_{\text{total}} \\ &= \text{payoff}_s + \text{payoff}_r + P_s + P_r \\ &= (U_s - P_s) + (U_r - P_r) + P_s + P_r = U_s + U_r, \\ \text{payoff}_{\text{total-non}} &= (1 - \lambda_s) R_s T (1 - \text{ber}_{sd}) \\ &\quad + (1 - \lambda_r) R_r T (1 - \text{ber}_{rd}), \\ \text{payoff}_{\text{total-coop}} &= (1 - \lambda_{s-sr}) (R_s - R_{sr}) T (1 - \text{ber}_{sd}) \\ &\quad + (1 - \mu - \lambda_s) R_{rs} T (1 - \text{ber}_{rs}) \\ &\quad + \nu R_{sr} T (1 - \text{ber}_{sr}) \\ &\quad + (1 - \lambda_{r-rs}) (R_r - R_{rs}) T (1 - \text{ber}_{rd}) \\ &\quad + (1 - \nu - \lambda_{sr}) R_{sr} T (1 - \text{ber}_{sr}) \\ &\quad + \mu R_{rs} T (1 - \text{ber}_{rs}) \\ &= (1 - \lambda_s) R_s T (1 - \text{ber}_{rs}) \\ &\quad + (1 - \lambda_{sr}) R_{sr} T (1 - \text{ber}_{sr}) \\ &\quad + (1 - \lambda_{s-sr}) (R_s - R_{sr}) T (1 - \text{ber}_{sd}) \\ &\quad + (1 - \lambda_{r-rs}) (R_r - R_{rs}) T (1 - \text{ber}_{rd}). \end{aligned} \quad (34)$$

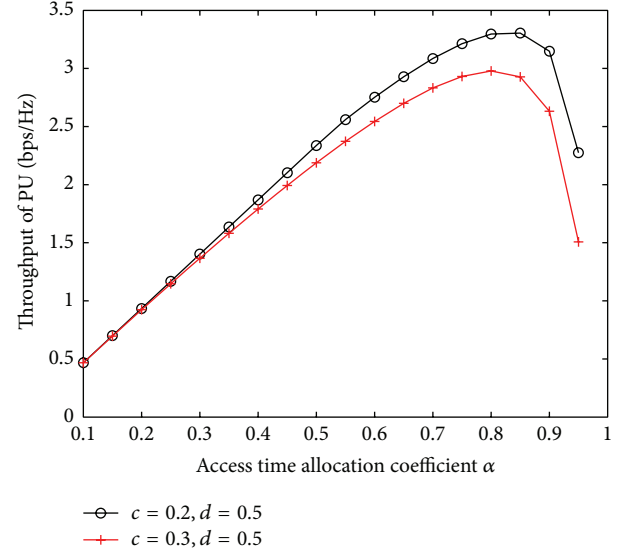


FIGURE 3: Throughput of PU versus time allocation coefficient α .

Assuming $r_r = r_{rs}$ and $r_s = r_{sr}$, the above formulation could be approximated as

$$\begin{aligned} \text{payoff}_{\text{total-coop}} &= (\lambda_{r-rs} - \lambda_s) R_{rs} T (1 - \text{ber}_{rd}) \\ &\quad + (\lambda_{s-sr} - \lambda_{sr}) R_{sr} T (1 - \text{ber}_{sd}) \\ &\quad + (1 - \lambda_{s-sr}) R_s T (1 - \text{ber}_{sd}) \\ &\quad + (1 - \lambda_{r-rs}) R_r T (1 - \text{ber}_{rd}). \end{aligned} \quad (35)$$

When $\lambda_{r-rs} - \lambda_s > 0$ and $\lambda_{s-sr} - \lambda_{sr} > 0$, the Nash equilibrium result (R_{rs}^*, μ_{\min}) , (R_{sr}^*, ν_{\min}) can satisfy Pareto optimum along with the increase of R_{rs} and R_{sr} .

6. Simulation Results

In the simulation, WS in VHF band spectrum is considered. Similar to [30], we normalize the distance between PU and BS, and SU is located at the distance $d_s \in (0, 1)$ from PU and $1 - d_s$ from the BS. The channel gains are in accordance with simple path loss model with coefficient $\xi = 3.5$. The maximum secondary transmission power P_{\max} is normalized to 1.

Figure 3 presents the throughput of PU on a certain channel versus the access time allocation coefficient α . The parameter of weights c is set as 0.2 and 0.3 for $d = 0.5$, respectively. The normalized distance between the SU and its corresponding receiver is set to 0.8. Furthermore, we do some changes on the parameter chosen. Through the analysis of output, we could find that the resulting graphs will become irregular if the corresponding parameters become greater or smaller. For example, if we take into consideration of the limit value parameters (e.g., d gets 0.1, the output will be closed to a straight line), some simulation results will become meaningless. Then, we determine the best optimal parameters to outperform the simulation. From Figure 3, it can be seen that there exists an optimal α to maximize the

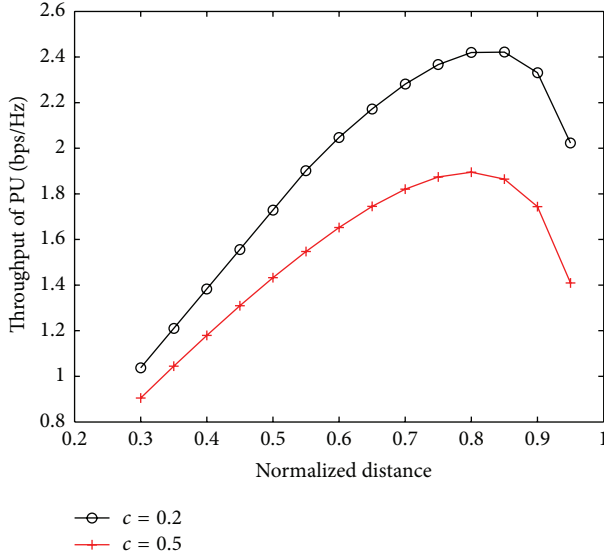


FIGURE 4: Throughput of PU versus the normalized distance.

throughput via cooperative transmission. Furthermore, the throughput of PU is higher with a smaller weight c . This is because the fact that the SU pays more attention to the throughput than the energy consumption. When SU wishes to consume more power to achieve cooperation; PU could attain more throughput accordingly. Figure 4 shows PU's throughput versus the normalized distance d_s for c setting to be 0.2 and 0.5, respectively. It is obvious that PU can achieve a higher throughput via cooperation than direct transmission, when appropriate d_s is chosen.

With the cooperation among SUs, we set the parameters as follows. The coordination of destination is (0, 0) and the source is (4, 0), with the unit coordinate of nautical mile. It is easy to observe that, if we set relay (x , y), the horizontal coordination and vertical coordination are varied in the range of (1, 10). The transmitting time T is 10 ms, and the λ_0 criterion price is 0.1. The transmitting rate of source and relay is all 320 Mb/s.

The simulation results show the comparison of total payoff, source payoff, and relay payoff, when axes x and y , respectively, indicate horizontal and vertical coordination. From Figures 5, 6, and 7, it could be seen that the payoff of cooperative scheme is obviously increased than noncooperative scheme. Along with the increased distance between relay and destination, the difference of the total payoff as well as relay payoff between cooperative and noncooperative becomes larger. But the payoff of source is very close between cooperation and noncooperation. It could be explained that it does not achieve the best offset from payoff fee in cooperation. Meanwhile, the payoff of source could be increased if the value of μ is larger.

7. Conclusion

In this paper, an innovative paradigm CCMCPSSs is developed to provide high-speed and low-cost communications for maritime community. We focus on the resource allocation

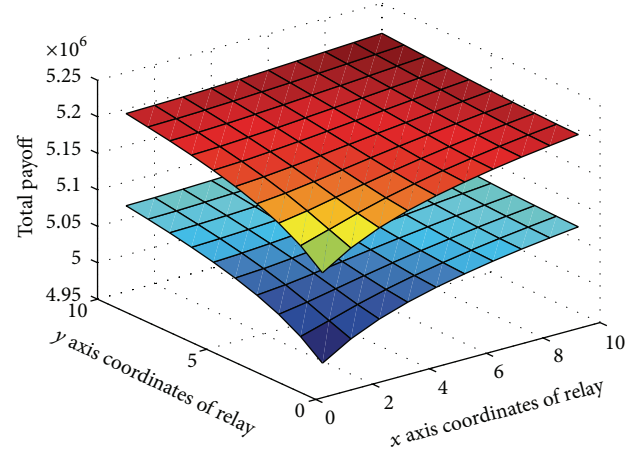


FIGURE 5: The total payoff comparison between cooperative and noncooperative schemes.

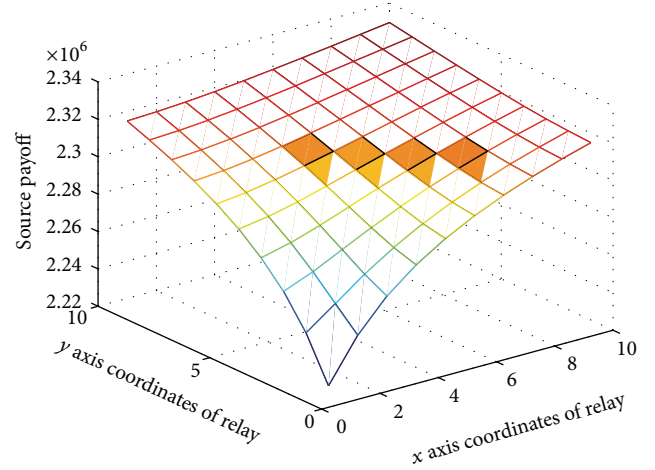


FIGURE 6: The source payoff comparison between cooperative and noncooperative schemes.

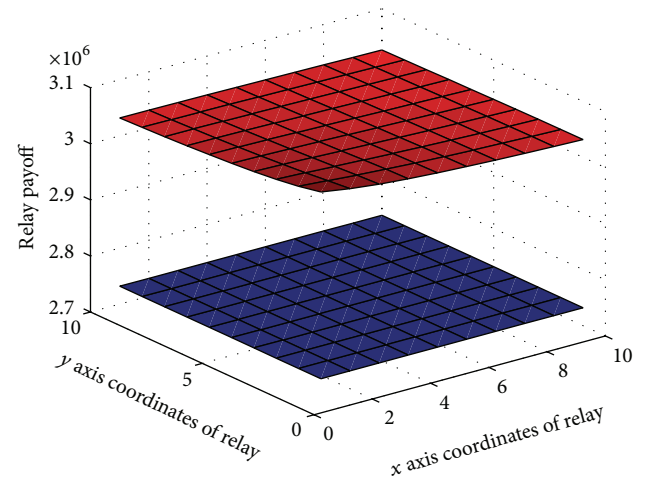


FIGURE 7: The relay payoff comparison between cooperative and noncooperative schemes.

between the two stages, that is, PU-to-SUs and SUs-to-SUs based on the game theory. In the first stage, Stackelberg game, which is a sequential game with priority, is introduced to obtain the optimal resource allocation between PU and SUs. In the second stage, SUs-to-SUs, an symmetrical system model is considered, while a price game based on the payoff function is proposed. Then, the game theoretic procedure that converges to Nash equilibrium is illustrated. Simulation results indicate that our proposed schemes could effectively increase the throughput as well as the payoffs of the system. For future work, the combination of the designing of MAC layer and resource allocation will be further considered.

Notations

α :	The portion of time slot for cooperative transmission
α^* :	The optimal payment selection of PU
λ_i :	The service price
λ_0 :	The criterion price
ber_{sd} :	The bit error rate of source node
$h^{p,b}(h_{s_j}^{s,b})(h_s^j)$:	The channel coefficients between PU and BS (PU and SU _j) (SU _j and receivers)
N_0 :	The variance of channel coefficients
P_s^j :	The power of SU _j for both relaying transmission and own transmission
$P_s^{*j}(\alpha)$:	The optimal payment selection of SUs
P_{\max} :	The power constraint of both the PU and SUs
$(R_{rs\max}^*, \mu_{\min})$:	The Nash equilibrium result that satisfies Pareto optimum of R_{rs}
$(R_{sr\max}^*, \nu_{\min})$:	The Nash equilibrium result that satisfies Pareto optimum of R_{sr}
T :	The time slot
$U_p(U_s(\alpha))$:	The utility of PU and SU
W :	The total bandwidth allocated to PU.

Disclosure

This work was partly presented at IEEE WASA 2015 [36].

Conflict of Interests

The authors declare that there is no conflict of interests regarding the publication of this paper.

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