# Research Article

# **Improved Bounds for Radio** k-Chromatic Number of Hypercube $Q_n$

# Laxman Saha, Pratima Panigrahi, and Pawan Kumar

Department of Mathematics, IIT Kharagpur, Kharagpur 721302, India

Correspondence should be addressed to Laxman Saha, laxman.iitkgp@gmail.com

Received 13 December 2010; Accepted 18 February 2011

Academic Editor: Brigitte Forster-Heinlein

Copyright © 2011 Laxman Saha et al. This is an open access article distributed under the Creative Commons Attribution License, which permits unrestricted use, distribution, and reproduction in any medium, provided the original work is properly cited.

A number of graph coloring problems have their roots in a communication problem known as the channel assignment problem. The channel assignment problem is the problem of assigning channels (nonnegative integers) to the stations in an optimal way such that interference is avoided as reported by Hale (2005). Radio *k*-coloring of a graph is a special type of channel assignment problem. Kchikech et al. (2005) have given a lower and an upper bound for radio *k*-chromatic number of hypercube  $Q_n$ , and an improvement of their lower bound was obtained by Kola and Panigrahi (2010). In this paper, we further improve Kola et al.'s lower bound as well as Kchikeck et al.'s upper bound. Also, our bounds agree for nearly antipodal number of  $Q_n$  when  $n \equiv 2 \pmod{4}$ .

### **1. Introduction**

Radio coloring is derived from the assignment of radio frequencies (channels) to a set of transmitters. The frequencies assigned depend on the geographical distance between the transmitters: the closer two transmitters are, the greater the potential for interference between their signals. Thus, when the distance between two transmitters is small, the difference in the frequencies assigned must be relatively large, whereas two transmitters at a large distance may be assigned relatively close frequencies.

Radio *k*-coloring of a graph is a variation of the channel assignment problem. For a simple connected graph *G* of order *n* and diameter *q*, and a positive integer *k* with  $1 \le k \le q$ , a *radio k-coloring f* of *G* is an assignment of non-negative integers to the vertices of *G* such that for every two distinct vertices *u* and *v* of *G* 

$$|f(u) - f(v)| \ge k + 1 - d(u, v),$$
 (1.1)

where d(u, v) is the distance between u and v in G. The *span* of a radio k-coloring f,  $span_k(f)$ , is the maximum integer assigned to a vertex of G. The *radio* k-chromatic number of G, denoted by  $rc_k(G)$ , is defined by  $rc_k(G) = min\{span_k(f) : f \text{ is a radio } k\text{-coloring of } G\}$ . Since  $rc_1(G)$  is the chromatic number  $\chi(G)$ , radio k-colorings are a generalization of ordinary vertex coloring of graphs. In the literature,  $rc_q(G)$  is termed as *radio* number of G, whereas  $rc_{q-1}(G)$  and  $rc_{q-2}(G)$  are called *antipodal* and *nearly antipodal* number of G, respectively. Since the problem is to find a radio k-coloring with minimum span, we use the least color 0 in every radio k-coloring appearing in this paper.

Finding radio k-chromatic number of a graph is an interesting yet difficult combinatorial problem with potential applications to FM channel assignment. The concept of radio k-coloring was introduced by Chartrand et al. in [1]. But so far, radio k-chromatic number is known for very limited classes of graphs and for specific values of k only. Radio numbers of  $C_n$  and  $P_n$  [2],  $C_n^2$  [3], and  $P_n^2$  [4] are known. Recently, radio numbers of complete *m*-array trees have been found by Li et al. [5]. Radio number of hypercubes is determined by Khennoufa and Tongi [6] and Kola and Panigrahi [7] independently. Khennoufa et al. have also found antipodal number of  $Q_n$ . For k = n - 1, n - 2, n - 3 and n - 4 the radio k-chromatic number of path  $P_n$  has been determined in [4, 8–10], respectively. Juan and liu [11] have determined the antipodal number of cycle  $C_n$ . Kchikech et al. [12] have given an upper and a lower bound for radio k-chromatic number of cartesian product of two graphs G and G'. Also, as a particular case, they have given an upper and a lower bounds for radio *k*-chromatic number of hypercube  $Q_n$ , but these bounds were very loose. Recently, Kola and Panigrahi [7] have improved their lower bound for  $rc_k(Q_n)$  and have shown that this bound is exact for radio number of  $Q_n$ . In this paper, we obtain an improvement of Kola et al.'s lower bound as well as Kchikeck et al.'s upper bound. Further, we show that these bounds agree for  $rc_{n-2}(Q_n)$ ,  $n \equiv 2 \pmod{4}$ .

#### 2. Improved Lower Bound

In this section, we give an improved lower bound of  $rc_k(Q_n)$  for  $2 \le k \le n-2$ . For a hypercube  $Q_n$  of dimension n, the vertex set can be taken as binary n-bit strings and two vertices being adjacent if the corresponding strings differ at exactly one bit.

*Definition 2.1.* For any two *n*-bit binary strings  $a = a_0a_1 \cdots a_{n-1}$  and  $b = b_0b_1 \cdots b_{n-1}$ , the Hamming distance  $d_H(a, b)$  between *a* and *b* is the number of bits in which they differ. In particular, if  $x, y \in \{0, 1\}$ , then  $d_H(x, y) = 0$  or 1 according as x = y or  $x \neq y$ .

If *u* and *v* are two vertices of  $Q_n$  with *a* and *b* as the corresponding strings, then  $d_{Q_n}(u, v) = d_H(a, b)$ . Two *n*-bit binary strings may differ in at most *n* positions, so diameter of  $Q_n$  is *n*.

The results in the following lemma may be found in [7].

**Lemma 2.2.** For any three vertices x, y, and z of the hypercube  $Q_n$ , the following holds:

- (i)  $d(x, y) + d(y, z) + d(x, z) \le 2n$ ,
- (ii) d(x, y) + d(y, z) + d(x, z) = 2n, if d(x, y) = n.

The following lower bound for  $rc_k(Q_n)$  was determined by Kola and Panigrahi [7].

**Theorem 2.3** (see [7]). For the hypercube  $Q_n$  of dimension  $n \ge 2$  and for any positive integer k with  $2 \le k \le n$ ,

$$\operatorname{rc}_{k}(Q_{n}) \geq \begin{cases} \left(\frac{3(k+1)-2n}{2}\right)2^{n-1} - \frac{k+1}{2}, & \text{if } k \text{ is odd,} \\ \left(\frac{3(k+1)-2n+1}{2}\right)2^{n-1} - \frac{k+2}{2}, & \text{if } k \text{ is even.} \end{cases}$$

$$(2.1)$$

In the theorem below, we give an improvement of the above lower bound of  $rc_k(Q_n)$  for all values of n and  $2 \le k \le n-2$ .

**Theorem 2.4.** For a hypercube  $Q_n$  of dimension  $n \ge 2$  and for any positive integer k with  $2 \le k \le n-2$ ,

$$rc_k(Q_n) \ge \left[\frac{3(k+1)-2n}{2}\right] \left(2^{n-1}-1\right).$$
 (2.2)

*Proof.* Let *f* be any radio *k*-coloring of  $Q_n$  and  $x_1, x_2, ..., x_{2^n}$  be an ordering of the vertices of  $Q_n$  such that  $f(x_{j+1}) \ge f(x_j), 1 \le j \le 2^n - 1$ . Obviously,  $f(x_1) = 0$  and span $(f) = f(x_{2^n})$ . Since *f* is a radio *k*-coloring of  $Q_n$ , for every *i* with  $0 \le i \le 2^n - 2$ , we have the following:

$$f(x_{i+1}) - f(x_i) \ge k + 1 - d(x_{i+1}, x_i),$$
  

$$f(x_{i+2}) - f(x_{i+1}) \ge k + 1 - d(x_{i+2}, x_{i+1}),$$
  

$$f(x_{i+2}) - f(x_i) \ge k + 1 - d(x_{i+2}, x_i).$$
  
(2.3)

Adding (2.3) we get,

$$2(f(x_{i+2}) - f(x_i)) \ge 3(k+1) - (d(x_i, x_{i+1}) + d(x_{i+2}, x_{i+1}) + d(x_{i+2}, x_i))$$
  
$$\ge 3(k+1) - 2n \quad \text{(from (i) of Lemma 2.2)}.$$
(2.4)

From the above inequality, we have  $f(x_{i+2}) - f(x_i) \ge \lfloor (3(k+1) - 2n)/2 \rfloor$ , for all  $i = 1, 3, ..., 2^n - 3$  and summing up these we get,  $f(x_{2^n-1}) \ge \lfloor (3(k+1) - 2n)/2 \rfloor (2^{n-1} - 1)$ . From this inequality and the fact that  $f(x_{2^n}) \ge f(x_{2^{n-1}})$ , we get the result.

#### 3. Improved Upper Bound

In this section, we give an improved upper bound of  $rc_k(Q_n)$ . For better presentation, we need the definition below.

*Definition* 3.1. For two positive integers *n* and *l* with l < n, a binary (n, l)-Gray code is a listing of all the *n*-bit binary strings such that the Hamming distance between two successive strings is exactly *l*. Further, a quasi (n, l)-Gray code is a listing of all the *n*-bit binary strings such that the Hamming distance between two successive strings is exactly *l* except between the two successive strings is exactly *l* and  $2^{n-1} - 1$  and  $2^{n-1}$  for which it is l - 1 or l + 1.

*Notation* 3.2. For any positive integer *n*, we define  $\delta(n)$  as

$$\delta(n) = \begin{cases} 0, & \text{if } n \equiv 2 \pmod{4}, \\ 1, & \text{if } n \not\equiv 2 \pmod{4}. \end{cases}$$
(3.1)

The following theorem gives a partition of the vertex set of  $Q_n$  that will be used to find the upper bound of  $rc_k(Q_n)$ .

**Theorem 3.3.** For a hypercube  $Q_n$  of dimension n, there exists a partition of the vertex set of  $Q_n$  with the partite sets  $U_1 = \{x_i : i = 0, 1, ..., 2^{n-1} - 1\}$  and  $U_2 = \{y_i : i = 0, 1, ..., 2^{n-1} - 1\}$  satisfying the followyng properties:

where  $\delta(n)$  is defined as in Notation 3.2.

*Proof.* Let  $Q'_{n-1}$  and  $Q''_{n-1}$  be two copies of  $Q_{n-1}$  induced by all the *n*-bit strings with starting bit 0 and 1, respectively. Let  $x_0, x_1, \ldots, x_{2^{n-1}-1}$  be the ordering of the vertices of  $Q'_{n-1}$  which induce a quasi  $(n-1, \lfloor n/2 \rfloor)$ -Gray code if  $n \not\equiv 2 \pmod{4}$  and a (n-1, n/2)-Gray code if  $n \equiv 2 \pmod{4}$ ; (such code exists, see [6]). We take  $y_i = \overline{x}_i$ ,  $i = 1, 2, \ldots, 2^{n-1} - 1$ , that is,  $\overline{x}_i$  is obtained from  $x_i$  by changing 0s to 1s and vice versa. Now,  $y_0, y_1, \ldots, y_{2^{n-1}-1}$  is an ordering of the vertices of  $Q'_{n-1}$  which induces the same Gray code as in the above. By taking  $U_1 = \{x_i : i = 0, 1, \ldots, 2^{n-1} - 1\}$  and  $U_2 = \{y_i : i = 0, 1, \ldots, 2^{n-1} - 1\}$ , we get the result.

The following upper bound for  $rc_k(Q_n)$  was determined by Kchikech et al. in [12].

**Theorem 3.4** (see [12]). For the hypercube  $Q_n$  of dimension  $n \ge 2$  and for any  $k \ge 2$ ,

$$rc_k(Q_n) \leq (2^n - 1)k - 2^{n-1} + 1.$$
 (3.2)

The theorem below gives an improvement of this upper bound of  $rc_k(Q_n)$ .

**Theorem 3.5.** For the hypercube  $Q_n$  of dimension  $n \ge 2$  and for any positive integer  $k, 2 \le k \le n-2$ ,

$$rc_{k}(Q_{n}) \leq \begin{cases} (2^{n-1}-1)\left(k+1-\left\lfloor\frac{n}{2}\right\rfloor\right)+\delta(n), & \text{if } 2\left\lfloor\frac{n}{2}\right\rfloor-2 \leq k \leq n-2, \\ k2^{n-2}-\left\lceil\frac{k}{2}\right\rceil+\delta(n), & \text{if } \left\lfloor\frac{n}{2}\right\rfloor-1 \leq k \leq 2\left\lfloor\frac{n}{2}\right\rfloor-3, \\ k2^{n-2}-k, & \text{if } 2 \leq k \leq \left\lfloor\frac{n}{2}\right\rfloor-2, \end{cases}$$
(3.3)

where  $\delta(n)$  is defined as in Notation 3.2.

$u, v \in U_1 \cup U_2$	Value of $ i - j $	d(u,v)	k+1-d(u,v)	f(u) - f(v)
$x_i, x_j \in U_1$ , or $y_i, y_j \in U_2$	i - j  = 1	[ <i>n</i> /2]	$k + 1 - \lfloor n/2 \rfloor$	$k + 1 - \lfloor n/2 \rfloor$
	$ i-j  \ge 2$	$\geqslant 1$	$\leqslant k$	$\geqslant k$
$x_i \in U_1, y_j \in U_2$	i - j = 0	п	$\leqslant 0$	0
	i - j  = 1	$\geq \lfloor n/2 \rfloor$	$\leq k + 1 - \lfloor n/2 \rfloor$	$\geq k + 1 - \lfloor n/2 \rfloor$
	$ i-j  \ge 2$	$\geqslant 1$	$\leqslant k$	$\geqslant k$

Table 1

*Proof.* Here, we consider the same partition of  $V(Q_n)$  as in Theorem 3.3. We first take  $2\lfloor n/2 \rfloor - 2 \le k \le n-2$ . For these values of k, we define a coloring f of  $V(Q_n)$  with  $f(x_0) = 0$  and

$$f(x_i) = f(x_{i-1}) + k + 1 - \left\lfloor \frac{n}{2} \right\rfloor, \quad i = 1, 2, \dots, 2^{n-1} - 1, \text{ except } i = 2^{n-2},$$

$$f(x_{2^{n-2}}) = f(x_{2^{n-2}-1}) + k + 1 - \left\lfloor \frac{n}{2} \right\rfloor + \delta(n),$$

$$f(y_i) = f(x_i), \quad i = 0, 1, 2, \dots, 2^{n-1} - 1.$$
(3.4)

We first check the radio condition for the vertex  $x_{2^{n-2}}$  with all other vertices. Let  $v \in V(Q_n) - \{x_{2^{n-2}}\}$ . From the definition of f,  $|f(x_{2^{n-2}}) - f(u)| \ge k$ , except  $v \in \{x_{2^{n-2}-1}, y_{2^{n-2}}, x_{2^{n-2}+1}, y_{2^{n-2}+1}\}$ . From Theorem 3.3(b) and Lemma 2.2(ii), we have  $d(x_{2^{n-2}-1}, x_{2^{n-2}}), d(y_{2^{n-2}-1}, x_{2^{n-2}}) \ge [n/2] - \delta(n)$  and therefore  $f(x_{2^{n-2}}) - f(v) = k + 1 + \delta(n) - [n/2] \ge k + 1 - d(x_{2^{n-2}}, v)$ , for  $v = x_{2^{n-2}-1}$  or  $y_{2^{n-2}-1}$ . Similarly, for  $v = x_{2^{n-2}+1}$  or  $y_{2^{n-2}-1}$ , we show  $f(x_{2^{n-2}}) - f(u) \ge k + 1 - d(x_{2^{n-2}}, u)$ . Since  $f(x_{2^{n-2}}) = f(y_{2^{n-2}})$  and  $d(x_{2^{n-2}}, y_{2^{n-2}}) = n$ , radio condition is trivially true for the pair  $(x_{2^{n-2}}, y_{2^{n-2}})$ . In a similar manner one can show that  $|f(y_{2^{n-2}}) - f(v)| \ge k + 1 - d(y_{2^{n-2}}, v)$ , for all  $v \in V(Q_n) - \{y_{2^{n-2}}\}$ . For the rest of the pairs of vertices in  $V(Q_n)$ , checking radio condition is straightforward from the definition of f. However, in Table 1, we compute the values of |f(u) - f(v)| and k + 1 - d(u, v) for every pair of vertices  $u, v \in V(Q_n) - \{x_{2^{n-2}}, y_{2^{n-2}}\}$ .

Therefore, *f* is a radio *k*-coloring of  $Q_n$  and the span of *f* is  $(2^{n-1}-1)(k+1-\lfloor n/2 \rfloor)+\delta(n)$ . Next, we consider the value of *k* as  $\lfloor n/2 \rfloor \leq k \leq 2\lfloor n/2 \rfloor - 3$  and define a coloring *f* with  $f(x_0) = 0$ , and

$$f(x_{i+1}) = f(x_i) + \left\lfloor \frac{k}{2} \right\rfloor, \text{ if } i \text{ is even and } i \neq 2^{n-2} - 1,$$

$$f(x_{i+1}) = f(x_i) + \left\lfloor \frac{k}{2} \right\rfloor, \text{ if } i \text{ is odd},$$

$$f(x_{i+1}) = f(x_i) + \left\lfloor \frac{k}{2} \right\rfloor + \delta(n), \quad i = 2^{n-2} - 1,$$

$$f(y_i) = f(x_i), \quad i = 0, 1, 2, \dots, 2^{n-1} - 1.$$
(3.5)

By a similar way as in the above, one can check that f is a radio k-coloring of  $Q_n$  with the span  $f(x_{2^{n-1}-1}) = 2^{n-2}\lfloor k/2 \rfloor + (2^{n-2}-1)\lfloor k/2 \rfloor + \delta(n) = k2^{n-2} - \lfloor k/2 \rfloor + \delta(n)$ .

Finally, we take  $2 \le k \le \lfloor n/2 \rfloor - 2$ . From Theorem 2.4, we have  $d(x_i, x_{i+1})$ ,  $d(x_i, y_i)$ ,  $d(x_i, y_{i+1})$ ,  $d(x_{i+1}, y_i)$ ,  $d(y_i, y_{i+1}) \ge k + 1$ . So, we use the same color for the vertices  $x_i, x_{i+1}, y_i$ 

and  $y_{i+1}$ . In this case, we define a coloring f as  $f(x_i) = (i/2)k = f(y_i)$ ,  $f(x_{i+1}) = f(x_i)$ ,  $f(y_{i+1}) = f(y_i)$ ,  $i = 0, 2, ..., 2^{n-1} - 2$ . It is easy to show that f is a radio k-coloring of  $Q_n$  with the span  $k2^{n-2} - k$ .

Observe that the bounds given in Theorems 2.4 and 3.5 agree for k = n - 2 with  $n \equiv 2 \pmod{4}$ . Therefore we have determined the nearly antipodal number of  $Q_n$  which is given in the theorem below.

**Theorem 3.6.** For  $n \equiv 2 \pmod{4}$ , the nearly antipodal number  $rc_{n-2}(Q_n)$  of the hypercube  $Q_n$  is  $(2^{n-1}-1)(n-2)/2$ .

# 4. Concluding Remarks

Improved lower and upper bounds of  $rc_k(Q_n)$  have been obtained in Theorems 2.4 and 3.5, respectively. It is easy to verify that the bound given in Theorem 3.5 is an improvement over the bound in Theorem 3.4. However, below, we check that lower bound given in Theorem 2.4 is really an improvement over the bound in Theorem 2.3. From Theorem 2.4, we have

$$\operatorname{rc}_{k}(Q_{n}) \geq \begin{cases} \left(\frac{3(k+1)-2n}{2}\right)(2^{n-1}-1), & \text{if } k \text{ is odd,} \\ \left(\frac{3(k+1)-2n+1}{2}\right)(2^{n-1}-1), & \text{if } k \text{ is even,} \end{cases}$$

$$= \begin{cases} \left(\frac{3(k+1)-2n}{2}\right)2^{n-1}+n-\frac{3(k+1)}{2}, & \text{if } k \text{ is odd,} \\ \left(\frac{3(k+1)-2n+1}{2}\right)2^{n-1}+n-\frac{3k+4}{2}, & \text{if } k \text{ is even.} \end{cases}$$

$$(4.1)$$

For odd integer k,

$$\left[ \left( \frac{3(k+1)-2n}{2} \right) 2^{n-1} + n - \frac{3(k+1)}{2} \right] - \left[ \left( \frac{3(k+1)-2n}{2} \right) 2^{n-1} - \frac{k+1}{2} \right]$$
  
=  $n - k - 1 \ge 1$ , if  $2 \le k \le n - 2$ . (4.2)

Again, if *k* is an even integer, then

$$\left[\left(\frac{3(k+1)-2n+1}{2}\right)2^{n-1}+n-\frac{3k+4}{2}\right] - \left[\left(\frac{3(k+1)-2n+1}{2}\right)2^{n-1}-\frac{k+2}{2}\right]$$

$$= n-k-1 \ge 1, \quad \text{if} \ 2 \le k \le n-2.$$
(4.3)

Theorem 3.6 gives the nearly antipodal number of  $Q_n$  for  $n \equiv 2 \pmod{4}$ . From Theorems 2.4 and 3.5, one gets that the difference between the lower and upper bound of  $\operatorname{rc}_{n-2}(Q_n)$ ,  $n \equiv 0 \pmod{4}$ , is equal to one, that is,  $(2^{n-1} - 1)(n-2)/2 \leq \operatorname{rc}_{n-2}(Q_n) \leq (2^{n-1} - 1)(n-2)/2 + 1$ .

#### References

- [1] G. Chartrand, D. Erwin, and P. Zhang, "A graph labeling problem suggested by FM channel restrictions," *Bulletin of the Institute of Combinatorics and Its Applications*, vol. 43, pp. 43–57, 2005.
- [2] D. D.-F. Liu and M. Xie, "Radio number for square paths," *Ars Combinatoria*, vol. 90, pp. 307–319, 2009.
  [3] D. D.-F. Liu and M. Xie, "Radio number for square of cycles," *Congressus Numerantium*, vol. 169, pp. 105–125, 2004.
- [4] D. Liu and X. Zhu, "Multilevel distance labelings for paths and cycles," *SIAM Journal on Discrete Mathematics*, vol. 19, no. 3, pp. 610–621, 2005.
- [5] X. Li, V. Mak, and S. Zhou, "Optimal radio labellings of complete *m*-ary trees," *Discrete Applied Mathematics*, vol. 158, no. 5, pp. 507–515, 2010.
- [6] R. Khennoufa and O. Togni, "The radio antipodal and radio numbers of the hypercube," Ars Combinatoria. In press.
- [7] S. R. Kola and P. Panigrahi, "An improved lower bound for the radio *k*-chromatic number of the hypercube  $Q_{n'}$ " *Computers & Mathematics with Applications*, vol. 60, no. 7, pp. 2131–2140, 2010.
- [8] R. Khennoufa and O. Togni, "A note on radio antipodal colourings of paths," *Mathematica Bohemica*, vol. 130, no. 3, pp. 277–282, 2005.
- [9] S. R. Kola and P. Panigrahi, "Nearly antipodal chromatic number *ac*'(*P<sub>n</sub>*) of the path *P<sub>n</sub>*," *Mathematica Bohemica*, vol. 134, no. 1, pp. 77–86, 2009.
- [10] S. R. Kola and P. Panigraĥi, "On radio (n 4)-chromatic number of the path  $P_n$ ," AKCE International Journal of Graphs and Combinatorics, vol. 6, no. 1, pp. 209–217, 2009.
- [11] J. S. Juan and D. D.-F. Liu, Antipodal labeling of cycles, Manuscript, 2006, http://www .calstatela.edu/faculty/dliu/dliu.htm.
- [12] M. Kchikech, R. Khennoufa, and O. Togni, "Radio k-labelings for cartesian products of graphs," *Electronic Notes in Discrete Mathematics*, vol. 22, pp. 347–352, 2005.



Advances in **Operations Research** 



**The Scientific** World Journal







Hindawi

Submit your manuscripts at http://www.hindawi.com



Algebra



Journal of Probability and Statistics



International Journal of Differential Equations





Complex Analysis





Mathematical Problems in Engineering



Abstract and Applied Analysis



Discrete Dynamics in Nature and Society



International Journal of Mathematics and Mathematical Sciences





Journal of **Function Spaces** 



International Journal of Stochastic Analysis

