

Research Article On Soft Quantum B-Algebras and Fuzzy Soft Quantum B-Algebras

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This paper aims to make a combination between the quantum B-algebras (briefly, \mathcal{X} -As) and two interesting theories (e.g., soft set theory and fuzzy soft set theory). Firstly, we propose the novel notions of soft quantum B-algebras (briefly, \mathbb{SQB} -As), a soft deductive system of \mathbb{QB} -As, and deducible soft quantum B-algebras (briefly, \mathbb{DSQB} -As). Then, we discuss the relationship between \mathbb{SQB} -As and \mathbb{DSQB} -As. Furthermore, we investigate the union and intersection operations of \mathbb{DSQB} -As. Secondly, we introduce the notions of a fuzzy soft quantum B-algebras (briefly, \mathbb{FSQB} -As), a fuzzy soft deductive system of \mathbb{QB} -As, and present some characterizations of \mathbb{FSQB} -As, along with several examples. Finally, we explain the basic properties of homomorphism image of \mathbb{FSQB} -As.

1. Introduction

In 1999, Molodtsov [1] introduced the notion called soft sets (briefly, SS) (i.e., which reduce the uncertainty and vagueness of knowledge). Maji et al. [2] presented the fuzzy soft sets (briefly, FSS). Since then, many researchers studied further on SS and FSS as in the following published articles (e.g., [3–9]).

In 2014, Rump and Yang [10] proposed the notion of \mathbb{QB} -As (i.e., a partial ordered implication algebras). Rump [11, 12] investigated many implication algebras (for example, pseudo-BCK-algebras, po-groups, BL-algebras, MV-algebras, GPE-algebras, and resituated lattices). Botur and Paseka [13] studied filters on integral QB-As, and Zhang et al. [14] established the quotient structures by using q-filters in QB-As and investigated the relation between basic implication algebras and QB-As. Han et al. [15] constructed the unitality of QB-As and explained the injective hulls of QB-As in [16]. By the framework of QB-As, there are many published papers on QB-As (e.g., [17–23]).

Regarding these developments, as the motivation of this paper, we will combine QB-As with SS and FSS (i.e., enrich the previous work on hybrid soft set and fuzzy soft set theories algebras with quantum structures). We introduce the notions of SQB-As and the soft deductive system of QB-As and consider the relation between SQB-As and DSQB-As. Furthermore, some conditions are given to ensure the operations union and intersection holds of soft deductive of QB-As. Then, we investigate the homomorphism image of deductive SQB-As. Lastly, we define FSQB-As and fuzzy soft deductive system of QB-As and give an example to illustrate its derive properties.

In the following, we have arranged the sections as follows. In Section 2, we briefly recall many notions related to QB-As, SS, and FSS as indicated in Definitions 1–7, which are used in the sequel. In Section 3, we propose the notions of SQB-As, soft deductive system of QB-As, and DSQB-As. In Section 4, we present the notions of FSQB-As and a fuzzy soft deductive system of QB-As and discuss the homomorphism image of FSQB-As. The conclusions are explained in Section 5.

2. Preliminaries

We give some basic notions of QB-As, SS, and FSS before defining QB-As in Section 3.

Definition 1 (cf. [10]).

(1) QB-As is a partially ordered set (X, ≤) with two binary operations → and → which satisfy (∀x, y, z ∈ X):

$$y \longrightarrow z \le (x \longrightarrow y) \longrightarrow (x \longrightarrow z),$$

$$y \longrightarrow z \le (x \longrightarrow y) \longrightarrow (x \longrightarrow z),$$

$$y \le z \Longrightarrow x \longrightarrow y \le x \longrightarrow z,$$

$$x \le y \longrightarrow z \Longleftrightarrow y \le x \longrightarrow z.$$
(1)

- (2) QB-A is a commutative (briefly, CQB-A) if $x \longrightarrow y = x \dashrightarrow y$ ($\forall x, y \in \mathcal{X}$).
- (3) A subset \mathcal{Y} of a QB-A \mathcal{X} is a subalgebra if $x \longrightarrow y, x \dashrightarrow y \in \mathcal{Y} \ (\forall x, y \in \mathcal{X}).$

In what follows, denote by $\mathcal X$ a QB-A unless otherwise specified.

Definition 2 (cf. [10]). Let \mathcal{X}_1 and \mathcal{X}_2 be two QB-As. Then, $\psi: \mathcal{X}_1 \longrightarrow \mathcal{X}_2$ is a morphism of QB-As if it satisfies $(\forall x, y \in \mathcal{X}_1)$:

$$\psi(x \longrightarrow y) \le \psi(x) \longrightarrow \psi(y),$$

$$\psi(x \dashrightarrow y) \le \psi(x) \dashrightarrow \psi(y).$$
(2)

We say morphism ψ is exact if the inequalities become equations.

Definition 3 (cf. [1]). Assume that \mathscr{X} be a set and \mathscr{K} be a set of parameters. $\mathscr{S}_{\mathscr{K}}$ (called SS) is a mapping given by $\mathscr{S}: \mathscr{K} \longrightarrow 2^{\mathscr{X}}$ (i.e., $2^{\mathscr{X}}$ is the power set of \mathscr{X}).

Definition 4 (cf. [3]). Assume that $\mathscr{S}_{\mathscr{H}_1}$ and $\mathscr{S}_{\mathscr{H}_2}$ are two SS over \mathscr{X} . $\mathscr{S}_{\mathscr{H}_1}$ is a subset of $\mathscr{S}_{\mathscr{H}_2}$ (denoted by $\mathscr{S}_{\mathscr{H}_1} \subset \mathscr{S}_{\mathscr{H}_2}$) if

- (1) $\mathscr{K}_1 \subset \mathscr{K}_2$
- (2) For every $k \in \mathcal{K}_1$, $S_{\mathcal{K}_1}(k)$ and $S_{\mathcal{K}_2}(k)$ are identical approximations

Definition 5 (cf. [3]). Assume that $\mathscr{S}_{\mathscr{H}_1}, \mathscr{S}_{\mathscr{H}_2}$, and $\mathscr{S}_{\mathscr{H}_3}$ are three SS over \mathscr{X} . $\mathscr{S}_{\mathscr{H}_3}$ is the intersection of $\mathscr{S}_{\mathscr{H}_1}$ and $\mathscr{S}_{\mathscr{H}_2}$ (denoted by $\mathscr{S}_{\mathscr{H}_3} = \mathscr{S}_{\mathscr{H}_1} \cap \mathscr{S}_{\mathscr{H}_2}$) if

- (1) $\mathcal{K}_3 = \mathcal{K}_1 \cap \mathcal{K}_2$
- (2) $\forall k \in \mathcal{K}_3, S_{\mathcal{K}_3}(k) = S_{\mathcal{K}_1}(k) \text{ or } S_{\mathcal{K}_2}(k)$ (as both are same sets)

Definition 6 (cf. [3]). Assume that $\mathscr{S}_{\mathscr{H}_1}, \mathscr{S}_{\mathscr{H}_2}$, and $\mathscr{S}_{\mathscr{H}_3}$ are three SS over \mathscr{X} . $\mathscr{S}_{\mathscr{H}_3}$ is called the union of $\mathscr{S}_{\mathscr{H}_1}$ and $\mathscr{S}_{\mathscr{H}_2}$ (denoted by $\mathscr{S}_{\mathscr{H}_3} = \mathscr{S}_{\mathscr{H}_1} \cup \mathscr{S}_{\mathscr{H}_2}$) if

(1)
$$\mathscr{K}_3 = \mathscr{K}_1 \cup \mathscr{K}_2$$
.

(2)
$$k \in \mathcal{H}_{3}$$
,

$$\mathcal{S}_{\mathcal{H}_{3}}(k) = \begin{cases} \mathcal{S}_{\mathcal{H}_{1}}(k), & k \in \mathcal{H}_{1} \setminus \mathcal{H}_{2} \\ \mathcal{S}_{\mathcal{H}_{2}}(k), & k \in \mathcal{H}_{2} \setminus \mathcal{H}_{1}, \\ \mathcal{S}_{\mathcal{H}_{1}}(k) \cup \mathcal{S}_{\mathcal{H}_{2}}(k), & k \in \mathcal{H}_{1} \cap \mathcal{H}_{2} \end{cases}$$
(3)

Definition 7 (cf. [2]). FSS (called FSS) $\hat{\mathscr{S}}_{\mathscr{H}}$ is a mapping given by $\hat{\mathscr{S}}: \mathscr{H} \longrightarrow I^X$ (i.e., I^X is the set of all fuzzy sets [24] of \mathscr{X}).

3. SQB-As

We define the \mathbb{SQB} -As and give several examples based on \mathbb{SQB} -As. Also, we will study the union and intersection operations between two \mathbb{SQB} -As as follows.

Definition 8. $S_{\mathcal{H}}$ is a SQB-As over \mathcal{X} if $S_{\mathcal{H}}(x) (\forall x \in \mathcal{K})$ are subalgebras of \mathcal{X} (i.e., in case $\mathcal{K} = \mathcal{X}$).

Example 1

(1) Suppose X (i.e., X = {k₁, k₂, k₃, 1}) with the order k₂, k₃ < k₁ < 1. Now, we show, by Table 1, the binary operation →.
Clearly X is a COB A We define S (∀x ∈ X)

Clearly, \mathscr{X} is a CQB-A. We define $\mathscr{S}_{\mathscr{H}}(\forall x \in \mathscr{K})$ (i.e., $\mathscr{K} = \mathscr{X}$) by

$$\mathscr{S}_{\mathscr{H}}(x) = \{ y \in \mathscr{X} | (x \longrightarrow y) \longrightarrow y \in \{k_1, 1\} \}.$$
(4)

From Table 1, we can get on $\mathscr{S}_{\mathscr{K}}(k_1) = \mathscr{X}, \ \mathscr{S}_{\mathscr{K}}(k_2) = \mathscr{S}_{\mathscr{H}}, \ (k_3) = \{k_1, k_3, 1\}, \text{ and } \mathscr{S}_{\mathscr{H}}(1) = \mathscr{X}, \text{ and then,} \\ \mathscr{S}_{\mathscr{H}}(x)(x \in \mathscr{H}) \text{ are all subalgebras of } \mathscr{X}. \text{ Consequently, } \mathscr{S}_{\mathscr{H}} \text{ is a } \mathbb{SQB}\text{-}As \text{ over } \mathscr{X}.$

(2) Suppose \mathscr{X} (i.e., $\mathscr{X} = \{k_1, k_2, k_3, 1\}$) with the order $k_1 < k_2 < k_3 < 1$. Now, we show, by Table 2, the binary operation \longrightarrow .

Clearly, \mathscr{X} is a CQB-A. We define $\mathscr{S}_{\mathscr{H}}(\forall x \in \mathscr{K})$ (i.e., $\mathscr{K} = \mathscr{X}$) by

$$\mathcal{S}_{\mathscr{K}}(x) = \{ y \in \mathscr{X} | x \mathscr{R} y \Longleftrightarrow x \longrightarrow (x \longrightarrow y) \in \{k_3, 1\} \}.$$
(5)

From Table 2, we can get on (HTML translation failed), and then, $\mathcal{S}_{\mathcal{H}}(x)(x \in \mathcal{K})$ are all subalgebras of \mathcal{X} . Consequently, $\mathcal{S}_{\mathcal{H}}$ is a SQB-As over \mathcal{X} .

We ensure the operations (i.e., union and intersection) are holding on \mathbb{SQB} -As by the following suggested theorem.

Theorem 1. Assume that $S_{\mathcal{H}_1}$ and $S_{\mathcal{H}_2}$ are \mathbb{SQB} -As over \mathcal{X} . Then,

- (1) If $\mathscr{K}_3 = \mathscr{K}_1 \cap \mathscr{K}_2$, then $\mathscr{S}_{\mathscr{K}_3} = \mathscr{S}_{\mathscr{K}_1} \cap \mathscr{S}_{\mathscr{K}_2}$ is called a $\mathbb{SQB-A}$ over \mathscr{X}
- (2) If $\mathscr{K}_1 \cap \mathscr{K}_2 = \emptyset$, then $\mathscr{S}_{\mathscr{K}_1} \cup \mathscr{S}_{\mathscr{K}_2}$ is called a SQB-A over \mathscr{X}

Journal of Mathematics

TABLE 1: The binary operation \longrightarrow .

\longrightarrow	k_1	k_2	k_3	1
k_1	1	k_1	k_1	1
k_2	1	1	k_1	1
k_3	1	1	k_1	1
1	k_1	k_2	k_3	1

TABLE 2: The binary operation \longrightarrow .

\longrightarrow	k_1	k_2	k_3	1
k_1	1	1	1	1
k_2	k_1	k_2	1	1
k_3	k_1	k_1	1	1
1	k_1	k_1	k_3	1

Proof

- (1) If $\mathscr{K}_3 = \mathscr{K}_1 \cap \mathscr{K}_2$ and by Definition 5, we obtain $\mathscr{S}_{\mathscr{K}_3}(x) = \mathscr{S}_{\mathscr{K}_1}(x)$ or $\mathscr{S}_{\mathscr{K}_3}(x) = \mathscr{S}_{\mathscr{K}_2}(x)$, for all $x \in \mathscr{K}_3$. Since $\mathscr{S}_{\mathscr{K}_1}$ and $\mathscr{S}_{\mathscr{K}_2}$ are SQB-As over \mathscr{X} , which implies that $\mathscr{S}_{\mathscr{K}_3}$ is a SQB-As over \mathscr{X} , that is, $\mathscr{S}_{\mathscr{K}_3}(x) = \mathscr{S}_{\mathscr{K}_1}(x)$ or $\mathscr{S}_{\mathscr{K}_3}(x) = \mathscr{S}_{\mathscr{K}_2}(x)$ are both subalgebras of $\mathscr{X} (\in \mathscr{K}_3)$, therefore, $\mathscr{S}_{\mathscr{K}_3} = \mathscr{S}_{\mathscr{K}_1} \cap \mathscr{S}_{\mathscr{K}_2}$ is a SQB-A over \mathscr{X} .
- (2) If $\mathscr{K}_3 = \mathscr{K}_1 \cup \mathscr{K}_2$ and by Definition 6, we obtain

$$\mathcal{S}_{\mathcal{H}_{3}}(x) = \begin{cases} \mathcal{S}_{\mathcal{H}_{1}}(x), & x \in \mathcal{H}_{1} \backslash \mathcal{H}_{2}, \\ \mathcal{S}_{\mathcal{H}_{2}}(x), & x \in \mathcal{H}_{2} \backslash \mathcal{H}_{1}, \\ \mathcal{S}_{\mathcal{H}_{1}}(x) \cup \mathcal{S}_{\mathcal{H}_{2}}(x), & x \in \mathcal{H}_{1} \cap \mathcal{H}_{2} \end{cases}$$
(6)

For $x \in \mathcal{K}_1 \setminus \mathcal{K}_2$ and since $\mathcal{S}_{\mathcal{H}_1}$ is a \mathbb{SQB} - \mathbb{A} , then we have $\mathcal{S}_{\mathcal{H}_3}(x) = \mathcal{S}_{\mathcal{H}_1}(x)$ is a subalgebra of \mathcal{X} . Similarly, for $x \in \mathcal{H}_2 \setminus \mathcal{H}_1$, then $\mathcal{S}_{\mathcal{H}_3}(x) = \mathcal{S}_{\mathcal{H}_2}(x)$ is a subalgebra of \mathcal{X} due to $\mathcal{S}_{\mathcal{H}_2}$ is a \mathbb{SQB} - \mathbb{A} . Again, for $\mathcal{H}_1 \cap \mathcal{H}_2 = \emptyset$, so $x \in \mathcal{H}_1 \cap \mathcal{H}_2$ or $x \in \mathcal{H}_2 \cap \mathcal{H}_1$, for all $x \in \mathcal{H}_3$. Thus, $\mathcal{S}_{\mathcal{H}_3} = \mathcal{S}_{\mathcal{H}_1} \cup \mathcal{S}_{\mathcal{H}_2}$ is a \mathbb{SQB} - \mathbb{A} over \mathcal{X} .

Remark 1. If $\mathscr{K}_1 \cap \mathscr{K}_2 \neq \emptyset$, then Theorem 1 (2) does not hold by the following example.

Example 2. Suppose \mathscr{X} (i.e., $\mathscr{X} = \{0, k_1, k_2, k_3, k_4, 1\}$). Now, we show, by Tables 3 and 4, the binary operations \longrightarrow and \rightsquigarrow , respectively.

Clearly, \mathcal{X} is a CQB-A. Then,

(i) We define
$$\mathscr{S}_{\mathscr{H}_{1}}(\forall x \in \mathscr{H}_{1})$$
 (i.e., $\mathscr{H}_{1} = \mathscr{X}$) by
 $\mathscr{S}_{\mathscr{H}_{1}}(x) = \{ y \in \mathscr{X} | x \mathscr{R} y \Longleftrightarrow x \longrightarrow (x \longrightarrow y) x$
 $\longrightarrow (x \leadsto y) \in \{ k_{3}, k_{4}, 1 \} \}.$
(7)

From Table 3, we can get $\mathscr{S}_{\mathscr{H}_1}(0) = \mathscr{X}$ and $\mathscr{S}_{\mathscr{H}_1}(k_1) = \mathscr{S}_{\mathscr{H}_1}(k_2) = \mathscr{S}_{\mathscr{H}_1}(k_3) = \mathscr{S}_{\mathscr{H}_1}(k_4) = \mathscr{S}_{\mathscr{H}_1}(1) = \{k_3, k_4, 1\}$, and then, $\mathscr{S}_{\mathscr{H}_1}(x) (x \in \mathscr{H}_1)$ are all subalgebras of \mathscr{X} . Consequently, $\mathscr{S}_{\mathscr{H}_1}$ is a SQB-As over \mathscr{X} .

TABLE 3: The binary operation \longrightarrow .

\longrightarrow	0	k_1	(HTML translation failed)	k_3	k_4	1
0	1	1	1	1	1	1
k_1	0	k_2	0	k_4	1	1
k_2	0	0	k_2	k_4	k_4	1
k_3	0	0	0	1	1	1
k_4	0	0	0	k_4	1	1
1	0	0	0	k_4	k_4	1

>	0	k_1	k_2	k_3	k_4	1
0	1	1	1	1	1	1
k_1	0	0	0	1	1	1
k_2	0	k_1	k_2	k_3	k_4	1
k_3	0	0	0	1	1	1
k_4	0	0	0	1	1	1
1	0	0	0	k_3	k_4	1

(ii) We define
$$S_{\mathcal{H}_2} (\forall x \in \mathcal{H}_2)$$
 (i.e., $\mathcal{H}_2 = \{k_2\}$) by
 $S_{\mathcal{H}_2}(x) = \{y \in \mathcal{H}_2 | x \mathcal{R} y \Longleftrightarrow x \longrightarrow y = k_2, x \leadsto y = k_2\}.$
(8)

From Table 4, we can get $\mathscr{S}_{\mathscr{K}_2}(k_2) = \{k_2\}$ is the subalgebra of \mathscr{X} . Consequently, $\mathscr{S}_{\mathscr{K}_2}$ is a SQB-As over \mathscr{X} .

From (i) and (ii) and $\mathscr{K}_1 \cap \mathscr{K}_2 = \{k_2\} \neq \emptyset$, then we have $\mathscr{S}_{\mathscr{K}_3}(k_2) = \mathscr{S}_{\mathscr{K}_1}(k_2) \cup \mathscr{S}_{\mathscr{K}_2}(k_2) = \{k_3, k_4, 1\} \cup \{k_2\} = \{k_2, k_3, k_4, 1\}$ is not a subalgebra over \mathscr{X} . Thus, $\mathscr{S}_{\mathscr{K}_3}$ is not a \mathbb{SQB} -A.

3.1. Soft Deductive Systems of \mathbb{QB} -As. Based on Definition 8, we will propose the notion of soft deductive systems of \mathbb{QB} -As as indicated below.

Definition 9. Assume that $\mathcal{X} = (\mathcal{X}, \dots, \mathcal{N}, \leq)$ be a \mathbb{SQB} -A. A nonempty subset $\mathcal{D} \subseteq \mathcal{X}$ is a deductive system of \mathcal{X} if it satisfies

(1) $\forall x \in \mathcal{D}, x \longrightarrow x \in \mathcal{D}, x \rightsquigarrow x \in \mathcal{D}$ (2) $\forall x, y \in \mathcal{X}, x \in \mathcal{D}, x \longrightarrow y \in \mathcal{D} \Longrightarrow y \in \mathcal{D}$

Definition 10. Let \mathcal{X} be a \mathbb{QB} -A and \mathcal{Y} a subalgebra of \mathcal{X} . A subset \mathcal{D} of \mathcal{X} is a deductive system of \mathcal{X} related to \mathcal{Y} (i.e., \mathcal{Y} -deductive system of \mathcal{X}), denoted by $\mathcal{D} \bowtie \mathcal{Y}$, and satisfies the following two conditions:

(1)
$$\forall x \in \mathcal{D}, x \longrightarrow x \in \mathcal{D}, x \rightsquigarrow x \in \mathcal{D}$$

(2) $\forall y \in \mathcal{Y}, x \in \mathcal{D}, x \longrightarrow y \in \mathcal{D} \Longrightarrow y \in \mathcal{D}$

Remark 2. According to Definitions 9 and 10, we obtain that any deductive system of \mathcal{X} is \mathcal{Y} -deductive system if \mathcal{Y} is a subalgebra of \mathcal{X} .

The converse of Remark 2 does not hold by Example 3 (i.e., \mathcal{Y} is a subalgebra of \mathcal{X} and \mathcal{Y} -deductive system is not a deductive system).

Example 3. Suppose \mathcal{X} (i.e., $\mathcal{X} = \{0, k_1, k_2, k_3, 1\}$) with partial order $0 < k_1 < k_3 < 1$ and $0 < k_1 < k_2 < 1$. Now, we show, by Tables 5 and 6, the binary operations \longrightarrow and \rightsquigarrow , respectively.

Clearly, \mathscr{X} is a \mathbb{CQB} -A. Consider a subalgebra $\mathscr{Y} = \{k_1, 1\}$ and a subset $\mathscr{D} = \{k_1, k_2, 1\}$; we can see that $\mathscr{D} \bowtie \mathscr{Y}$. However, \mathscr{D} is not a deductive system of \mathscr{X} since $k_3 \longrightarrow 1 = 1 \in \mathscr{D}$ and $k_3 \notin \mathscr{D}$.

Definition 11. Assume that $\mathscr{S}_{\mathscr{X}}$ is a \mathbb{QB} - \mathbb{A} over \mathscr{X} . $\mathscr{S}_{\mathscr{D}}$ (i.e., \mathbb{S}) over \mathscr{X} is a soft deductive system of $\mathscr{S}_{\mathscr{X}}$, denoted by $\mathscr{S}_{\mathscr{D}} \bowtie \mathscr{S}_{\mathscr{X}}$, and satisfies the following two conditions:

- (1) $\mathcal{D} \subseteq \mathcal{K}$
- (2) $\forall x \in \mathcal{D}, S_{\mathcal{D}}(x) \triangleright \triangleleft S_{\mathcal{K}}(x)$

Now, we will give an example to illustrate Definition 11 as follows.

Example 4. Suppose \mathcal{X} (i.e., $\mathcal{X} = \{k_1, k_2, k_3, k_4, 1\}$) with partial order $k_1 < k_2 < k_3 < k_4 < 1$. Now, we show, by Tables 7 and 8, the binary operations \longrightarrow and \rightsquigarrow , respectively.

Clearly, \mathscr{X} is a CQB-A. We define $\mathscr{S}_{\mathscr{K}}(\forall x \in \mathscr{K})$ (i.e., $\mathscr{K} = \mathscr{X}$) by

$$\mathcal{S}_{\mathcal{K}}(x) = \{ y \in \mathcal{X} | x \mathcal{R} y \Longleftrightarrow (x \longrightarrow y) \leadsto y = 1 \}.$$
(9)

From Tables 7 and 8, we can get on $\mathcal{S}_{\mathcal{H}}(k_1) = \mathcal{S}_{\mathcal{H}}(k_2) = 1$, $\mathcal{S}_{\mathcal{H}}(k_3) = \{k_2, 1\}$, $\mathcal{S}_{\mathcal{H}}(k_4) = \{k_2, k_3, 1\}$, and $\mathcal{S}_{\mathcal{H}}(1) = \mathcal{X}$, and then, $\mathcal{S}_{\mathcal{H}}(x) (x \in \mathcal{H})$ are all subalgebras of \mathcal{X} . Consequently, $\mathcal{S}_{\mathcal{H}}$ is a SQB-As over \mathcal{X} .

Next, for a subset $\mathcal{D} = \{k_2, k_4\}$, we define $\mathcal{S}_{\mathcal{D}} (\forall x \in \mathcal{D})$ by

$$\mathcal{S}_{\mathcal{D}}(x) = \{1\} \cup \{y \in \mathcal{X} | y \le x\}.$$

$$(10)$$

Then, we obtain $\mathscr{S}_{\mathscr{D}}(k_2) = \{k_1, k_2, 1\} \triangleright \triangleleft \{1\} = \mathscr{S}_{\mathscr{K}}(k_2)$ and $\mathscr{S}_{\mathscr{D}}(k_4) = \mathscr{X} \triangleright \triangleleft \{k_2, k_3, 1\} = \mathscr{S}_{\mathscr{K}}(k_4)$. Consequently, $\mathscr{S}_{\mathscr{D}}$ is a soft deductive system of $\mathscr{S}_{\mathscr{K}}$.

Theorem 2. Assume that $S_{\mathcal{X}}$ is a SQB-A over \mathcal{X} and $S_{\mathcal{D}_1}$ and $S_{\mathcal{D}_2}$ are two SS. Then,

(1) If
$$\mathscr{D}_{1} \cap \mathscr{D}_{2} \neq \emptyset$$
, then $\mathscr{S}_{\mathscr{D}_{1}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}}, \mathscr{S}_{\mathscr{D}_{2}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}}$
 $\mathscr{S}_{\mathscr{R}} \Longrightarrow \mathscr{S}_{\mathscr{D}_{1}} \cap \mathscr{S}_{\mathscr{D}_{2}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}}$
(2) If $\mathscr{D}_{1} \cap \mathscr{D}_{2} = \emptyset$, then $\mathscr{S}_{\mathscr{D}_{1}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}}, \mathscr{S}_{\mathscr{D}_{2}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}} \Longrightarrow$
 $\mathscr{S}_{\mathscr{D}_{1}} \cup \mathscr{S}_{\mathscr{D}_{2}} \vdash \overrightarrow{\operatorname{dS}}_{\mathscr{R}}$

Proof

- (1) Follow from Definition 5.
- (2) If $\mathscr{S}_{\mathscr{D}_1} \Join \mathscr{S}_{\mathscr{D}_2} \Join \mathscr{S}_{\mathscr{D}_2} \Join \mathscr{S}_{\mathscr{D}_2}$, then, by Definition 6, we have $\mathscr{D}_3 = \mathscr{D}_1 \cap \mathscr{D}_2$ (i.e., $x \in \mathscr{D}_3$), $\mathscr{S}_{\mathscr{D}_1} \cup \mathscr{S}_{\mathscr{D}_2} = \mathscr{S}_{\mathscr{D}_3}$, and

TABLE 5: The binary operation \longrightarrow .

\longrightarrow	0	k_1	k_2	k_3	1
0	1	1	1	1	1
k_1	0	1	k_2	1	1
k_2	k_1	k_1	1	1	1
k_3	0	k_1	k_2	1	1
1	0	k_1	k_2	k_3	1

-**>	0	k_1	k_2	k_3	1
0	1	1	1	1	1
k_1	k_2	1	k_2	1	1
k_2	0	k_1	1	1	1
k_3	0	k_1	k_2	1	1
1	0	k_1	k_2	k_3	1

TABLE 7: The binary operation \longrightarrow .

\longrightarrow	k_1	k_2	k_3	k_4	1
k_1	1	1	1	1	1
k_2	k_3	1	1	1	1
k_3	k_2	k_2	1	1	1
k_4	k_2	k_2	k_{c}	1	1
1	k_1^-	k_2	k_3	k_4	1

TABLE 8: The binary operation -----.

-**>	k_1	k_2	k_3	k_4	1
k_1	1	1	1	1	1
k_2	k_4	1	1	1	1
k_3	k_2	k_2	1	1	1
k_4	k_1	k_2	k_3	1	1
1	k_1	k_2	k_3	k_4	1

$$\mathcal{S}_{\mathcal{D}_{3}}(x) = \begin{cases} \mathcal{S}_{\mathcal{D}_{1}}(x), & x \in \mathcal{D}_{1} \backslash \mathcal{D}_{2} \\ \mathcal{S}_{\mathcal{D}_{2}}(x), & x \in \mathcal{D}_{2} \backslash \mathcal{D}_{1}, \\ \mathcal{S}_{\mathcal{D}_{1}}(x) \cup \mathcal{S}_{\mathcal{D}_{2}}(x), & x \in \mathcal{D}_{1} \cap \mathcal{D}_{2} \end{cases}$$
(11)

Since $\mathcal{D}_1 \cap \mathcal{D}_2 = \emptyset$, we obtain either $x \in \mathcal{D}_1 \setminus \mathcal{D}_2$ or $x \in \mathcal{D}_2 \setminus \mathcal{D}_1$. Then, we have the following:

Case 1: if $x \in \mathcal{D}_1 \setminus \mathcal{D}_2$, since $\mathcal{S}_{\mathcal{D}_1} \vdash \mathcal{A} \mathcal{S}_{\mathcal{H}}$, then $\mathcal{S}_{\mathcal{D}_3}(x) = \mathcal{S}_{\mathcal{D}_1}(x) \vdash \mathcal{A} \mathcal{S}_{\mathcal{H}}(x)$ Case 2: if $x \in \mathcal{D}_2 \setminus \mathcal{D}_1$ and $\mathcal{S}_{\mathcal{D}_2} \vdash \mathcal{A} \mathcal{S}_{\mathcal{H}}$, then $\mathcal{S}_{\mathcal{D}_3}(x) = \mathcal{S}_{\mathcal{D}_3}(x) \vdash \mathcal{A} \mathcal{S}_{\mathcal{H}}(x)$

Consequently, for all $x \in \mathcal{D}_3$, we have $\mathcal{S}_{\mathcal{D}_3}(x) \triangleright \triangleleft \mathcal{S}_{\mathcal{K}}(x)$, which implies that $\mathcal{S}_{\mathcal{D}_1} \cup \mathcal{S}_{\mathcal{D}_2} = \mathcal{S}_{\mathcal{D}_3} \widecheck{\triangleright} \triangleleft \mathcal{S}_{\mathcal{K}}$. \Box

Remark 3. If $\mathscr{K}_1 \cap \mathscr{K}_2 \neq \emptyset$, then Theorem 2 (2) does not hold by the following example.

Example 5. Suppose \mathcal{X} (i.e., $\mathcal{X} = \{0, k_1, k_2, k_3, k_4, 1\}$). Now, we show, by Table 9, the binary operations \longrightarrow .

0

1

 k_3 0 k_1 k_2 k_4 1 1 0 1 1 1 1 1 k_1 k_3 1 k_2 k_3 k_2 1 k_2 k_4 k_2 k_1 1 k_1 1 k_3 k_1 k_1 1 1 k_1 1 k_4 k_2 1 1 1 k_2 1

 k_2

 k_3

1

 k_4

TABLE 9: The binary operation \longrightarrow .

Clearly, \mathcal{X} is a CQB-A. Then,

 k_1

(i) We define $\mathcal{S}_{\mathcal{H}}(\forall x \in \mathcal{H})$ (i.e., $\mathcal{H} = \mathcal{X}$) by $\mathcal{S}_{\mathcal{H}}(x) = \{ y \in \mathcal{X} | x \mathcal{R} y \iff (x \longrightarrow y) \longrightarrow y \in \{k_1, k_2, 1\} \}.$ (12)

From Table 9, we can get $\mathscr{S}_{\mathscr{H}}(0) = \{k_1, k_2, 1\}$, $\mathscr{S}_{\mathscr{H}}(k_1) = \mathscr{X}, \mathscr{S}_{\mathscr{H}}(k_2) = \mathscr{S}_{\mathscr{H}}(k_3) = \{k_1, 1\}$, and $\mathscr{S}_{\mathscr{H}}(k_4) = \{k_1, k_2, k_3, 1\}, \mathscr{S}_{\mathscr{H}}(1) = \mathscr{X}$, and then, $\mathscr{S}_{\mathscr{H}}(x) (\forall x \in \mathscr{H})$ are all subalgebras of \mathscr{X} . Consequently, $\mathscr{S}_{\mathscr{H}}$ is a SQB-As over \mathscr{X} .

(ii) We define $\mathscr{S}_{\mathscr{K}_1}(\forall x \in \mathscr{K}_1)$ (i.e., $\mathscr{K}_1 = \{k_1, k_2, k_3\}$) by

$$\mathcal{S}_{\mathcal{H}_1}(x) = \{ y \in \mathcal{X} | x \mathcal{R} y \longleftrightarrow x \longrightarrow y = 1 \}.$$
(13)

Then, we can get $\mathcal{S}_{\mathcal{H}_1}(k_1) = \{k_1, 1\} \triangleright \triangleleft X = \mathcal{S}_{\mathcal{H}}(k_1)$, $\mathcal{S}_{\mathcal{H}_1}(k_2) = \{k_2, 1\} \triangleright \triangleleft \{a, 1\} = \mathcal{S}_{\mathcal{H}}(k_2)$, and $\mathcal{S}_{\mathcal{H}_1}(k_3) = \{k_2, k_3, 1\} \triangleright \triangleleft \{k_1, 1\}$. Therefore, $\mathcal{S}_{\mathcal{H}_1}$ is a soft deductive system over $\mathcal{S}_{\mathcal{H}}$.

(iii) We define
$$\mathscr{S}_{\mathscr{K}_2}(\forall x \in \mathscr{K}_2)$$
 (i.e., $\mathscr{K}_2 = \{k_1\}$) by
 $\mathscr{S}_{\mathscr{K}_2}(x) = \{y \in \mathscr{X} | x \mathscr{R} y \Longleftrightarrow y \longrightarrow x = k_1\}.$ (14)

Then, we can get $\mathscr{S}_{\mathscr{H}_2}(k_1) = \{k_2, k_3, 1\} \triangleright \triangleleft X = \mathscr{S}_{\mathscr{H}_2}(k_1)$. Therefore, $\mathscr{S}_{\mathscr{H}_2}$ is a soft deductive system over $\mathscr{S}_{\mathscr{H}}$.

From (i)–(iii), we have $\mathscr{S}_{\mathscr{K}_3} = \mathscr{S}_{\mathscr{K}_1} \cup \mathscr{S}_{\mathscr{K}_2}$ which is not a soft deductive system of $\mathscr{S}_{\mathscr{K}}$, where $\mathscr{S}_{\mathscr{K}_3}(k_1) = \mathscr{S}_{\mathscr{K}_1}(k_1) \cup \mathscr{S}_{\mathscr{K}_2}(k_1) = \{k_1, k_2, k_3, 1\}$ is not a $\mathscr{S}_{\mathscr{K}}(a)$ -deductive system because $k_2 \longrightarrow k_4 = k_1 \in \{k_1, k_2, k_3, 1\}$ and $k_4 \notin \{k_1, k_2, k_3, 1\}$.

3.2. DSQB-As. We will give the notion of DSQB-As and investigate homomorphism image of DSQB-As as indicated below.

Definition 12. Assume that $\mathscr{S}_{\mathscr{K}}$ is a $\mathbb{SQB-A}$ over \mathscr{X} . If $\mathscr{S}_{\mathscr{K}}(x)(\forall x \in \mathscr{K})$ is a deductive system of \mathscr{X} , then $\mathscr{S}_{\mathscr{K}}$ is called a $\mathbb{DSQB-A}$ over X.

Example 6 (continued from Example 1 (2)). Clearly, $S_{\mathscr{K}}$ is $\mathbb{DSQB-A}$ over \mathscr{X} .

Definition 13

Suppose X be a QB-A with the greatest element 1 (i.e., X just only a poset); for any x ∈ X, the order of element x is defined as

$$\mathcal{O}(x) = \min\left\{p, q \in N | x \xrightarrow{p} x = 1, x \xrightarrow{q} x = 1\right\}, \quad (i),$$
(15)

where *N* is a natural number and $x \longrightarrow^{p} x = (((x \longrightarrow x) \longrightarrow \cdots) \longrightarrow x), \quad x \xrightarrow{q} x = (((x \dashrightarrow x) \longrightarrow \cdots) \dashrightarrow x).$

(2) If $p, q \in N$ does not exist to satisfy the above condition (i), then $x (\forall x \in \mathcal{X})$ is called infinite order.

Remark 4. Assume that $S_{\mathscr{H}}$ and $S_{\mathscr{H}_1}$ be two SQB-As over \mathscr{X} such that $\mathscr{H}_1 \subseteq \mathscr{H} \subseteq \mathscr{X}$. If $S_{\mathscr{H}}$ is a DSQB-A over \mathscr{X} , then $S_{\mathscr{H}}$ is a DSQB-A.

 $\mathscr{S}_{\mathscr{K}_1}$ is a DSQB-A. The converse of Remark 4 does not hold by the following Example 7.

Example 7 (continued from Example 2). We define $S_{\mathscr{K}}(\forall x \in \mathscr{K})$ (i.e., $\mathscr{K} = \mathscr{X}$) by

$$\mathcal{S}_{\mathscr{K}}(x) = \{ y \in \mathscr{X} | \mathcal{O}(x) = \mathcal{O}(y) \}.$$
(16)

Then, we get on $\mathscr{S}_{\mathscr{R}}(0) = \mathscr{S}_{\mathscr{R}}(k_3) = \mathscr{S}_{\mathscr{R}}(k_4) = \mathscr{S}_{\mathscr{R}}(k_1) = \{0, k_3, k_4, 1\}, \ \mathscr{S}_{\mathscr{R}}(k_1) = \{k_1\}, \text{ and } \mathscr{S}_{\mathscr{R}}(k_2) = \{k_2\}.$ However, $k_3 \longrightarrow k_1 = 0 \in \{0, k_3, k_4, 1\}$ and $k_1 \notin \{0, k_3, k_4, 1\}$ imply that $\mathscr{S}_{\mathscr{R}}$ is not \mathbb{D} SQB-A. If we take $\mathscr{R}_1 = \{k_3, k_4, 1\} \subseteq \mathscr{R}$ and we define $\mathscr{S}_{\mathscr{R}_1} = \{y \in \mathscr{X} | \mathscr{O}(x) = \mathscr{O}(y)\} (\forall x \in \mathscr{R}_1)$, then $\mathscr{S}_{\mathscr{R}_1}$ is \mathbb{D} SQB-A.

Definition 14. Assume that $\mathscr{S}_{\mathscr{K}}$ is \mathbb{SQB} - \mathbb{A} over \mathscr{X} with the greatest element 1. If $\mathscr{S}_{\mathscr{K}}(x) = \mathscr{X}(\forall x \in \mathscr{K})$, then $\mathscr{S}_{\mathscr{K}}$ is called whole \mathbb{DSQB} - \mathbb{A} .

Example 8. Suppose \mathcal{X} (i.e., $\mathcal{X} = \{0, k_1, k_2, 1\}$) with partial order $0 < k_1 < k_2 < 1$. Now, we show, by Tables 10 and 11, the binary operations \longrightarrow and \rightsquigarrow , respectively.

Clearly, \mathscr{X} is a CQB-A. We define $\mathscr{S}_{\mathscr{H}}(\forall x \in \mathscr{K})$ (i.e., $\mathscr{K} = \mathscr{X}$) by

$$\mathcal{S}_{\mathcal{K}}(x) = \{ y \in \mathcal{X} | \mathcal{O}(x) = \mathcal{O}(y) \}.$$
(17)

From Tables 10 and 11, we can get on $S_{\mathscr{K}}(x) = \mathscr{X}(\forall x \in \mathscr{K})$. Thus, $S_{\mathscr{K}}$ is a whole \mathbb{DSQB} -A over \mathscr{X} .

Now, we will study homomorphism image of \mathbb{DSQB} -As by the following two theorems.

Theorem 3. Assume that $\psi: \mathcal{X} \longrightarrow \mathcal{Y}$ be a surjective exact morphism of QB-A and \mathcal{X} is a QB-As. If $\mathcal{S}_{\mathcal{X}}$ is a DSQB-A over \mathcal{X} , then $\psi(\mathcal{S}_{\mathcal{X}})$ is alsoDSQB-A over \mathcal{Y} .

Proof. Since $\mathscr{S}_{\mathscr{H}}(x) (x \in \mathscr{H})$ is a deductive system of \mathscr{X} and ψ is surjective, then $\psi(\mathscr{S}_{\mathscr{H}})(x) = \psi(\mathscr{S}_{\mathscr{H}}(x))$ is a deductive system of \mathscr{Y} which implies that $\psi(\mathscr{S}_{\mathscr{H}})$ is a \mathbb{DSQB} -A over \mathscr{X} .

Theorem 4. Assume that $\psi: X \longrightarrow Y$ be a surjective exact morphism of QB-A and $\mathcal{S}_{\mathscr{H}}$ a DSQB-A over \mathscr{X} . Then,

TABLE 10: The binary operation \longrightarrow .

\longrightarrow	0	k_1	k_2	1
0	1	1	1	1
k_1	k_1	1	1	1
k_2	k_1	k_1	1	1
1	Ō	k.	k_{2}	1

TABLE 11: The binary operation ------

-**>	0	k_1	k_2	1
0	1	1	1	1
k_1	k_2	1	1	1
k_2	0	k_1	1	1
1	0	k_1	k_2	1

- If S_K(x) = ker(ψ), for all x ∈ K, then ψ(S_K) is the whole DSQB-A over Y
- (2) If is whole DSQB-A over X, then ψ(S_K) is the whole
 DSQB-A over Y

Proof

- (1) Assume that $\mathscr{S}_{\mathscr{H}}(x) = \ker(\psi)$, where $\ker(\psi) = \{x \in \mathscr{X} | \psi(x) = x \longrightarrow x, \quad \psi(x) = x \dashrightarrow x\}$. Since ψ is surjective, then, from Theorem 3, we have $\psi(\mathscr{S}_{\mathscr{H}})(x) = \psi(\mathscr{S}_{\mathscr{H}}(x)) = \psi(\mathscr{X}) = \psi(\mathscr{X} \in \mathscr{H})$. Thus, $\psi(\mathscr{S}_{\mathscr{H}})$ is the whole $\mathbb{DSQB-A}$ over \mathscr{Y} .
- (2) Clearly, $\mathscr{S}_{\mathscr{H}}(x) = \mathscr{X}$ since $\mathscr{S}_{\mathscr{H}}$ is whole $\mathbb{D}SQ\mathbb{B}$ -A over $\mathscr{X}(x \in \mathscr{H})$. Thus, $\psi(\mathscr{S}_{\mathscr{H}})(x) = \psi(\mathscr{S}_{\mathscr{H}}(x)) = \psi(\mathscr{X}) = \mathscr{Y}(x \in \mathscr{H})$. By Theorem 3, we have $\psi(\mathscr{S}_{\mathscr{H}})$ is the whole $\mathbb{D}SQ\mathbb{B}$ -A over \mathscr{Y}

4. FSQB-As

We give the definition of \mathbb{FSQB} -As; a concrete example is given to illustrate its derive properties. Furthermore, we study the homomorphism image and preimage of \mathbb{FSQB} -As. Now, we first propose the definition of fuzzy quantum B-algebra (briefly, \mathbb{FQB} -A) as indicated below.

Definition 15. We call FQB-A (or a fuzzy set $\hat{\mu}$ in QB-A) if it satisfies ($\forall x, y \in \mathcal{X}, \mathcal{X}$ is QB-A):

$$\hat{\mu}(x \longrightarrow y) \ge \min\{\hat{\mu}(x), \hat{\mu}(y)\},$$

$$\hat{\mu}(x \leadsto y) \ge \min\{\hat{\mu}(x), \hat{\mu}(y)\}.$$
(18)

Definition 16. We call $\hat{\mu}$ is a fuzzy deductive system of \mathcal{X} if it satisfies $(\forall x, y \in \mathcal{X})$:

$$\widehat{\mu}(x \longrightarrow x) \ge \widehat{\mu}(x),
\widehat{\mu}(x \longrightarrow x) \ge \widehat{\mu}(x),
\widehat{\mu}(y) \ge \min\{\widehat{\mu}(x \longrightarrow y), \widehat{\mu}(x)\}.$$
(19)

Definition 17. Assume that $\widehat{\mathscr{S}}_{\mathscr{K}}$ be a FSS over \mathscr{X} . Then,

- If there exists μ̂ ∈ ℋ such that Ŝ_ℋ[μ] is a FQB-A (i.e., fuzzy deductive system) in a QB-A over ℋ, then Ŝ_ℋ is called a Ŝ_ℋ-A (i.e., fuzzy soft deductive system FSDS) which depends on a parameter set μ̂ over ℋ
- (2) If S_ℋ[µ] is a FQB-A (i.e., fuzzy deductive system) of X based on all parameters, then we say that S_ℋ is a FSQB-A (i.e., FSDS) of X

In the following, a concrete example is given to illustrate Definition 17.

Example 9. Suppose that there are five-class cars:

 $X = \{BMW, Audi, Toyota, Jeep, Cadilac\}.$ (20)

Let \oplus and \otimes be two soft machines to characterize two cars, defined by the following manner.

BMW $\oplus x$ = Cadilac, forall $x \in \mathcal{X}$,

$$Audi \oplus y = \begin{cases} Jeep, & y = BMW, \\ Cadilac, & y \in \{Audi, Toyouta, Jeep, Cadilac\}, \\ Toyota \oplus z = \begin{cases} Toyota, & z = BMW, \\ Jeep, & z = Audi, \\ Cadilac, & z \in \{Toyota, Jeep, Cadilac\}, \end{cases}$$
$$Jeep \oplus s = \begin{cases} Toyota, s = BMW, \\ Jeep, s \in \{Audi, Toyoya\}, \\ Cadilac, s \in \{Jeep, Cadilac\}, \end{cases}$$
$$BMW, t = BMW, \\ Audi, t = Audi, \\ Toyota, t = Toyota, \\ Jeep, t = Jeep, \\ Cadilac, t = Cadilac, \end{cases}$$
$$(21)$$

$$BMW \otimes x = Cadilac \text{ forall } x \in \mathcal{X},$$

$$Audi \otimes y = \begin{cases} Jeep, \ y = BMW, \\ Cadilac, \ y \in \{Audi, \text{ Toyouta, Jeep, Cadilac}\}, \\ Cadilac, \ z \in \{BMW, Audi\}, \\ Cadilac, \ z \in \{Toyouta, Jeep, Cadilac\}, \end{cases}$$

$$Jeep \otimes s = \begin{cases} Audi, \ s = BMW, \\ Jeep, \ s \in \{Audi, \text{ Toyouta}\}, \\ Cadilac, \ s \in \{Jeep, Cadilac\}, \end{cases}$$

$$Cadilac, \ s \in \{Jeep, Cadilac\}, \\ BMW, \ t = BMW, \\ Audi, \ t = Audi, \\ Toyota, \ t = Toyota, \\ Jeep, \ t = Jeep, \\ Cadilac, \ t = Cadilac. \end{cases}$$

(22)

Then, $(\mathcal{X}, \oplus, \otimes, \leq)$ is a QB-A. Now, we consider a set of parameters: $\hat{\mu} = (\text{Excellent}, \text{Good}, \text{Moderate}) \in \in \mathcal{K}$. Then, we have the following:

 We define S_𝔅[μ̂] over 𝔅 (i.e., S_𝔅[Excellent], S_𝔅[Good], and S_𝔅[Moderate] are fuzzy sets) by Table 12.

Therefore, we can see that $\hat{\mathscr{S}}_{\mathscr{H}}[\text{Excellent}]$, $\hat{\mathscr{S}}_{\mathscr{H}}[\text{Good}]$, and $\hat{\mathscr{S}}_{\mathscr{H}}[\text{Moderate}]$ are all FSQB-As based on parameters "Excellent," "Good," and "Moderate" over \mathscr{X} . Thus, $\hat{\mathscr{S}}_{\mathscr{H}}$ is a FSQB-A over \mathscr{X} .

(2) We define $\hat{\mathscr{S}}_{\mathscr{H}_1}[\hat{\mu}]$ over \mathscr{X} (i.e., $\hat{\mathscr{S}}_{\mathscr{H}_1}[\text{Excellent}]$, $\hat{\mathscr{S}}_{\mathscr{H}_1}[\text{Good}]$, and $\hat{\mathscr{S}}_{\mathscr{H}_1}[\text{Moderate}]$ are fuzzy sets) by Table 13.

However, $\hat{\mathcal{S}}_{\mathcal{H}_1}[\hat{\mu}]$ is not a FSQB-A based on a parameter "Excellent" over (HTML translation failed), where

$$\begin{split} \widehat{\mathscr{S}}_{\mathscr{H}_1}[\text{Excellent}] &(\text{Toyota} \oplus \text{BMW}) = \widehat{\mathscr{S}}_{\mathscr{H}_1}[\text{Excellent}] \\ &(\text{Toyota}) = 0.1 \not\ge 0.2 = \min\{0.2, 0.4\} = \min\{\widehat{\mathscr{S}}_{\mathscr{H}_1} \\ &[\text{Excellent}] &(\text{Audi}), \widehat{\mathscr{S}}_{\mathscr{H}_1}[\text{Excellent}] &(\text{BMW})\}. \\ &\text{Also,} \\ &\text{we obtain that } \widehat{\mathscr{S}}_{\mathscr{H}_1}[\widehat{\mu}] \text{ is a } \mathbb{FSQB-A} \\ &\text{based on both} \\ &\text{the parameter "Good" and "Moderate" over } \mathscr{X}. \end{split}$$

(3) We define S_{*K*₂}[μ̂] over *X* (i.e., S_{*K*₂}[Excellent] and S_{*K*₂} [Good] are fuzzy sets) by Table 14.

Then, $\hat{\mathscr{S}}_{\mathscr{H}_2}[\hat{\mu}]$ is a FSDS on parameters "Excellent." However, $\hat{\mathscr{S}}_{\mathscr{H}_2}[\mu]$ is not a fuzzy deductive system of \mathscr{X} based on parameter "Good," where $\hat{\mathscr{S}}_{\mathscr{H}_2}[\text{Good}](\text{Toyota}) = 0.3 < 0.5 = \min\{\hat{\mathscr{S}}_{\mathscr{H}_2} [\text{Good}](\text{Jeep} \text{Toyota}), \hat{\mathscr{S}}_{\mathscr{H}_2} [\text{Good}](\text{Jeep})\}.$

 (4) We define δ_{ℋ3}[μ̂] over X (i.e., δ_{ℋ3} [Excellent] and δ_{ℋ2} [Moderate] are fuzzy sets) by Table 15. Then, δ_{ℋ3}[μ̂] is a FSDS of X.

Now, we will present several characterizations of $\mathbb{FSQB}\text{-}\mathsf{As.}$

By Definition 17, if $\widehat{\mathscr{S}}_{\mathscr{X}}$ is a FSQB-A of QB-A over \mathscr{X} based on all parameters, then we say that $\widehat{\mathscr{S}}_{\mathscr{X}}$ is a FSQB-A of \mathscr{X} , that is,

$$\widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](x \longrightarrow y) \ge \min\{\widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](x), \widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](y)\}, \\
\widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](x \longrightarrow y) \ge \min\{\widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](x), \widehat{\mathscr{S}}_{\mathscr{H}}[\widehat{\mu}](y)\}.$$
(23)

Proposition 1. Assume \mathcal{X} be a QB-A. If $\hat{\mathcal{S}}_{\mathcal{H}}$ is FSQB-A over \mathcal{X} , then, for all $t \in [0, 1]$, $(\hat{\mathcal{S}}_{\mathcal{H}})_t \neq \emptyset$ is the subalgebra of \mathcal{X} , in which

$$\left(\widehat{\mathscr{S}}_{\mathscr{H}}\right)_{t} = \left\{ \left(\widehat{\mathscr{S}}_{\mathscr{H}}\left[\widehat{\mu}\right]\right)_{t} | \widehat{\mu} \in \mathscr{H} \right\}.$$
(24)

Proof. Let $(\hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}])_t \neq \emptyset$. Then, $\forall x, y \in (\hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}])_t$; since $\hat{\mathcal{S}}_{\mathscr{H}}$ is a FSQB-A, then $\hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}](x) \ge t$, $\hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}](y) \ge t$. So,

TABLE 12: Fuzzy sets $\widehat{\mathcal{S}}_{\mathscr{K}}[\widehat{\mu}]$ over \mathscr{X} .

$\hat{\mathcal{S}}_{\mathscr{K}}$	BMW	Audi	Toyota	Jeep	Cadilac
Excellent	0.2	0.2	0.5	0.6	0.8
Good	0.1	0.2	0.3	0.5	0.7
Moderate	0.1	0.1	0.4	0.4	0.6

TABLE 13: Fuzzy sets $\widehat{\mathcal{S}}_{\mathscr{K}_1}[\widehat{\mu}]$ over \mathscr{X} .

$\widehat{\mathcal{S}}_{\mathscr{K}_1}$	BMW	Audi	Toyota	Jeep	Cadilac
Excellent	0.4	0.2	0.1	0.6	0.8
Good	0.2	0.2	0.3	0.5	0.7
Moderate	0.1	0.1	0.4	0.5	0.9

TABLE 14: Fuzzy sets $\widehat{\mathscr{S}}_{\mathscr{K}_{2}}[\widehat{\mu}]$ over \mathscr{X} .

$\widehat{\mathcal{S}}_{\mathscr{K}_2}$	BMW	Audi	Toyota	Jeep	Cadilac
Excellent	0.2	0.2	0.2	0.2	0.6
Good	0.2	0.2	0.3	0.5	0.7

TABLE 15: Fuzzy sets $\widehat{\mathcal{S}}_{\mathcal{K}_3}[\widehat{\mu}]$.

$\widehat{\mathcal{S}}_{\mathscr{K}_3}$	BMW	Audi	Toyota	Jeep	Cadilac
Excellent	0.3	0.3	0.3	0.3	0.3
Moderate	0.1	0.1	0.1	0.1	0.7

$$\hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}] (x \longrightarrow y) \ge \min \left\{ \hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}] (x \longrightarrow y), \ \hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}] (x \leadsto y) \right\} \\ \ge \min \left\{ \hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}] (x), \ \hat{\mathcal{S}}_{\mathscr{H}}[\hat{\mu}] (y) \right\} \ge t.$$
(25)

Similarly, we have $\widehat{\mathscr{S}}_{\mathscr{K}}[\widehat{\mu}](x \rightsquigarrow y) \ge t$. Therefore, $x \longrightarrow y, x \rightsquigarrow y \in (\widehat{\mathscr{S}}_{\mathscr{K}}[\widehat{\mu}])_t$. This implies that $(\widehat{\mathscr{S}}_{\mathscr{K}}[\widehat{\mu}])_t$ is the subalgebra of \mathscr{X} .

Analogously, we can get Proposition 2 as follows. \Box

Proposition 2. Assume that $S_{\mathcal{H}_1}$ and $S_{\mathcal{H}_2}$ are two FSQB-A over \mathcal{X} . Then, $S_{\mathcal{H}_1} \cap S_{\mathcal{H}_2}$ and $S_{\mathcal{H}_1} \cup S_{\mathcal{H}_2}$ are FSQB-As over \mathcal{X} .

Definition 18. Let (α, β) be a fuzzy soft map from QB-A over \mathcal{X} to QB-A over \mathcal{Y} . Then,

- If α is an exact morphism from X to Y, then (α, β) is called a FSQB-A exact morphism from X to Y
- (2) If α is an isomorphism from X to Y and β is a bijective from K₁ to K₂, then (α, β) is a called an isomorphism between FSQB-As

Proposition 3. Let \mathcal{X} and \mathcal{Y} be two QB-As. $\mathcal{S}_{\mathcal{H}}$ is a FSQB-A over \mathcal{Y} and (α, β) a FSQB-A exact morphism from \mathcal{X} to \mathcal{Y} ; then, $(\alpha, \beta)^{-1} \mathcal{S}_{\mathcal{H}}$ is FSQB-A over \mathcal{X} .

8

$$\begin{aligned} \alpha^{-1} \left(\mathscr{S}_{\mathscr{H}} \right) [\widehat{\mu}] (x \longrightarrow y) \\ &= \mathscr{S}_{\mathscr{H}} \beta[\widehat{\mu}] (\alpha(x \longrightarrow y)) \\ &= \mathscr{S}_{\mathscr{H}} \beta[\widehat{\mu}] (\alpha(x) \longrightarrow \alpha(y)) \\ &\geq \min \{ \mathscr{S}_{\mathscr{H}} \beta[\widehat{\mu}] \alpha(x), \, \mathscr{S}_{\mathscr{H}} \beta[\widehat{\mu}] \alpha(y) \} \\ &= \min \{ \alpha^{-1} \left(\mathscr{S}_{\mathscr{H}} \right) [\widehat{\mu}] (x), \, \alpha^{-1} \left(\mathscr{S}_{\mathscr{H}} \right) [\widehat{\mu}] (x) \}. \end{aligned}$$
(26)

Consequently, $(\alpha, \beta)^{-1} \mathscr{S}_{\mathscr{X}}$ is a FSQB-A over \mathscr{X} . Similarly, we can get Proposition 4 as follows.

Proposition 4. Let \mathcal{X} and \mathcal{Y} be two QB-As. $\mathcal{S}_{\mathcal{X}}$ is a FSQB-As over \mathcal{X} and (α, β) a FSQB-As isomorphism from \mathcal{X} to \mathcal{Y} ; then, $(\alpha, \beta)\mathcal{S}_{\mathcal{H}}$ is the FSQB-As over \mathcal{Y} .

5. Conclusions

In this paper, we introduce the concept of \mathbb{QB} -As, and some examples are given to illustrate this definition. Also, we investigate the union and intersection operations between two \mathbb{QB} -As and give some conditions for the operation holds. With the help of the definition of \mathbb{SQB} -As, we define soft deductive systems of \mathbb{SQB} -As and then investigate the relation between them. As a further step, we define \mathbb{DSQB} -As and investigate the homomorphism image of \mathbb{DSQB} -As. Moreover, we define \mathbb{FSQB} -As. Finally, a concrete example is given to illustrate its derive properties; besides, homomorphism image and preimage of \mathbb{FSQB} -As are discussed.

As a future work, it makes sense to apply \mathbb{QB} -As to medical diagnosis (for example, [25, 26]) in practice. Furthermore, it would be interesting if we study hybrid soft lattice-ordered quantum B-algebras.

Data Availability

No data were used to support this study.

Conflicts of Interest

The authors declare that they have no conflicts of interest.

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