Optimization of Bus Bridging Service under Unexpected Metro Disruptions with Dynamic Passenger Flows

Jiadong Wang1, Zhenzhou Yuan1, and Yonghao Yin2

1Key Laboratory of Transport Industry of Big Data Application Technologies for Comprehensive Transport, Ministry of Transport, Beijing Jiaotong University, Beijing 100044, China
2State Key Lab of Rail Traffic Control and Safety, Beijing Jiaotong University, Beijing, 100044, China

Correspondence should be addressed to Jiadong Wang; 16124210@bjtu.edu.cn

Received 1 April 2019; Accepted 1 July 2019; Published 24 July 2019

Academic Editor: Francesco Galante

Copyright © 2019 Jiadong Wang et al. This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.

A metro disruption is a situation where metro service is suspended for some time due to unexpected events such as equipment failure and extreme weather. Metro disruptions reduce the level of service of metro systems and leave numerous passengers stranded at disrupted stations. As a means of disruption management, bus bridging has been widely used to evacuate stranded passengers. This paper focuses on the bus bridging problem under operational disruptions on a single metro line. Unlike previous studies, we consider dynamic passenger flows during the disruption. A multi-objective optimization model is established with objectives to minimize total waiting time, the number of stranded passengers and dispatched vehicles with constraints such as fleet size and vehicle capacity. The NSGA-II algorithm is used for the solution. Finally, we apply the proposed model to Shanghai Metro to access the effectiveness of our approaches in comparison with the current bridging strategy. Sensitivity analysis of the bus fleet size involved in the bus bridging problem was conducted.

1. Introduction

Metro disruptions, usually caused by emergency events such as infrastructure blockages, accidents and extreme weather, have frequently occurred in lots of cities [1]. A disruption is a relatively large external incident leading to large delays or cancellations of a number of trains in the timetable [2]. When the disruption lasts for a long time, for the purpose of reducing the impacts of disruptions, the affected metro line tends to operate in a short turning mode, on residual railway sections beyond the track crossovers. Metro disruptions prevent trains operate according to planning timetables. In addition, disrupted stations are unable to provide services to passengers resulting that numerous passengers are stranded. Pender et al. [3] found through the survey that bus bridging is the most common response to metro disruption especially line blockages. As an alternative mode for metro, a current approach to bus bridging is that buses travel parallelly to the disrupted metro line. However this approach usually unlikely achieves satisfactory effects. Due to the limited capacity of public transport, passengers at disrupted stations especially intermediate stations need to wait for several buses to get on the bus.

When operational disruptions occur, the most commonly used disruption management strategies mainly include adjusting timetable and rescheduling rolling stock. Considering an over-crowded passenger flow following a disruption, Gao et al. [4] proposed a stop-skipping strategy and built an optimization model to reschedule a metro train. Since metro lines are generally equipped with storage lines and backup trains, Yin et al. [5] proposed a MILP model for rescheduling metro trains and timetable by using backup trains. In addition to timetables, operational disruptions will also affect rolling stock plans. Veelenturf et al. [6] proposed a disruption management method integrating rescheduling of the timetable and rolling stock by taking dynamic passenger demand into account. Considering that passengers affected by the disruptions may change their travel plans [1], Cadarso et al. [7] proposed a two-stage approach to combine the timetable and rolling stock with passengers’ behavior in an integrated optimization model.
When the disruption lasts for a long time, for example, 30 minutes or more, it is unlikely to clear up the stranded passengers by adjusting the metro system itself. In this situation the bus bridging service is required. Kepaptsoglou and Karlaftis [8] first defined the bus bridging as the temporary routes to “bridge” disrupted stations and proposed a methodology for designing a bus bridging network. Codina et al. [9] established a nonlinear programming model to optimize a congested bus bridging system. The output included frequencies and the number of bus for each candidate bus lines. Jin et al. [10] proposed a systematic approach to optimizing bus bridging services. Firstly a column generation algorithm was used to generate candidate bus routes, then a path-based multimmodity network flow model was established to optimize bus bridging network. Similarly, Van der Hurk et al. [11] proposed a formulation to minimize passenger inconvenience cost including transfers and waiting cost with a minimal frequency constraint. Deng et al. [12] developed a model to design bus bridging routes considering passenger route choice behavior and the bridging station capacity.

The above research focused on a frequency-based modeling framework with the assumption that buses run on routes with fixed headway, which is similar to the traditional transit network design problem (TNDP) [13]. Nevertheless, when the metro service is disrupted, the transit agency must arrange the bus bridging service as soon as possible. Thus, it is necessary to take into consideration time spent required to dispatch buses. Recently, bus bridging operation strategies proposed to deal with unexpected disruptions did not require buses to operate based on fixed frequencies [1]. They proposed a type of models based on Bus Evacuation Problem (BEP) [14]. For example, Wang et al. [15] established a feeder-bus dispatch planning model to minimize the total transit operation cost in the evacuation. Hu et al. [16] proposed an efficient bus bridging strategy and constructed a multi-circle bus dispatching model minimizing the total evacuation time without taking dynamic passenger demands into account. Considering uncertainties under unexpected situation, Lv et al. [17] proposed a bus evacuation model based on interval chance-constrained integer programming method. The above research designated buses to shuttle between fixed stations resulting that the operation flexibility and evacuation efficiency may be limited. Recently, Gu et al. [1] proposed a flexible strategy allocating and scheduling buses to different bridging routes. They formulated a two-step model to optimize the bus bridging schedule minimizing total evacuation time and passenger delay, respectively. Although dynamic passenger demand was taken into account with a rolling horizon approach, the essence is taking static passenger flows at the end of each rolling period as the input of the proposed model.

Some other researchers studied travel behavior of passengers and travel demand after disruptions, which could provide references for this study [18–20]. For example, Wang et al. [18] studied the demand modelling of affected passengers and formulated it as a bulk queuing problem with the theory of stochastic process. The results showed the bridging demand continued for hours. Yin et al. [19] proposed a three-layer discrete choice behavior model to predict dynamic passenger flows. It was found that as time passed, an increasing number of passengers would alter their planned routes or quit metro journeys during the disruption. As a result, there are continuously passengers who transfer to buses during the disruption. Duan et al. [20] conducted studies using mobile phone data to investigate the impact on passengers of the collision accident of Metro Line 10 in Shanghai, China, on 27 September 2011. The result revealed that in the second evacuation stage over 60% of stranded passengers left the accident stations by ground transportation. They suggested that buses should be dispatched to connect disrupted stations, especially the non-interchange stations to the nearest operating stations. Related work has also been done to integrate bus bridging with other aspects. For example, Yang et al. [21] develop a joint approach to integrate passenger flow control and bus bridging to alleviate over-saturated situation for commuting metro lines. Kang et al. [22] coordinated last train timetabling and bus bridging service to deal with the stranded passengers. Zhang et al. [23] studied the optimal bus bridging starting time under uncertain recovery time, considering the uncertainty of the recovery time, but did not give a specific vehicle scheduling scheme.

In this research, we investigate how to efficiently operate bus bridging service in response to a metro operational disruption. Based on the strategy proposed by Gu et al. [1], a multi-objective programming model was formulated. Compared with previous studies, contributions of our research are summarized as follows: (i) an optimization-based model was proposed considering dynamic passenger flow demands during the disruption. (ii) The model was formulated with the objective to minimize total waiting time, the number of stranded passengers and dispatched buses, rather than evacuation time. (iii) A large-scale case study was conducted to evaluate the method. The remainder of this paper is organized as follows. A detailed description of bus bridging problem under unexpected metro disruptions is presented in Section 2. An optimization model is formulated in Section 3. A solution approach based on NSGA-II algorithm is described in Section 4. Benefits of our bus bridging strategy are analyzed with a large-scale case study in Section 5. Finally, conclusions and future research outlook are summarized in Section 6.

2. The Problem Description

2.1. Problem Description. Consider a metro disruption that several stations suspend service between two stations with track crossovers (see Figure 1.) Trains operate in short turning mode on residual metro lines. Buses need to be dispatched for emergency evacuation from bus depots and to evacuate the passengers to the turnover stations.

We suppose that the bus bridging stops are near metro station exits. Bus bridging stops are set on both sides of the roads corresponding to opposite directions of the same metro station. Assume that metro operators can predict disruption recovery time precisely and start bus bridging service immediately after disruption occurs. The metro line returns to normal operation at time $T$. Let $[0, T]$ denote
the disruption duration period. When the metro disruption resumes, passengers return to the metro system and will be rapidly evacuated. Therefore, we limit the effective time of bus bridging to \([0, T]\). In order to represent dynamic passenger flows, we divide \([0, T]\) equally into several time intervals of one minute. We hereafter refer to time interval as time and use \(t\) to index it. The passenger demand at station \(i\) at time \(t\) is denoted as \(\lambda_i(t)\). Let \(M\) be the set of buses available. Let \(D\) be the set of depots. Let \(N\) be the set of stations, including the set of disrupted stations \(I\) and the set of turnover stations \(J\). We desire to dispatch buses in the following manner: bus \(m \in M\) is dispatched from depot \(d \in D\), visits one disrupted station \(i \in I\) and transports passengers with the capacity constraint \(C\) to turnover station \(j \in J\), come back to visit another disrupted station \(k \in I\), then go to turnover station \(l \in J\) and so on, until the disruption returns to the normal operation. We refer the station that a bus is dispatched to its designated station \([1]\).

We assume that a significant number of passengers choose to transfer to buses; whenever a bus visits a station, it gets fully loaded immediately. The journey that a bus transports passengers from an affected station to the turnover station is defined as a bridging trip of the bus. Multiple bridging trips constitute the bridging route of a bus. Note that the origin of the first bridging trip should be its designated station. \(a_{mr}^{ij}\) denotes the departure time of bus \(m\) in the \(r\) th bridging trip at station \(i\). Different buses may visit a station one after the other, and let \(b_{fr}^i\) indicate the departure time of the \(f\) th pickup at station \(i\). Figure 2(a) illustrates an example of a metro disruption with four stations and three buses in one depot. The number in the circle indicates stations while the number beside the link denotes the travel time between stations or depots. Figure 2(b) presents bridging route for each bus. Solid arrows indicate bridging trips loading passengers and dotted arrows denote deadheading trips. For example, bus 3 visits station 3, station 2 and station 4 sequentially. For station 2, the first pickup is at the fourth minute by bus 1 and the second pickup is at 12th minute by bus 3, i.e. \(b_{11}^1 = 4, b_{21}^2 = 12\).

As for stations, we focus on waiting time. As shown in Figure 3, the horizontal axis represents time, and the yellow curve and blue curve indicate the cumulative arrival curve and the cumulative departure curve of passengers, respectively. Assuming that passenger arrival rate is constant during the time period \([0, H]\), \(p_0\) indicates the number of passengers who quit metro and transfer to buses at the beginning of the disruption. Note that we only consider passengers who arrive within \(H\) minutes after the disruption. \(p\) indicates the sum of arrival passengers during the disruption. The area bounded by the two curves and line \(y = T\) represents total waiting time for passengers. In Section 3, we propose a variant model of Vehicle Routing Problem (VRP). One of the objectives of the proposed model is to minimize total waiting time. The decision is to plan the bus bridging routing, that is, for each bus to determine which bridging trips are selected and what is the optimal order of them. We use binary variables \(x_{ij}^{mr}\) to determine whether bus \(m\) transport passengers from station \(i\) to station \(j\) in the \(r\) th bridging trip and binary variables \(y_{dij}^{fr}\) to determine whether bus \(m\) is dispatched from depot \(d\) to station \(i\).
To illustrate the modelling framework, consider a metro disruption case with five disrupted stations. Assuming that there are totally five buses available, a time space diagram of each bus is as shown in Figure 4(a). Buses are shown by different colors. We seek to obtain the optimal bus bridging solution with the minimum total waiting time. Figure 4(b) shows the optimized bus bridging solution. As can be seen, the optimized bus trajectories are more dispersed. Table 1 presents their specific bridging trips. The first number in the bracket represents the origin of the bridging trip and the second number represents the destination. For example, bus 1 as red trajectory shows is dispatched from depot 1 to station 2 and travels through bridging trips \((2,1) \rightarrow (3,1) \rightarrow (4,1)\) sequentially. Figure 5 shows the cumulative arrival curves and cumulative departure curves of station 3 in the upward direction before and after optimization. As can be seen from Figure 5(a), before optimization there are three pickups at station 3, at 22th, 47th, and 50th minute, respectively. The total waiting time is 896 hours, with 762 passengers stranded. In contrast, in the optimized solution, the total waiting time is 123 hours, with 288 passengers stranded. As shown in Figure 5(b), the cumulative departure curve is much closer to the cumulative arrival curve than the initial plan.

### 2.2. Assumptions

To simplify the problem, the following assumptions are made.

**Table 1: Optimization results in illustrative example.**

<table>
<thead>
<tr>
<th>Bus index</th>
<th>Depot</th>
<th>Bridging process</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>D1</td>
<td>((2,1) \rightarrow (3,5) \rightarrow (4,1))</td>
</tr>
<tr>
<td>2</td>
<td>D2</td>
<td>((3,5) \rightarrow (4,5) \rightarrow (5,1))</td>
</tr>
<tr>
<td>3</td>
<td>D2</td>
<td>((3,1) \rightarrow (2,5) \rightarrow (4,5) \rightarrow (5,1))</td>
</tr>
<tr>
<td>4</td>
<td>D3</td>
<td>((4,1) \rightarrow (2,1) \rightarrow (2,5) \rightarrow (5,1))</td>
</tr>
<tr>
<td>5</td>
<td>D3</td>
<td>((1,5) \rightarrow (5,1) \rightarrow (1,5))</td>
</tr>
</tbody>
</table>

**Figure 4:** Comparing time space diagram of buses before and after optimization in illustrative example.

**Figure 5:** Comparing cumulative arrival departure curves of station 3 in upward direction before and after optimization in illustrative example.
(1) Neglect delays due to traffic congestions and bus travel times between stations are constant.
(2) The parking capacity of the bus stop is large enough and the queue delay caused by the simultaneous arrival of multiple buses is not considered.
(3) Buses immediately depart from a station after arriving there and boarding time for passengers at each station is equal.

2.3. Notations. For illustrative purpose, main notations for the model are listed as follows.

Indices
- $i, k$: index of disruption stations, $i, k \in I$
- $j, l$: index of turnover stations, $j, l \in J$
- $g$: index of directions, $g = 1$ represents upward direction, $g = 2$ represents downward direction
- $m$: index of buses, $m \in M$
- $d$: index of bus depots, $d \in D$
- $r$: index of bridging trips for a bus, $r = 1, 2, \ldots R$
- $f$: index of pickups at a station, $f = 1, 2, \ldots F$

Sets
- $I$: set of disrupted stations, $I = I_1 \cup I_2$
- $J$: set of turnover stations, $J = J_1 \cup J_2$
- $N$: set of all stations, $N_g = I_g \cup J_g$, $N = N_1 \cup N_2$
- $D$: set of bus depots
- $M$: set of buses

Parameters
- $R$: the maximum allowed number of bridging trips for each bus
- $F$: the maximum allowed number of pickups at each station
- $V_d$: number of buses available in depot $d$
- $C$: bus capacity
- $T$: time when the metro line resumes operation
- $t_{ij}$: bus travel time between station $i$ and station $j$
- $t_{di}$: bus travel time between depot $d$ and station $i$
- $\lambda_i(t)$: number of arrival passengers at time $t$ at station $i$
- $q^{mr}_i$: departure time of the $r$ th bridging trip of bus $m$ at station $i$
- $b^t_i$: departure time of the $f$ th bus visiting station $i$
- $W^t_i$: number of passengers waiting for the $f$ th bus at station $i$
- $L^t_i$: number of passengers left behind by the $f$ th bus at station $i$

Decision Variables
- $x^{mr}_{ij}$: binary variable that equals 1 if the $r$ th bridging trip of bus $m$ traverses arc $(i, j)$, else 0
- $y^{mr}_{di}$: binary variable that equals 1 if bus $m$ is dispatched to station $i$ from bus depot $d$, else 0
- $p^{mr}_{ij}$: number of passengers form station $i$ to turnover station $j$ by the $r$ th bridging trip of bus $m$

3. Model Formulations

3.1. Objective Functions. A multi-objective model is formulated to design the bus dispatching plans. The first objective is to minimize total waiting time of passengers and the number of stranded passengers. Total waiting time represents interests of affected passengers, while from the respective of transit agencies, evacuating as many passengers as possible is critical. The number of stranded passengers is negative related to evacuation efficiency. The fewer stranded passengers indicate the higher evacuation efficiency. For model simplicity, set a virtual bus with zero capacity at time $T$, that is $b^T_i = T$. Thus the total stranded passengers at a station during the disruption can be represented by the number of passengers who failed to board on the $F$ th bus denoted by $L^F_i$. We convert $L^F_i$ into time by multiplying it by a penalty coefficient $P$. Added to the waiting time, the number of stranded passengers constitutes the first goal given by

$$
\min Z_1 = \sum_{i \in N} \sum_{f=2}^{F} \sum_{t \in [b^1_i, b^F_i]} \lambda_i(t) \left( b^t_i - t \right)^2 + \sum_{i \in N} \sum_{f=2}^{F} \sum_{t \in [b^1_i, b^F_i]} \left( b^t_i - b^{f-1}_i \right) + \sum_{i \in N} P \cdot L^F_i
$$

where the first two terms represent total waiting time. The first term calculates the waiting time of all passengers arriving during the disruption. The second term represents the waiting time of the passengers who were left behind by overcrowded buses. The number of passengers who failed to board the $f$ th bus at station $i$ is given by

$$
L^F_i = \max \left( W^F_i - C, 0 \right)
$$

where $C$ represents bus capacity, and $W^F_i$ denotes the number of passengers who wait for the $f$ th bus at station $i$. $W^F_i$ is given by

$$
W^F_i = L^F_{i-1} + \sum_{t \in [b^1_i, b^F_i]} \lambda_i(t)
$$

Equation (3) involves two terms: the first one is the number of passengers who were left behind by the $f$-1 th bus at station
i and the second one is the number passengers who arrived at station i during the period between the departure time of two consecutive buses, i.e. \( b_i^{f-1}, b_i^f \).

Buses are mainly sourced from spare vehicles at bus depots or borrowed from the regular bus operating lines. Bus resources are tight during peak hours. In order to minimize the impact on the operation of regular bus lines, the second objective is to minimize the number of buses dispatched in bus bridging services given by

\[
\min Z_2 = \sum_{m \in M} \sum_{d \in D} \sum_{i \in N} y_{di}^m
\]  

(4)

3.2. Model Formulations. The objective functions and constraints are as follows.

Equation (1)(4)

\[
\sum_{i \in N} \sum_{j \in J} x_{ij}^{mr} \leq 1 \quad \forall m \in M, \ r = 1, 2, \ldots, R
\]

(5)

\[
\sum_{i \in N} \sum_{j \in J} x_{ij}^{mr} = \sum_{i \in N} \sum_{j \in J} x_{ij}^{mr} = 0
\]

(6)

\[
\forall m \in M, \ r = 1, 2, \ldots, R
\]

\[
\sum_{j \in J} x_{ij}^{mr} = 0 \quad \forall m \in M, \ r = 1, 2, \ldots, R
\]

(7)

\[
\sum_{i \in N} \sum_{j \in J} x_{ij}^{mr} = \sum_{i \in N} x_{ij}^{mr} - x_{ij}^{mr} \geq 0
\]

(8)

\[
\forall j \in J, \ \forall m \in M, \ r = 2, \ldots, R
\]

\[
\sum_{d \in D} \sum_{i \in N} y_{di}^m \leq 1 \quad \forall m \in M
\]

(9)

\[
\sum_{d \in D} y_{di}^m = \sum_{j \in J} x_{ij}^{m1} \quad \forall i \in N, \ \forall m \in M
\]

(10)

\[
\sum_{m \in M} \sum_{d \in D} y_{di}^m \leq V_d \quad \forall d \in D
\]

(11)

\[
p_{ij}^{mr} \leq C x_{ij}^{mr}
\]

(12)

\[
\forall i, j \in N, \ \forall m \in M, \ r = 1, 2, \ldots, R
\]

\[
a_{ij}^{m1} = \sum_{d \in D} t_{di} y_{di}^m \quad \forall i \in N, \forall m \in M
\]

(13)

\[
a_{ij}^{mr} = a_{ij}^{m(r-1)} + \sum_{k \in N \setminus J} \sum_{j \in J} t_{kj} x_{kj}^{m(r-1)} + t_{ji} \left( \sum_{k \in N \setminus J} x_{kj}^{m(r-1)} + \sum_{k \in J} x_{kj}^{mr} - 1 \right)
\]

(14)

\[
\forall i \in N, \ m \in M, \ r = 2, \ldots, R
\]

\[
\sum_{i \in N} \sum_{j \in N} x_{ij}^{mr} \leq \sum_{i \in N} \sum_{j \in N} x_{ij}^{m(r-1)}
\]

(15)

\[
\forall m \in M, \ r = 2, \ldots, R
\]

\[
\sum_{i \in N} \sum_{j \in N} y_{ij}^m \leq \sum_{i \in N} \sum_{j \in N} x_{ij}^{m(r-1)}
\]

Equation (1) is the first objective to minimize total waiting time of passengers. Equation (4) is the second objective to minimize total dispatched buses. Constraints (5)-(8) ensure the construction of valid bridging routes. Constraint (5) makes sure that one bus can link two stations at most per bridging trip. Equation (6) prevents that a bus visits two stations in different directions per bridging trip. The subscripts indicate directions; 1 represents upward direction and 2 represents downward direction. For example, \( I_1 \) corresponds to set of bus stops in the downward direction. Equation (7) prevents that buses transport passengers from the turnover stations to disrupted stations. Constraint (8) indicates that the origin of the current bridging trip does not have to be the same as the destination of the previous one. Constraints (9)-(12) are dispatching constraints. Constraint (9) ensures that any bus can only be dispatched to one disrupted station. Constraint (10) indicates the relationship between variables \( x \) and \( y \), and ensures that for any bus the origin of the first bridging trip should be its designated station. Constraint (11) indicates that the number of buses dispatched from any depot cannot exceed its fleet size. Constraint (12) prevents any bus from picking up passengers from a station unless it is visiting the station and enforces the bus capacity restriction. Equations (13)-(14) depict the calculation of the departure time. Equation (13) indicates that for each bus the departure time of its first bridging trip is equal to the travel time between its depot and its designated station. Equation (14) indicates the relationship between the departure time of the \( r \)-th bridging trip and the \( r \)th bridging trip. The second term in the right side denotes the travel time of the \( r \)-1 bridging trip. The third term denotes the deadheading time between two consecutive bridging trips. Constraint (15) prevents that a bus restarts after it quits bus bridging. Constraints (16)-(17) ensure that the departure time of each bridging trip cannot exceed the recovery time. \( M \) is a large number. Equation (18) indicates the relationship of departure time for buses and pickup time for stations. For any station, the sequence of
pick-up time is the union of departure time of all buses visiting
the station. Constraint (19) specifies binary and nonnegative
integer types of decision variables.

4. Solution Algorithm

The proposed model is a multi-objective optimization model
based on VRP, belonging to NP-hard problems. It is difficult
to obtain an optimal solution in a limited time using an
accurate algorithm. A usual approach to solve multi-objective
optimization model is to assign a weight to each objective
and transform it into a single-objective one. However,
different transit agencies often have dissimilar preference
in the weights. Hence, it is more practicable to offer the
Pareto optimal solution to decision makers. Multi-objective
evolutionary algorithms optimize multiple objectives simul-
taneously and they are widely used in the transportation
field. Multi-objective evolutionary algorithms include several
such as Non-dominated Sorting Genetic Algorithm (NSGA-
II) [24], Multi-Objective Simulated Annealing (MOSA) [25],
Multi-Objective Tabu Search Algorithm (MOTS) [26], and
Multi-Objective Particle Swarm Optimization (MOPSO)
[27]. However, NSGA-II performs better in terms of finding
a diverse set of solutions and in converging to near the
true pareto-optimal set compared with others. In addition,
it is convenient to code for the bus bridging problem [28].
Therefore, we use NSGA-II algorithm to find the Pareto
solution set of the problem.

4.1. Code Representation. Filter out all the bridging trips
that satisfy the constraints (5)-(8) and number these bridging
paths sequentially. The chromosome length is set as the
product of fleet size and the maximum allowed number of
bridging trips of each bus \( R \). Each code on the chromosome
indicates the trip number chosen as the \( r \) th bridging trip of
the bus. Consider the illustrated example with five stations
in Section 2.1. For upward direction, there are four alternative
bridging trips: (1,5), (2,5), (3,5) and (4,5). Similarly there
are four counterparts in the downward direction. Number
these bridging trips from one to eight. Set \( R = 5 \), as Figure 6
shows, the codes numbered 3, 4, and 8 correspond to the
bridging trip (3,5), (4,5), and (5,1), respectively. The first
zero code and the subsequent code belonging to the bus are
treated as empty codes. The total number of bus bridging
trips is equal to the sum of occurrences of non-zero codes.
As shown in Figure 6, Bus 2 has a total of three bridging
trips.

4.2. Fitness Function. Use the objective function of the model
as the fitness function, i.e., the minimum total waiting time
and the minimum number of dispatched buses.

4.3. Algorithm Process. The concrete steps of the NSGA-II
algorithm in the proposed model are as follows:

Step 1 (initiation). Let \( s = 0 \) and set population size \( Z \),
Cross probability \( P_c \), Mutation probability \( P_m \). The maximum
iterations \( S \). Randomly generate initial parent population \( P_0 \).
Calculate the minimum total waiting time and the minimum
dispatching buses for all individuals in \( P_0 \).

Step 2. Perform a fast non-dominated sort operation for \( P_s \)
to obtain non-dominated sort values for all individuals.

Step 3. Perform selection, crossover, and mutation opera-
tions for \( P_s \) to generate individuals of descendant population
\( Q_s \). Calculate the minimum total waiting time and the
minimum using buses for all individuals in \( Q_s \).

Step 4. Combine \( P_s \) and \( Q_s \) into new populations called \( R_s \).
Perform a fast non-dominated sort operation for \( R_s \) to obtain
non-dominated sort values for all individuals.

Step 5 (elite retention operation). Let \( P_{s+1} = \phi \), \( r = 1 \). If the
sum of the number of individuals with the non-dominated
dominance sorting value of \( r \) is not larger than \( Z \), put all
individuals in \( R_s \) with the non-dominated dominance
sort value of \( r \) into \( P_{s+1} \). Otherwise, calculate the crowding
distance of individuals in the non-dominated sort value of
\( r \). Put those individuals into \( P_{s+1} \) in turn of their crowding
distances from the largest to the smallest, until the number of
individuals in the group \( P_{s+1} \) is \( Z \).

Step 6. Perform a fast non-dominated sort operation for \( P_{s+1} \)
to obtain non-dominated sort values for all individuals.

Step 7 (update Pareto solution set). All individuals with a
non-dominant ranking value of 1 in \( P_{s+1} \) constitute the Pareto
solution set.

Step 8 (determine whether the termination condition is met).
If \( s = S \), then output the Pareto solution set, terminate the
algorithm, otherwise, let \( s = s + 1 \) and go to Step 3.

5. Numerical Experiment

5.1. Case Setup. The numerical experiment considers an
emergent disruption on a metro line (No.9) in Shanghai,
China (see Figure 7). Five stations suspend services from 9:00
am to 10:30 am on a weekday. Station 1 and station 7 are
equipped with track crossovers. Totally 60 available spare
buses are reserved in seven nearby depots.

Table 2 presents passenger flow demands at each station.
Continuous passenger flows until 10:00 am are taken into
account. It should be noted that the destinations of bus
bridging are set as either upward or downward turnover
stations. Passenger demands between disrupted stations are
not considered. The assumption is reasonable since we conduct

\[ \text{Bus index} \]

\[ \begin{array}{c|c|c|c|c}
\text{Bus 1} & \text{Bus 2} & \text{Bus 3} & \ldots & \text{Bus n} \\
\hline
\end{array} \]

\[ \text{Bridging paths of bus 2} \]

\[ \begin{array}{cccccc}
3 & 4 & 8 & 0 & 0 & 0 \\
\end{array} \]

\[ \text{Figure 6: Coding of bridging path.} \]
### Table 2: Passengers flow demands in the numerical case.

<table>
<thead>
<tr>
<th>stations</th>
<th>Initial stranded passengers</th>
<th>Arrival rates (units: min⁻¹)</th>
<th>Initial stranded passengers</th>
<th>Arrival rates (units: min⁻¹)</th>
</tr>
</thead>
<tbody>
<tr>
<td>S1</td>
<td>300</td>
<td>19</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>S2</td>
<td>200</td>
<td>9</td>
<td>100</td>
<td>10</td>
</tr>
<tr>
<td>S3</td>
<td>140</td>
<td>8</td>
<td>150</td>
<td>11</td>
</tr>
<tr>
<td>S4</td>
<td>120</td>
<td>8</td>
<td>130</td>
<td>8</td>
</tr>
<tr>
<td>S5</td>
<td>170</td>
<td>7</td>
<td>210</td>
<td>7</td>
</tr>
<tr>
<td>S6</td>
<td>200</td>
<td>11</td>
<td>150</td>
<td>7</td>
</tr>
<tr>
<td>S7</td>
<td>—</td>
<td>—</td>
<td>400</td>
<td>19</td>
</tr>
</tbody>
</table>

### Table 3: Specific data of buses at depots and travel times between depots and stations.

<table>
<thead>
<tr>
<th>Depots</th>
<th>Number of buses</th>
<th>Stations (units: min)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>S1</td>
<td>S2</td>
</tr>
<tr>
<td>D1</td>
<td>8</td>
<td>11</td>
</tr>
<tr>
<td>D2</td>
<td>5</td>
<td>24</td>
</tr>
<tr>
<td>D3</td>
<td>7</td>
<td>25</td>
</tr>
<tr>
<td>D4</td>
<td>12</td>
<td>12</td>
</tr>
<tr>
<td>D5</td>
<td>9</td>
<td>16</td>
</tr>
<tr>
<td>D6</td>
<td>10</td>
<td>21</td>
</tr>
<tr>
<td>D7</td>
<td>9</td>
<td>19</td>
</tr>
</tbody>
</table>

### Table 4: Travel time between stations in the numerical case (units: min).

<table>
<thead>
<tr>
<th>Stations</th>
<th>S1</th>
<th>S2</th>
<th>S3</th>
<th>S4</th>
<th>S5</th>
<th>S6</th>
<th>S7</th>
</tr>
</thead>
<tbody>
<tr>
<td>S1</td>
<td>—</td>
<td>6</td>
<td>15</td>
<td>23</td>
<td>28</td>
<td>34</td>
<td>38</td>
</tr>
<tr>
<td>S2</td>
<td>6</td>
<td>—</td>
<td>14</td>
<td>21</td>
<td>22</td>
<td>27</td>
<td>25</td>
</tr>
<tr>
<td>S3</td>
<td>15</td>
<td>11</td>
<td>—</td>
<td>9</td>
<td>13</td>
<td>18</td>
<td>22</td>
</tr>
<tr>
<td>S4</td>
<td>25</td>
<td>19</td>
<td>8</td>
<td>—</td>
<td>9</td>
<td>10</td>
<td>14</td>
</tr>
<tr>
<td>S5</td>
<td>28</td>
<td>20</td>
<td>10</td>
<td>8</td>
<td>—</td>
<td>8</td>
<td>12</td>
</tr>
<tr>
<td>S6</td>
<td>36</td>
<td>25</td>
<td>15</td>
<td>9</td>
<td>7</td>
<td>—</td>
<td>7</td>
</tr>
<tr>
<td>S7</td>
<td>38</td>
<td>29</td>
<td>19</td>
<td>11</td>
<td>9</td>
<td>7</td>
<td>—</td>
</tr>
</tbody>
</table>

---

**Figure 7:** Disrupted scenario of the metro line in the case study.

- **Turnover station**
- **Disrupted rail links**
- **Disrupted station**
- **Operating rail links**

The case study on a commuting line where the majority of passengers transfer at turnover stations. Table 3 shows the amount of available buses at depots and their travel time between depots and stations. Table 4 presents bus travel time between stations. Bus capacity is 80 passengers. Parameters $R$ and $F$ in the proposed model are set as 10 and 90, respectively. Parameters $Z$, $S$, $p_c$, $p_m$ in solution algorithm are set at 500, 100, 0.9, and 0.4, respectively. All experiments are performed on a PC with a 1.6 GHz Intel Core i5 CPU and 8.0 GB memory, and the algorithm is in MATLAB. The computational time for solving the proposed model is 7min.

### 5.2 Optimal Results and Analysis.

A total of 14 Pareto approximate optimal solutions were obtained. Figure 8 presents the first objective function value and the number of dispatched buses. For example, in Pareto solution 1, all 60 buses are dispatched to handle the emergency. Table 5 presents the total waiting time and the number of stranded passengers. The standard bus bridging plan described in Section 1 is set for comparison. In the standard plan, the buses are dispatched firstly to the nearby turnover stations and then shuttle on
Table 5: Comparing of optimization results before and after optimization.

<table>
<thead>
<tr>
<th></th>
<th>Waiting time (units: hours)</th>
<th>Number of stranded passengers</th>
<th>Number of dispatched buses</th>
</tr>
</thead>
<tbody>
<tr>
<td>Optimal plan</td>
<td>1504</td>
<td>0</td>
<td>60</td>
</tr>
<tr>
<td>Standard plan</td>
<td>4724</td>
<td>1737</td>
<td>60</td>
</tr>
</tbody>
</table>

routes parallel with the disrupted lines. Buses load passengers at each station along the routes.

As can be seen, in the standard plan, there are as many as 1737 passengers stranded at the end of disruption and the total waiting time is much higher than that of the optimal plan.

Figure 9 shows the distribution of total waiting time of each station in the upward direction before and after optimization. As can be seen, the total waiting time in the optimal plan was much lower than that in the standard plan especially for downstream stations. For example, at station 5 and station 6, the total waiting time in the optimal plan decreased by 82.6% and 82.1%, respectively. In the standard plan a large number of passengers got aboard at upstream stations, while the passengers at downstream stations were unlikely to get on buses due to insufficient remaining capacity. Similarly, in Figure 10, for example, passengers at station 4 were unlikely to get on board until 40 minutes after the disruption and no passengers at station 6 could board buses. Our bridging strategy achieved a more equitable condition for passengers among all stations.

Figure 10 describes the accumulative arrival/boarding curves of passengers at each station in the upward direction to station 7. All passengers board buses when the boarding curve (red/blue) coincides with arrival curve (grey). As can be seen from Figure 10, the red curve is closer to the grey one than blue one which indicates the total waiting time (enclosed area by the two curves) in the proposed plan is shorter than that in the standard plan. Total evacuation time when all passengers get on board of the optimal plan is shorter than that of the standard plan for any station except for station 2. For example, in station 5, total evacuation time was 72 minutes in the optimal plan, while it was not until the last minute before the operation resumed that all passengers got on board in the standard plan.

In the standard plan, multiple vehicles were firstly dispatched to turnover stations, so passengers could be quickly evacuated in the early stage of the evacuation. However, due to quite long turnaround time of buses, few buses arrived again in a long time period (about 40 minutes) until buses in the opposite direction turned over, resulting in long waiting time. In contrast, in the optimal plan, with the node-to-node transportation mode buses were able to operate flexibly on different bridging trips. These trips covered different travel times, resulting buses turned over at various time. Therefore, the pickup time distribution of stations was more discrete, which more accorded with linear passenger arrival.

5.3. Sensitivity Analysis. We further investigate the impact of recovery time and fleet size on results. Figure 11 presents total waiting time and the number of the stranded passengers given varying bus fleet sizes under certain recovery time. As the fleet size grew, the total waiting time and the number of stranded passengers both lessoned with smaller decreasing amplitude which is in accordance with empirical laws. For example, when recovery time was 90 minutes (see Figure 11(c)), total waiting time decreased by 30.9% when the fleet size increased from 30 to 40, while it only decreased by 2.5% when the fleet size increased from 80 to 90. According to Figure 11(c), if the goal is achieving a complete evacuation before the recovery time, the lowest level of availability for buses should be maintained at 60.

Figure 12 shows the variation of average waiting time with variable fleet size under different recovery time. The horizontal axis indicates fleet size. Four curves represent the variation of average waiting time under different recovery
Figure 10: Cumulative arrival and boarding process at stations in upward direction.
time. It is shown that as bus fleet size increases the average waiting time decreases, but the decrease rate is smaller with larger fleet size. The variation range gets larger under the longer recovery time. For example, when the recovery time is 30 minutes the average waiting time decreases from 11 mins to 5 mins, while the recovery time is 2 hours it decreases from 34 mins to 5 mins. Here are some suggestions for transit agencies: when the recovery time is quite long, e.g., 90 minutes, as many buses as possible should be dispatched, while when it is shorter, e.g., 30 minutes, the fleet size can be smaller.

Since passenger demand increases until time $H$ (one hour in the case study) and we only consider the period before recovery time $T$, when $T$ is above or below $H$, the passenger demand is different. For example, we can calculate from Table 2: when $T$ is 90 minutes the passenger demand is 9710 while when $T$ is 30 minutes the demand is 5990. It can be seen that under different recovery time the absolute quantity of stranded passengers cannot reflect the actual evacuation efficiency. Therefore, we look into an indicator: evacuation efficiency factor. Define evacuation efficiency factor as the proportion of the number of passengers who have got on board before recovery time to total passenger flows. Evacuation efficiency factor is given by:

\[
\eta = \left(1 - \frac{\Phi}{\Lambda}\right) \times 100\% \tag{20}
\]

In Equation (20), $\eta$ indicates evacuation efficiency factor, $\Phi = \sum_{i \in N} L^F_i$ represents the total number of stranded passengers, and $\Lambda = \sum_{i \in N} \sum_{t \in [0,T]} \bar{\lambda}_i(t)$ denotes the total passenger flows.

Figure 13 presents the variation of evacuation efficiency factor with different fleet size. The horizontal axis represents fleet size and four curves represent evacuation efficiency factors under different recovery time. As the fleet size increased,
Figure 12: Performance of waiting time with varying bus fleet size.

Figure 13: Performance of evacuation efficiency factor with varying bus fleet size.

\( \eta \) also increased. For example, when recovery time was 90 minutes and the fleet size was 30 vehicles, the corresponding \( \eta \) was only 77.4%, while when the fleet size increased to 60, \( \eta \) reached 100%. In addition, when the recovery time was less than 60 minutes, the evacuation efficiency factor was significantly lower than that of longer recovery time at the same fleet size. For example, when the recovery time was 30 minutes, even if the fleet size increased to 100, it was still unlikely to evacuate all passengers. When there were 60 available buses, \( \eta \) reached nearly 90%. It was unnecessary for transit agencies to dispatch numerous buses.

6. Conclusions

Considering dynamic passenger flow demands and recovery time, we formulated an optimization-based model to provide bus bridging services with two objective functions, one that minimized total waiting time and the number of stranded passengers and one that minimized the total number of dispatched buses. We solved the multi-objective model with NSGA-II algorithm. A case study on Shanghai Metro Line 9 was conducted finally. We can draw some conclusions as follows.

(a) A comparison with the standard bridging plan was carried out. The results show that the proposed model can significantly reduce total waiting time and the number of stranded passengers especially at the downstream stations. Due to the node-to-node bus bridging mode, more passengers can board buses than those in the standard bridging plan. The proposed model achieves a more equitable distribution of passenger condition between stations.

(b) Sensitivity analysis of fleet size under different recovery time was conducted. It was found that as the fleet size increased, average waiting time and the number of stranded passengers showed a decreasing trend, but when the fleet size reached a certain value, for instance, 70 in the case study, the decline was not significant. Similar analysis will facilitate transit agencies to determine suitable fleet size.

(c) Evacuation efficiency under different recovery time was evaluated. It was found that as the fleet size increased, the evacuation efficiency also increased. When recovery time was smaller than 60 minutes, even if the fleet size increased to 100, it was still difficult to evacuate all passengers. The travel time of buses dispatched to disrupted stations accounted for a large proportion of the disruption duration, leading to limited evacuation efficiency. In conclusion, transit agencies should dispatch as many buses as possible under recovery time more than 60 minutes, and vice versa.

Future research should focus on developing a more efficient algorithm applicable to real-world situations. Besides, since it is difficult to predict dynamic passenger flows under unexpected disruption, we assumed that passenger arrival follows a uniform distribution. We would like to develop a bus bridging optimization model under real-time dynamic passenger flows in the future.

Data Availability

All data generated and analyzed during the current study are available from the corresponding author on reasonable request.

Conflicts of Interest

The authors declare that they have no conflicts of interest.

Acknowledgments

This work is financially supported by the National Key Basic Research Program of China (2012CB725403).
References


