Research Article

Distance Measurements Related to Cartesian Product of Cycles

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Graph theory and its wide applications in natural sciences and social sciences open a new era of research. Making the graph of computer networks and analyzing it with aid of graph theory are extensively studied and researched in the literature. An important discussion is based on distance between two nodes in a network which may include closeness of objects, centrality of objects, average path length between objects, and vertex eccentricity. For example, (1) disease transmission networks: closeness and centrality of objects are used to measure vulnerability to particular disease and its infectivity; (2) routing networks: eccentricity of objects is used to find vertices which form the periphery objects of the network. In this manuscript, we have discussed distance measurements including center, periphery, and average eccentricity for the Cartesian product of two cycles. The results are obtained using the definitions of eccentricity, radius, and diameter of a graph, and all possible cases (for different parity of length of cycles) have been proved.

1. Introduction

Applications of graph theory to computer science, physics, chemistry, biology, social sciences, and statistics open up a new dimension for researchers [1–5]. One of the attributes is distance and its related measurements in the graph. Weighted distance, topological distance, eccentricity, radius, diameter, metric dimension, indices, etc., are such distance-related terms and have received much attention of researchers [6–8]. One of the fundamental questions related to distance measurement is community detection and location of their emergency facilitation within the network [9, 10]. The study of networks such as (1) social networks like Facebook, Twitter, LinkedIn, etc., and (2) biological networks like protein–protein interaction, gene transcription, ecological networks, etc., and statistical inference on these network models have been done extensively in [11, 12]. In distance-based networks, several vertices can have different closeness as well as exactly the same closeness with respect to a particular facility like hospital, electricity, etc. [13–17]. In this paper, we consider the distance measure, vertex eccentricity, and its associated definitions center and periphery. The indices related to vertex eccentricity are discussed in [18, 19]. For an undirected graph, Goddard et al. in [20, 21] have shown the following:

\[ \text{rad}(M) \leq \text{diam}(M) \leq 2\text{rad}(M). \]  

(1)

Other proved results are

(1) \( \text{rad}(K_m) = \text{diam}(K_m) = 1 \) for \( m \geq 2 \)
(2) \( \text{rad}(C_m) = \text{diam}(C_m) = \lfloor m/2 \rfloor \)
(3) \( \text{rad}(K_{p,m}) = \text{diam}(K_{p,m}) = 2, p, m \geq 2 \)
(4) \( \text{rad}(P_m) = \lfloor (m - 1)/2 \rfloor, \text{diam}(P_m) = m - 1 \)

Diameter of a tree, random graphs, and bridge graphs are determined in [22, 23], respectively.
The eccentricity denoted by $ec(x)$ of a node $x$ in a connected graph $M$ is defined as

$$ec(x) = \text{Max}\{d(y,x): y \in V(M)\}. \quad (2)$$

The radius denoted by $rad(M)$ of a connected graph $M$ is defined as

$$rad(M) = \text{Min}\{ec(x): x \in V(M)\}. \quad (3)$$

The diameter denoted by $diam(M)$ of a graph is defined as

$$diam(M) = \text{Max}\{ec(x): x \in V(M)\}. \quad (4)$$

In [24], authors introduced average eccentricity denoted by $Avg_{ec}$ for a graph $M$ with $m$ number of vertices as

$$Avg_{ec}(M) = \frac{1}{m} \sum_{x \in V(M)} ec_{_M}(x). \quad (5)$$

Buckley [25] defined the eccentric set of a nontrivial connected graph and proved the criteria for a nonempty set of positive integers to be an eccentric set of some graphs. Buckley also defined central subgraphs embedding and proved that the central subgraph of a tree is isomorphic to $K_1$ or $K_2$. Dankelmann and Osaye [26] proved results on average eccentricity for different parameters related to eccentricity: weight function, total weight function, eccentric sequence of tree for given diameter, and $k$-star. Dankelmann et al. in [27] prove the bounds for eccentricity and average eccentricity of the graph, subgraph, and its complement when the graph is replaced by a spanning tree or spanning graph. Yu et al. [28, 29] characterize the extremal unicyclic graphs among other $m$-unicyclic graphs with minimal and second minimal average eccentricity. Ilic [30, 31] discusses the graph transformations which change the eccentricity of a graph. He also solved four conjectures about average eccentricity, clique number, domination, and independent number using the system AutoGraphix.

Theorem 1. The family of the Cartesian product $C_p □ C_q$ is self-centered.

Proof. We will prove the result for some choices of $p$ and $q$ as given in the following cases:

Case 1. When $p \equiv 0 \pmod{2}, q \equiv 0 \pmod{2}, p \geq q$.

Consider the cycle $(x_{(1,1)}x_{(1,2)}x_{(1,3)} \cdots x_{(1,r)} \cdots x_{(1,p)})$ and choose an arbitrary vertex $x_{(1,1)}$ on this cycle:

$$V(C_p □ C_q) = \{x_{(r,s)}: 1 \leq r \leq p, 1 \leq s \leq q\},$$

$$E(C_p □ C_q) = \{x_{(r,s)}x_{(r+1,s)}: 1 \leq r \leq p - 1, 1 \leq s \leq q\} \cup \{x_{(r,s)}x_{(r,s+1)}: 1 \leq r \leq p, 1 \leq s \leq q - 1\} \quad (6).$$

Theorem 2. The family of the Cartesian product $C_p □ C_q$ is self-centered.

Proof. We will prove the result for some choices of $p$ and $q$ as given in the following cases:

Case 1. When $p \equiv 0 \pmod{2}, q \equiv 0 \pmod{2}, p \geq q$.

Consider the cycle $(x_{(1,1)}x_{(1,2)}x_{(1,3)} \cdots x_{(1,r)} \cdots x_{(1,p)})$ and choose an arbitrary vertex $x_{(1,1)}$ on this cycle:

$$d(x_{(1,1)}x_{(1,1)}) = r - 1, \quad 1 \leq r \leq \frac{p}{2} + 1,$$

$$d(x_{(1,1)}x_{(1,2)}) = \frac{p}{2} - 1, \quad r = \frac{p}{2} + 2,$$

$$d(x_{(1,1)}x_{(1,r)}) = 1 (7)$$
In order to locate a vertex at an extreme distance from \( x_{(1,1)} \) in \( C_p \square C_q \), we have to consider values \( 1 \leq r \leq (p/2) + 1 \). Since each \( x_{(1,r)} \) is the neighbor of \( x_{(2,r)} \) and \( x_{(q,r)} \), therefore, using equations (7) and (8), we have

\[
d(x_{(1,1)}, x_{(2,r)}) = d(x_{(1,1)}, x_{(q,r)}) = r, \quad 1 \leq r \leq \frac{p}{2} + 1,
\]

(9)

\[
d(x_{(1,1)}, x_{(q,r)}) = d(x_{(1,1)}, x_{(2,r)}) = p - r, \quad 2 + \frac{p}{2} \leq r \leq p.
\]

(10)

Further, \( x_{(2,r)} \) is the neighbor of \( x_{(3,r)} \) and \( x_{(q,r)} \) is the neighbor of \( x_{(q-1,r)} \), and therefore equations (9) and (10) give

\[
d(x_{(1,1)}, x_{(3,r)}) = d(x_{(1,1)}, x_{(q-1,r)}) = r + 1, \quad 1 \leq r \leq \frac{p}{2} + 1,
\]

\[
d(x_{(1,1)}, x_{(q-1,r)}) = d(x_{(1,1)}, x_{(3,r)}) = p + 3 - r, \quad \frac{p}{2} + 2 \leq r \leq p.
\]

(11)

Moreover, \( x_{(3,r)} \) is the neighbor of \( x_{(4,r)} \) and \( x_{(q-1,r)} \) is the neighbor of \( x_{(q-2,r)} \), therefore,

\[
d(x_{(1,1)}, x_{(4,r)}) = d(x_{(1,1)}, x_{(q-2,r)}) = r + 1, \quad 1 \leq r \leq \frac{p}{2} + 1,
\]

\[
d(x_{(1,1)}, x_{(q-2,r)}) = d(x_{(1,1)}, x_{(4,r)}) = p + 4 - r, \quad \frac{p}{2} + 2 \leq r \leq p.
\]

(12)

Continuing the same procedure, after \((q/2)\) steps, the vertex \( x_{((q/2),r)} \) would be in the neighborhood of \( x_{((q/2)+1,r)} \), which implies

\[
d(x_{(1,1)}, x_{((q/2),r)}) = r + \frac{q}{2} - 1, \quad 1 \leq r \leq \frac{p}{2} + 1,
\]

\[
d(x_{(1,1)}, x_{((q/2)+1,r)}) = p - r + \frac{q}{2} + 1, \quad \frac{p}{2} + 2 \leq r \leq p.
\]

(13)

This means \( x_{((q/2)+1,(p/2)+1)} \) is farthest from \( x_{(1,1)} \). Therefore, \( ec(x_{(1,1)}) = (p/2) + (q/2) \).

Similarly, \( x_{((q/2)+2,(p/2)+1)} \), \( x_{((q/2)+3,(p/2)+1)} \), and \( x_{((q/2)+4,(p/2)+1)} \) are farthest from \( x_{(2,1)}, x_{(3,1)}, x_{(4,1)}, \ldots, x_{(q,1)} \) in \( C_p \square C_q \) respectively.

Since the graph is symmetric, each vertex on either cycle has the same eccentricity.

\[
rad(C_p \square C_q) = \text{diam}(C_p \square C_q) = \frac{p}{2} + \frac{q}{2}
\]

(14)

Consequently, each vertex is a central vertex as well as a peripheral vertex.

Case 2. When \( p \equiv 1, q \equiv 1 \) (mod2), \( p \geq q \).

Consider the cycle \( x_{(1,1)}x_{(1,2)}x_{(1,3)}\ldots x_{(1,p)} \) and select a vertex \( x_{(1,1)} \) from it, and then

\[
d(x_{(1,1)}, x_{(1,r)}) = r - 1, \quad 1 \leq r \leq \frac{p - 1}{2} + 1.
\]

(15)

When values of \( r \) vary between \((p - 1)/2 + 2\) and \( p \), the distance \( d(x_{(1,1)}, x_{(1,r)}) \) varies between \((p - 1)/2\) and 1:

\[
d(x_{(1,1)}, x_{(1,r)}) = p + 1 - r, \quad \frac{p - 1}{2} + 2 \leq r \leq p.
\]

(16)

Thus, to locate a farthest vertex from \( x_{(1,1)} \) in \( C_p \square C_q \), we only consider \( 1 \leq r \leq (p - 1)/2 + 1 \).

Each \( x_{(1,r)} \) is the neighbor of \( x_{(2,r)} \) and \( x_{(q,r)} \). Therefore, using equations (15) and (16), we have

\[
d(x_{(1,1)}, x_{(1,r)}) = d(x_{(1,1)}, x_{(2,r)}) = r, \quad 1 \leq r \leq \frac{p - 1}{2} + 1.
\]

(17)

\[
d(x_{(1,1)}, x_{(q,r)}) = d(x_{(1,1)}, x_{(2,r)}) = p - 2 - r, \quad \frac{p - 1}{2} + 2 \leq r \leq p.
\]

(18)
Further, $x_{(2,r)}$ is the neighbor of $x_{(3,r)}$ and $x_{(q,r)}$ is the neighbor of $x_{(q-1,r)}$, and so equations (17) and (18) imply

$$d(x_{(1,1)}, x_{(q-1,r)}) = d(x_{(1,1)}, x_{(3,r)}) = 1 + r, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q-1,r)}) = d(x_{(1,1)}, x_{(3,r)}) = p + 3 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (19)$$

Moreover, $x_{(3,r)}$ is the neighbor of $x_{(4,r)}$ and $x_{(q-1,r)}$ is the neighbor of $x_{(q-2,r)}$. Hence,

$$d(x_{(1,1)}, x_{(q-2,r)}) = d(x_{(1,1)}, x_{(4,r)}) = 2 + r, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q-2,r)}) = d(x_{(1,1)}, x_{(4,r)}) = p + 4 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (20)$$

Continuing the same procedure, after $(q - 1/2)$th steps, the vertex $x_{(q-1/2,r)}$ would be in the neighbor of $x_{(q+1/2,r)}$. Therefore,

$$d(x_{(1,1)}, x_{(q+1/2,r)}) = \frac{q-1}{2} + r - 1, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q+1/2,r)}) = p - r + \frac{q-1}{2} + 1, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (21)$$

It means that $x_{(q+1/2, (p-1/2)+1)}$ is farthest from $x_{(1,1)}$. Therefore, ec$(x_{(1,1)}) = (p - 1/2) + (q - 1/2)$. Similarly, $x_{(q+3/2, (p-1/2)+1)}, x_{(q+5/2, (p-1/2)+1)}, x_{(q+7/2, (p-1/2)+1)}, \ldots, x_{(q+1/2, (p-1/2)+1)}$ are farthest vertices from $x_{(2,1)}, x_{(3,1)}, x_{(4,1)}, \ldots, x_{(q,1)}$, respectively.

Hence, $\text{rad}(C_p \sqcap C_q) = \text{diam}(C_p \sqcap C_q) = (p - 1/2) + (q - 1/2)$.

Case 3. Consider $p \equiv 1 \pmod{2}$, $q \equiv 0 \pmod{2}$, $p \geq q$.

The vertices $x_{(1,1)}$ and $x_{(1,r)}$ on the cycle $(x_{(1,1)}, x_{(1,2)}, x_{(1,3)}, \ldots, x_{(1,r)}, \ldots, x_{(1,p)})$ have the following distances:

$$d(x_{(1,1)}, x_{(1,r)}) = r - 1, \quad 1 \leq r \leq \frac{p-1}{2} + 1, \quad (22)$$

$$d(x_{(1,1)}, x_{(1,r)}) = p + 1 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (23)$$

This means to locate the farthest vertex from $x_{(1,1)}$ in $C_p \sqcap C_q$, only these values $1 \leq r \leq (p - 1/2) + 1$ are considered.

Since, each $x_{(1,r)}$ is the neighbor of $x_{(2,r)}$ and $x_{(q,r)}$, therefore using equations (22) and (23), we have

$$d(x_{(1,1)}, x_{(q,r)}) = d(x_{(1,1)}, x_{(2,r)}) = r, \quad 1 \leq r \leq \frac{p-1}{2} + 1, \quad (24)$$

$$d(x_{(1,1)}, x_{(q,r)}) = d(x_{(1,1)}, x_{(2,r)}) = p + 2 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (25)$$

Further, $x_{(2,r)}$ is the neighbor of $x_{(3,r)}$ and $x_{(q,r)}$ is the neighbor of $x_{(q-1,r)}$. Therefore, equations (24) and (25) give

$$d(x_{(1,1)}, x_{(q-1,r)}) = d(x_{(1,1)}, x_{(3,r)}) = r + 1, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q-1,r)}) = d(x_{(1,1)}, x_{(3,r)}) = p + 3 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (26)$$

Moreover, $x_{(3,r)}$ is the neighbor of $x_{(4,r)}$ and $x_{(q-1,r)}$ is the neighbor of $x_{(q-2,r)}$. Therefore,

$$d(x_{(1,1)}, x_{(q-2,r)}) = d(x_{(1,1)}, x_{(4,r)}) = r + 2, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q-2,r)}) = d(x_{(1,1)}, x_{(4,r)}) = p + 4 - r, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (27)$$

Continuing the same procedure, after $(q/2)$th steps, the vertex $x_{(q/2,1)}$ would be in the neighbor of $x_{(q/2)+1,i}$ which implies

$$d(x_{(1,1)}, x_{(q/2+1,i)}) = r + \frac{q-1}{2}, \quad 1 \leq r \leq \frac{p-1}{2} + 1,$$

$$d(x_{(1,1)}, x_{(q/2+1,i)}) = p - r + \frac{q+1}{2}, \quad \frac{p-1}{2} + 2 \leq r \leq p. \quad (28)$$
It means \( x_{(q/2)+1,(p-1/2)+1} \) is farthest from \( x_{(1,1)} \). Therefore, \( ec(x_{(1,1)}) = (p - 1/2) + (q/2) \). Similarly, farthest vertices from \( x_{(2,r)} \), \( x_{(3,r)} \), \ldots, \( x_{(q,r)} \) are \( x_{((q/2)+2,(p-1/2)+1)}, x_{((n/2)+3,(p-1/2)+1)}, \ldots, x_{((q/2),(p-1/2)+1)} \), respectively. Hence, \( rad(C_p \square C_q) = diam(C_p \square C_q) = (p - 1/2) + (q/2) \).

**Case 4.** Consider \( p \equiv 0 \text{ (mod2)}, q \equiv 1 \text{ (mod2)}, p \geq q \).

Since the graph is symmetric, by switching the roles of \( p \) and \( q \), we get the same case as Case 3. Therefore, we have discussed all the possible cases.

Now, it is concluded that the family of \( C_p \square C_q \) is self-centered for all possible values of \( p \) and \( q \).

2.1. Illustration. Consider the graph \( C_6 \square C_3 \) shown in Figure 2 in which eccentricity of every vertex is shown by blue circled numbers. Clearly from Figure 2, all vertices have eccentricity 3. Therefore, the center and periphery for \( C_6 \square C_3 \) is the graph itself.

3. Average Eccentricity of \( C_p \square C_q \), \( p \geq q \)

There are \( q \) circles in \( C_q \square C_q \) and each has \( p \) vertices. Thus, the total number of vertices in \( C_q \square C_q \) is equal to the product of \( p \) and \( q \). The average eccentricity of \( C_p \square C_q \), \( p \geq q \), will be discussed in the following cases:

**Case 1.** For \( \{p, q\} \equiv 0 \text{ (mod2)} \),

\[
\text{Avgec}(C_p \square C_q) = \frac{1}{pq} \sum_{x \in V(M)} ec_M(x) = \frac{1}{pq} \left\{ q \left\{ p \left( \frac{p + q}{2} \right) \right\} \right\} = \frac{p + q}{2}. 
\]

**Case 2.** When \( \{p, q\} \equiv 1 \text{ (mod2)} \),

\[
\text{Avgec}(C_p \square C_q) = \frac{1}{pq} \sum_{x \in V(M)} ec_M(x) = \frac{1}{pq} \left\{ q \left\{ p \left( \frac{p + q - 2}{2} \right) \right\} \right\} = \frac{p + q - 2}{2}. 
\]

**Case 3.** When \( p \equiv 1 \) and \( q \equiv 0 \text{ (mod2)} \),

\[
\text{Avgec}(C_p \square C_q) = \frac{1}{pq} \sum_{x \in V(M)} ec_M(x) = \frac{1}{pq} \left\{ q \left\{ p \left( \frac{p + q - 1}{2} \right) \right\} \right\} = \frac{p + q - 1}{2}. 
\]

Therefore,

\[
\text{Avgec}(C_p \square C_q) = \begin{cases} 
\frac{q + p}{2}, & [p, q] \equiv 0 \text{ (mod2)}, \\
\frac{q + p - 1}{2}, & p \equiv 1 \text{ and } q \equiv 0 \text{ (mod2)}, \\
\frac{q - 1 + p - 1}{2}, & [p, q] \equiv 1 \text{ (mod2)}. 
\end{cases}
\]

**Case 4.** When \( p \equiv 1 \) and \( q \equiv 0 \text{ (mod2)} \).

Since the graph is symmetric, by changing the roles of \( p \) and \( q \), we get Case 3.

Thus, all cases have been discussed, and our result is completed.

4. Conclusion

In this manuscript, we have discussed distance measurements including center, periphery, and average eccentricity for the Cartesian product of two cycles. The results are obtained using the definitions of eccentricity, radius, and diameter of a graph, and all possible cases (for different parity of length of cycles) have been proved. One of the attributes in applications of graph theory is distance and its related measurements in the graph. Weighted distance, topological distance, eccentricity, radius, diameter, metric dimension, indices, etc., are such distance-related terms and have received much attention of researchers. Along with the distance, the graph operations make the structures somewhat similar to the practical situation. One can be interested in researching the distance-related measurements for different graph operations: corona product, strong product, lexicographic product, etc. The other direction might be of
extending these results for one point union of graphs with nonisomorphic copies.

Data Availability

All data used for preparation of this manuscript are listed in references.

Conflicts of Interest

The authors declare that there are no conflicts of interest.

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