A Cost Effective Route Guidance Method for Urban Areas Using Histograms

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Analysing shortest path for real time traffic environment is crucial with dynamic updating. VANET technology can be used for analysing traffic but generates a huge amount of data to be exchanged, which demands more processing power and resources. In this paper, a new histogram-based route guidance algorithm (HBA) has been proposed based on light weight processing. The proposed algorithm enables selecting the shortest path between any source and destination using the histogram models, which capture the higher order distribution function of the number vehicles in every lane. Furthermore, the histogram model is used to estimate the traffic delays at intersections and roundabouts. The data entity collection through sensors used for histogram modelling is presented in detail. The experimental results show that the proposed algorithm provides a good prediction of road traffic status and a better solution for the congestion problem in the urban areas.

1. Introduction

Due to the growth of urban areas [1–3], the vehicle density per area square mile has been increased tremendously and thus impacts the roadway trajectory directly to have more congestion and heavy traffic jams. This challenge cannot be solved only by expanding the existing road traffic infrastructure by itself and it may be infeasible with restricted roads. To remedy congestion and traffic jams, it is important to take precaution measures and the existing roadways should be utilized properly. At the moment, Intelligent Transportation System [4–7] (ITS) is becoming an increasingly important element in the road traffic system. Intelligent transport system includes wireless automated technologies to integrate with the vehicles, drivers, and other infrastructures to communicate with each other to have better driving experience. This system helps to optimize transportation travel time [8–10] and enables finding desired routes for the vehicles. For example, when implementing ITS in roadways it alerts the drivers from a potential traffic jam and provides the best shortest path [11–13] to the destination. When comparing with the static approach the best shortest path likely to be evaluated from the dynamic approach which analyses the real time traffic situations [14]. Informing traffic patterns [15, 16] in prior makes the trip more robust. Precautious measure congestion based cost function should be estimated for every lane to predict the best possible route from source to destination [17, 18].

In general, the cost function can be derived from the past road traffic traces. Due to the large participant, the data generated from vehicles and infrastructure will be substantially large, thus making the system unstable. However, it requires heavy investment in terms of computation and power factors.

In this paper, we consider the abovementioned problem and propose a road traffic model which analyses the vehicles travelling in roads and queuing up in different road cross junctions. The proposed model captures the arrival rate distribution in VANET and build histograms on a time window basis to derive the shortest path for every source to destination trip. The paper is organised as follows: Section 2 discuss the existing works of route guidance in ITS environment. Section 3 describes the proposed algorithm and the
queuing modelling. Section 4 presents the network modelling and the route guidance strategy is discussed in Section 5. In Section 6, the performance analysis is discussed with various scenarios. Finally, Section 7 concludes the paper.

2. Related Work

Congestion in roadways has been identified as a major cause for the problems created in urban areas, whereby creating distress to travellers. Many researchers have been prudent to analyse the congestion problem in urban areas. Real time traffic data gives the insight of the real time events and can be used in classifying congestion [19] for route forecasting. Such technology would also lead to improvements in existing transportation while reducing the cost of new infrastructure. Barrachina et al. [20] presented a paper for continuously estimating the vehicular density using V2I for traffic redirection. It allows the intelligent transport system to estimate the density by computing the number of beacons received per roadside unit. Pan et al. [21] proposed proactive five different rerouting algorithms that include the dynamic shortest path (DSP), the A+ shortest path with repulsion (AR+), the random k shortest path (RkSP), the entropy-balanced kSP (EBkSP), and the flow-balanced kSP (FBkSP) and studied the effective travel time. However, the algorithms were active all the time and thus the vehicles trend to take some of the shortest route more often and leaving behind other routes. This could again make the network unstable due to congestion in some particular roads. Rajabi-Bahaabadi et al. [22] propose an efficient path finding algorithm that is suitable for road networks using nondominant sorting algorithm and its parameters are enhanced by the Taguchi method. Nevertheless, one limitation to the V2I architecture is that it contains a single point-of-processing, meaning that the central server process a large amount of data in real time that must be difficult by the server when all of the vehicles are sending speed and location data, as well as querying for fastest paths (or other information). The proposed traffic modelling architecture attempts to reduce these limitations by offering the benefits of the histogram graphs to the vehicles. In the proposed algorithm, the transportation network is broken into individual lanes, each of which consists of one or more edges. The lanes estimate the congestion identically by counting the vehicles and then exchange the information with other RSU units. The model uses less data exchange rate as it does not indulge vehicles in route estimation process. The proposed model is explained in detail in further chapters.

3. Proposed Algorithm

The proposed histogram-based algorithm (HBA) is divided into five stages: (1) Traffic Monitoring, (2) Histogram Computing, (3) Queuing Model (4), Network Modelling (V2I), and (5) Route Guidance. The hierarchy of each stage in the proposed algorithm is shown in Figure 1. The description of each stage is presented in the following sections.

3.1. Traffic Monitoring. Modelling and analysis of the road traffic in a statistical way give the insight of understanding the relationship between the higher order distribution function and the roadway navigation algorithms. Every lane has its own characteristics based on the vehicle injection behaviour over the time period, for example, some intersections will have more congestion during the office hours and some intersections will be more congested during the school timings. The congestion happening in urban areas is heterogeneous and changes over the period of time. In the proposed algorithm, this behaviour of changing patterns has been studied and taken into consideration. The traffic flow can be predicted by observing the vehicles travelling in every lane. To estimate the traffic flow of individual lanes, a density function should be derived as follows. ITS uses wireless sensor nodes that is intended for monitoring the road environment. The major advantage of using wireless sensor nodes in the road side will enable to capture the real time behaviour of roadways. The sensors will sense and collect data from a certain lane, process them, and transmit them to the RSU (Road Side Unit). Sensors are used in every lane to count the number of vehicles on a window basis. This information is used to build histograms by the RSU and used for further processing. As shown in Figure 2, a signalized lane is connected by two intersections in each end is selected and taken as an example. Each lane has its own sensors coupled with the corresponding RSU. It is done by providing unique identification code (e.g., SeqNum) for both sensors and RSU units. When the information is broadcasted, only the corresponding RSU will process the data based on the unique code. In the diagram, we clearly see that when the sensors transmit the data, it is received only by the corresponding RSU unit.
The principle behind the estimation of density is to find the cost function for all individual lanes. Imagine there can be different routes exist between a source and destination. To select a best shortest path, we need to estimate the cost of every route based on the congestion by using the real time data and cost can be a factor that estimates the total travelling time in that particular lane.

3.2. Histogram Computation. In this paper, we use histograms to model traffic distribution in urban areas. We sense the road traffic flows on every lane for a predefined period (also known as an observation period) and build histograms on window basis. A new histogram per lane is built in every observation period \( O \), where the length of the observation period depends on the lane characteristics. During every observation period, the number of vehicles in each lane is counted over \( N \) equal sampling periods of length \( T \) each, where \( N = O/T \). The frequency of these numbers is then computed to construct the histogram model.

As the number of vehicles in any sampling period can vary over a large set of integers \( I = [0, 1, \ldots, K] \), where \( K \) is the maximum number of vehicles that can be held in the street, we divide this \([0 \cdots K]\) range into a limited number of equal width classes \( N_c \) and compute the grouped probability distribution function over each class. Table 1, for example, shows the distribution of the number of vehicles in one lane over \( N_c = 6 \) classes during an observation period \( O = 10 \) minutes and sampling period \( T = 15 \) seconds. The number of sampling periods \( N = 10 \times 60/15 = 40 \) sampling period. For each class \( i \), the table shows the lower limit \( C_{i-} \), the upper limit \( C_{i}^{+} \), the middle point \( C_i \), the number of intervals during which this class range was encountered, and the probability

<table>
<thead>
<tr>
<th>Class number</th>
<th>Class range ([C_{i-}, C_{i}^{+}])</th>
<th>Class midpoint (C_i)</th>
<th>Frequency of occurrence</th>
<th>(P(C_i))</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>([0, 10])</td>
<td>5</td>
<td>8</td>
<td>0.2</td>
</tr>
<tr>
<td>1</td>
<td>([10, 20])</td>
<td>15</td>
<td>16</td>
<td>0.4</td>
</tr>
<tr>
<td>2</td>
<td>([20, 30])</td>
<td>25</td>
<td>6</td>
<td>0.15</td>
</tr>
<tr>
<td>3</td>
<td>([30, 40])</td>
<td>35</td>
<td>6</td>
<td>0.15</td>
</tr>
<tr>
<td>4</td>
<td>([40, 50])</td>
<td>45</td>
<td>4</td>
<td>0.1</td>
</tr>
<tr>
<td>5</td>
<td>([50, 60])</td>
<td>55</td>
<td>0</td>
<td>0</td>
</tr>
</tbody>
</table>
of occurrence of this class \( P(C_i) \) during the observation period. The probability \( P(C_i) \) is calculated as shown in

\[
P(C_i) = P \left( \frac{\text{No of sampling periods with vehicle counts in the range } (C^{-i} \leq C_i \leq C^+)}{\text{Total number of sampling periods}} \right),
\]

(1)

The histogram is then derived from the table to measure the congestion density of a particular lane during an observation window as shown in Figure 3. To ensure that the statistical model of histogram is suitable for practical use, we assume that the number of vehicles within an observation period/window is a stationary process and independent from window to window.

In general, any statistical histogram \( C \) can be defined using two main attributes:

(i) A set of quantized number of cars per class \( C \), which can be derived by averaging the upper and lower limits of number of cars in each class.

(ii) A set of probabilities \( P \) corresponding to \( C \):

\[
C = (C, P) = \begin{cases} 
C = C_i, & i = 0, 1, \ldots, N_C - 1, \\
P = P(C_i), & i = 0, 1, \ldots, N_C - 1,
\end{cases}
\]

(2)

where \( N_C \) is the number of classes.

The histogram presented in this paper represents the statistical distribution of vehicles in an observation window basis against their corresponding probabilities (or frequencies). The number of vehicles can be easily obtained using sensors along lanes, and the histograms can be computed using the Road Side Units (RSUs) and then disseminated to vehicles moving on the road through an ad hoc based infrastructure-to-vehicle network.

3.3. Queuing Models for Vehicular Road Traffic. Urban roadways commonly have three types of road junctions [25], including road junction with traffic-signals, road junction with round-about, and road junction with fly-over. Initially, we consider a road junction with traffic signal interconnected by two-way roads. Let \( r(t) \) \( t \in \{1, 2, \ldots, N\} \) be the vehicle speed, in kilometre per hour, during a sampling period \( t \). According to the 2-second rule for road safety, the minimum distance between consecutive vehicles on the street is given by:

\[ d(t) = \frac{r(t) \cdot 1000 \cdot 2}{3600} \text{ meters} \]

If \( \bar{V} \) is the average length of vehicle in meters, then vehicles from the street will move into the road cross junction at a rate of \( D(t) = \frac{r(t)(d(t) + \bar{V})}{3600} \text{ vehicles per second} \). Furthermore, if \( W \) is the length of the road between two adjacent road junctions, then the maximum number of vehicles that can be held by the road during a sampling period \( t \) is given by \( K(t) = \frac{W}{d(t) + \bar{V}} \text{ vehicles} \).

The road traffic signal is modelled as a Markov modulated Bernoulli process (MMPB) consisting of ON state and OFF state. The ON state represents the green signal and the OFF state represents the red signal (in this case, yellow signal is not taken into account). If vehicles from different directions are equally accessing to the road cross junction, which is controlled by traffic lights with equal access time operating on opposite two-way traffic basis, then vehicles from one direction moves into the inter-section at an average rate of \( \frac{D(t)}{2} \text{ vehicle per second} \). Hence, the queue at the cross junctions can be modelled as finite buffer HD/D/1/K queue.

Where, HD stands for the histogram Deterministic Inter-arrival distribution that represents vehicle flow in the queue, \( D \) represents the deterministic process of vehicles moving from the street into the cross junction at a rate of \( D(t)/2 \) vehicle per second. The server discipline is first come first served.

The usage of the HD/D/1/K queuing model can be justified as follows. From statistical analysis point of view, the performance of the queue depends on the interarrival of
vehicles moving on the street, which is relatively determined by the number of vehicles moving into the system within a single unit of time. Although the number of vehicles moving into the queuing system is time variable, it is much slower comparing to the time for a vehicle across the queue. In this case, if the histogram window is chosen properly, the vehicle interarrival process within the window can be treated as a stable process. We also assume that vehicles are independent of each other and move from road into the cross junction at a constant speed since it is determined by the road traffic signal.

Likewise, a street connected with a round-about junction is modelled as finite buffer HD/D/1/K queue, where D represents the deterministic process of vehicles moving from street into the round-about junction at an average rate of $D(t)/4$. In other words, we assume that vehicles randomly access the round-about junction from 4 directions with equal distribution with a mean value of $D(t)/4$. Similarly, a road junction with fly-over can be modelled as free-access buffer without traffic control mechanism. Therefore, a roadway in urban city can be modelled as a tandem chain of HD/D/1/K queues [26].

Let $V(t) | t \in \{1, 2, \ldots, N\}$ represent the number of vehicles that enter the street during the $t$th sampling period. Let $\tau(t)$ be the number of vehicles that exit from the street during the same period. The vehicle input flow is supplied through a street/buffer of finite capacity. The pending vehicles that cannot move from the street and enter the road intersection during the observation period will be queued in the street. The system is assumed to be in stationary if the pending traffic convergences to a finite value. Let $Q[t]$ represents the queue length in sampling period $t \in \{1, 2, \ldots, N\}$. In the classic analysis of M/D/1/K queuing system, the $Q[t]$ can be expressed as follows:

$$Q[t] = \theta^K_0 (Q[t-1] + V(t) - \tau[t]), \quad (3)$$

where Operator $\theta$ limits the queue lengths so that they cannot either underflow or overflow the lane capacity $K(t)$ in the $t$th sampling period. The $\theta$ operator is defined as shown in

$$\theta^K_0 (x) = \begin{cases} 0, & \text{for } x < a, \\ x - a, & \text{for } a \leq x < b + a, \\ b, & \text{for } x \geq b + a. \end{cases} \quad (4)$$

Since the number of vehicles moving from the street into the cross junction during the $t$th sampling window is given by $D(t) \times T$ vehicles, then the number of vehicles that exit from the street during the same period is given by $\tau[t] = D(t)/2 \times T$, where $T$ is the duration of the sampling period. Likewise, if the cross junction is a round-about, then $\tau[t] = D(t)/4 \times T$. When $V[t] \leq \tau[t]$, the buffer size is not increased. Otherwise, the buffer occupancy increases until queuing buffer is full. Therefore, equation (5) can be expressed as

$$Q[t] = \theta^K_0 (Q[t-1] + V(t) - \tau) = \theta^K_0 (Q[t-1] + V(t)). \quad (5)$$

The bound operator $\theta^K_0$ is the statistical generalization of the previously defined $\theta^K_0$ operator. $C$ is the statistical histogram defined in (2). The operator $\otimes$ stands for the standard statistical convolution, which is described as follows. If $X$ and $Y$ are two independent random variables with $n$ and $m$ intervals, respectively, the convolution $X \otimes Y$ is a new random variable $Z$ with $n + m - 1$ intervals and

$$P_Z[i] = \sum_{k=0}^{i} P_X(i-k) \times P_Y(k). \quad (7)$$

The steady state probability of the queue length can be calculated using Algorithm 1, where $C$: arrival process (histogram) as defined in (2), $\tau$: service rate class of $\tau$, $\tau = \tau/L_c$, where $L_c$ is the class length, $K$: the maximum queue length class of $K$, $K = \lfloor K/L_c \rfloor$, where $L_c$ is the class length, and $E(X)$: the mean value (or expectation) of gpd $X$.

Once the steady state probabilities of the queue length are calculated, the expected queue length can be calculated as follows:

$$E(Q) = \sum_{i=0}^{K} i^* P(q_i). \quad (8)$$

In the case of a cross junction with an output rate of $\tau$ and a queue length characterized by a gpd $Q$, the queuing delay $W$ is proportional to $Q$ and has the same pmf, where $W = 1/\tau \times Q$. Therefore, the expected queuing delay at each intersection can be calculated as follows:

$$E[W] = \frac{1}{\tau} E[Q]. \quad (9)$$

**4. Network Modelling (I2V)**

The network modelling for the proposed algorithm uses the infrastructure-to-vehicle (I2V) communications for transferring the histograms across the network. Once an RSU
computes the histogram for a corresponding lane, it informs the other RSU units to build up the histogram map of the entire network. Finally, this information is given to the vehicles travelling in the network and a sequence is added along with it. The vehicles cache the information and verify the sequence number; if the information is updated then the new information is processed and the map is updated with new histogram. Figure 4(a) shows an experimental result of the statistical histograms built in the city of Abu Dhabi obtained by on-side road traffic sensor system.

5. Route Guidance

As discussed in the above scheme, the histograms were computed in every lane. From this information, we can find the shortest path from source to destination across the network as shown in Figure 4(b). The proposed model adapts Dijkstra’s algorithm to find the shortest path. In the proposed model, the road networks are considered as a weighted graph model, the lanes are perceived as edges, and the intersections are marked as nodes. The weight of each link, in other words the link cost, is computed as follows:

\[
\text{LinkCost} = \frac{\text{LinkLength}}{\text{Speed}} + \text{Queuing delay},
\]

where the Queuing delay is derived from (8). The histograms are computed across every lane and the link density function is updated periodically. The vehicles are routed according to the lane weights.

6. Performance Metrics

The performance analysis of the HBA-based algorithm is compared with two different existing algorithms namely, the intervehicular communication (IVC) [28] algorithm and the optimal dynamic shortest path algorithm (SPA). The working model of IVC algorithm introduces the concept of searcher vehicle and candidate vehicle, whose origin node matches the destination node of the searcher vehicle which travels in the opposite direction. SPA algorithm computes the shortest path by exchanging traffic information between these vehicles using V2V communication.

In the SPA-Dynamic algorithm, the cost of a link is computed by monitoring the traces of the vehicles across the entire network. Hence, the link cost gets updated with every single vehicle passing through the link. All links’ costs are disseminated in real time to all vehicles in the network so that they can compute the shortest path to destinations. The SPA-Dynamic approach is infeasible in practice as it demands dissemination of continuous real time data in VANET environment. Moreover, this algorithm also consumes more time and requires high resource utilization. Therefore, the SPA-Dynamic algorithm is considered to be an optimal algorithm in simulation (but not in practice). However, we are interested in adding the SPA-Dynamic approach in our performance analysis to prove the efficiency of the proposed algorithm and to depict how close the proposed algorithm can perform with the theoretical model that holds the best optimal results.

A microscopic simulator I-Sim [29] is employed to perform the comprehensive comparative analysis for the proposed system. This simulator effectively creates urban roadway traffic scenarios and generates vehicle trips randomly. The microscopic visualization of the simulator enables in analysing the traffic behaviour of roadways on a broader percept. While running the simulation, every vehicle stores all individual road trip information from source to destination, which enables determining the efficiency of the proposed algorithm. This simulator is inbuilt with a car following technique that adapts a finite state model which determines the acceleration/deceleration rate, velocity, and the distance between the vehicles. The simulator is efficient enough to designing the wide range of roadways that reflect the urban roadway system.

To study the effectiveness of the proposed system, a grid network shown in Figure 5 is used to model the urban roadway. In the topology, the grid has 49 nodes and each node acts as an intersection. The edge between any two nodes acts as a lane and its length varies according to the scenario. In each scenario, the simulator generates the vehicles randomly and the vehicle burst rate is set pertaining to the given scenario. While running the simulation, the vehicle routes are computed based on the link cost and the routes are updated in the OD Matrix. If the congestion increases in the route lane,
a new route is computed based on the histogram model and the vehicles are rerouted with the updated shortest path.

As shown in Table 2, different scenarios were created to study the effectiveness of the proposed algorithm by varying the source volume injection rate, lane length, and maximum speed limit.

The total number of origins in the grid network is set as 14 sources, in which the vehicles start their trip randomly. In order to simulate the real time urban traffic, only 28 intersections are marked as destinations. The vehicles start the trip from any one of the sources and reach the assigned destination by taking the shortest route. Every link connected between intersections will have 3 lanes and same speed limit. Every scenario is simulated for 60 minutes and the average travel time is computed.

6.1. Source Volume Scenario. In this scenario, the link length and the lane speed are kept as constant and the volume demand differs with different vehicle burst rate from the source. Three scenarios were carried out using the link speed of 80 km/hr; horizontal and vertical lane lengths are 300 x 150. Figure 6 shows average travel time versus the source volume generation.

In the experiment, it shows that the average travel time taken by every algorithm increases as the source volume increases. This is due to the congestion happening in the entire traffic network area. The HBA proposed algorithm gives better results than the IVC algorithm and gives close results with the SPA, which is considered to be an optimal algorithm. However, the SPA approach is infeasible in practice as dissimilating real time data continuously and instantaneously in VANET environment is impossible due to time and computation constrains.

6.2. Link Speed Scenario. In general, as the link speed increases, the network accommodates more vehicles. Furthermore, the relative network productivity increases with the increase of speed. Therefore, as the link speed increases, the network accommodates more vehicles. As shown in Figure 7, the performance of HBA gives good results when the lane speed increases, the vehicles traveling through the shortest path link are purely driven by histograms that are available in every lane as a cost function. Finally, an increase in speed significantly increases the efficiency of the proposed algorithm.

6.3. Link Length Scenario. Figure 8 shows the average travel time performance measure with various network sizes. For all the three scenarios, the source volume burst rate is fixed
Table 2: Experimental scenario details.

<table>
<thead>
<tr>
<th>Source volume scenario (veh/hr)</th>
<th>Link speed scenario (km/hr)</th>
<th>Link length scenario (meters)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Horizontal link length = 300 m</td>
<td>Link length = 150 m</td>
<td>Source volume = 2000 veh/hr</td>
</tr>
<tr>
<td>Vertical link length = 150 m</td>
<td></td>
<td>Horizontal link length = 300 m</td>
</tr>
<tr>
<td>Vertical link length = 150 m</td>
<td></td>
<td>Link speed = 80 km/hr</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Vehicle burst rate</th>
<th>Link maximum speed</th>
<th>Link distance</th>
</tr>
</thead>
<tbody>
<tr>
<td>Experiment 1</td>
<td>500</td>
<td>60</td>
</tr>
<tr>
<td>Experiment 2</td>
<td>1000</td>
<td>80</td>
</tr>
<tr>
<td>Experiment 3</td>
<td>2000</td>
<td>100</td>
</tr>
</tbody>
</table>

Figure 6: Source volume scenario.

Figure 7: Link speed scenario.

Figure 8: Link length scenario.

The speed is fixed. The HBA travels from source to destination with less travel time comparatively with the IVC algorithm.

By analysing the figure, it shows that when link length increases and the speed is kept as constant the relative average travel time increases. It shows that the minimum average travel time is associated with less grid length in (300 × 150) scenario. This is mainly because the path is progressively determined with the ratio of the histogram graph.

7. Conclusion

This paper presented a novel route guidance algorithm, which forecasts traffic delays using lightweight histogram models. The histogram model extended the deterministic model, which usually describe the traffic with a discrete number of classes with a quantifying probability. Adapting the queuing model in the system makes the model more invigorate and helps to estimate the real time behaviours of roadways. A microscopic simulation is used to assess the effectiveness of the algorithm comparatively against two algorithms from literature. The performance was analysed with various scenarios reflecting the network source volumes, link speeds, and lengths. The proposed algorithm performed as 2000 veh/hr and the link speed as 80 km/hr. The average travel time for HBA ranges within 11.23–16.02 minutes as the travel time is directly proportional to the lane length when
comparably with the optimal dynamic shortest path algorithm and outperformed the IVC algorithm, which results in lesser travel time from source to destination.

The results also showed that the proposed road traffic model can be easily obtained and transmitted in real time by individual vehicles while they are travelling on streets or queuing in road cross junctions and provides a good prediction on road traffic status.

The proposed model can be enhanced further by combining different histogram statistics with respect to speed, utilization, and throughput. This will be investigated in further research.

Competing Interests

The authors declare that they have no competing interests.

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