Research Article

Flocking for Multiagent Systems with Partial Information Exchange Based on Inclusion Principle

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This paper investigates the flocking problem of multi-agents with partial information exchange, which means that only part, but not all, of the agents are informed of the group objective. A distributed flocking model based on the inclusion principle is provided to simplify the design and analysis of multi-agent systems. Furthermore, to reduce the communication energy consumption, an improved flocking algorithm based on the model is proposed to achieve stable flocking for all the agents. The stability of the multi-agent system is then established, with the help of the Lyapunov stability theorem and LaSalle’s invariance principle. Especially, considering the individual heterogeneity in both nature and engineering applications, we also investigate the flocking problem of multi-agents with different sensing radiuses and equilibrium distances. Finally, two kinds of simulation results are presented to demonstrate the validity of the proposed results. This work provides an insight not only into the properties of the different species of individual flocking, but also into the theoretical framework for the engineering design of multi-agent systems considering individual heterogeneity.

1. Introduction

Flocking, a common phenomenon in nature, can eventually achieve a group objective through local communication among the neighbouring agents. Examples of these agents include flocks of birds, schools of fish, groups of ants, and crowds of people. For many decades, flocking problems have received a great deal of attention from researchers in biology, social science, physics, control science, computer science, and so on [1–10], due to the emergence of swarming for a group of agents with local interactions. This phenomenon of swarming finds a broad range of applications in swarm robots [11], mobile sensor networks [12], and unmanned aircraft systems (UAS) [13–15].

In 1986, Reynolds produced a computer model by simulating flocks of birds, which consists of three basic rules of cohesion, separation, and alignment [9]. Soon after this, Vicsek et al. first studied the collective behavior of bird swarms from a theoretical perspective [16–19]. For instance, Vicsek et al. [16] proposed a self-driven particle model in 1995, which mainly concentrated on the emergence of flocking, whereas Toner and Tu [17] preferred the continuum mechanics method. From the particle-based model proposed by Vicsek et al. [16], Helbing et al. [18] conducted a series of experiments to explore the mechanism of escape panic. In 2003, Jadbabaie et al. [19] studied the linear Vicsek model without noise and were the first to provide a theoretical explanation and mathematical proof of this model. In 2006, inspired by the Reynolds rules, Olfati-Saber [7] designed two typical flocking algorithms: flocking in an obstacle-free environment and flocking with obstacle avoidance. Moreover, he performed 2D flocking in the presence and absence of obstacles and 3D flocking in free space. In his algorithms, it was assumed that all the agents can obtain the information of the group objective (or virtual leader). Such an assumption, however, requires significant energy consumption during the flocking process. This makes it almost impossible to be implemented in engineering applications, such as multitarget consensus circle pursuit for multiagent systems via a distributed multiflocking method. To overcome this limitation, Su et al. [20] extended the typical flocking algorithm, considering that only partial agents are informed of the group objective, and then, numerical simulations indicated that most agents will asymptotically track the group objective,
even if only a fraction of the group are informed agents. However, there are some uninformed agents that still fail to track the group objective during the flocking process. Therefore, it is necessary to improve the typical flocking algorithm further in this work.

The aforementioned works have greatly promoted the research process of flocking control for multi-agent systems and established an important foundation for future researches.

Recently, some studies were conducted that are closely related to the research topic of this paper. For example, Luo et al. [21] investigated the multi-target tracking problem of multi-agent systems, which, however, is still a deadlock problem. In order to overcome this problem, Pei et al. [22] proposed a local multi-flocking algorithm, which successfully resulted in the multi-target consensus pursuit for multi-agent systems. Xi et al. [23] provided a theoretical framework for the analysis of the output consensus problem of high-order linear time-invariant multi-agent systems, during the flocking process. Cao and Ren [24] presented a distributed consensus tracking algorithm for first-order kinematics and a distributed swarm tracking algorithm for second-order dynamics, which mainly concentrated on solving the distributed coordinated tracking problem. Based on the continuous-time information-weighted Kalman consensus filter (IWKCF), Luo and Li [25] proposed a distributed topology optimization scheme to decrease the communication complexity of flocking for multi-agent systems, in which each agent produces a local optimally rigid graph with its neighbouring agents. However, most of the above studies have merely focused on the homogeneous multi-agent system. Considering the individual characteristics, multi-agent systems can be divided into two types: homogeneous and heterogeneous [26]. Individual heterogeneity exists widely in both nature and engineering applications, such as multi-UAS in military confrontations, different growth backgrounds, social distancing, and economic strength in human society. Therefore, it is of great theoretical and practical significance for the research of the flocking problem of multi-agents with different equilibrium distances and sensing radiiues. Note that the equilibrium distance is proportional to the sensing radius (according to the works of [7]). For the sake of brevity, we will discuss only the different equilibrium distances later.

As further research is being conducted on multi-agent systems, we can determine that the motion consensus of multi-agents is similar to the synchronization problem of complex networks in essence, which is the solution to achieve the consistency of the state of all agents [27–29]. Furthermore, the inclusion principle and its extended pair-wise decomposition can take full advantage of the interconnection relationship between subsystems, which have been widely applied in fields such as multi-area interconnected power systems [30–32], formation control of unmanned aerial vehicles [33], Petri nets [34], etc. Hence, the inclusion principle and its extended pair-wise decomposition can be chosen as a proper tool to simplify the design and analysis of flocking for multi-agent systems.

In this paper, we first divide a complex swarm system into multiple pair-wise subsystems based on the inclusion principle. Then, from the Lennard-Jones potential function [35] and self-organization process [36], an improved flocking algorithm is designed for each subsystem. Thereby, flocking of multi-agents with partial information exchange is achieved via the coordinated control of the pair-wise subsystems, which means that only a part, but not all, of the agents can be informed of the group objective. Compared with the typical flocking algorithm [7] and the existing distributed topology optimization scheme [25], this work, to some extent, has decreased the calculated amount and greatly reduces the communication energy consumption during the flocking process. Furthermore, considering the individual heterogeneity in both nature and engineering applications, the multi-agent system can be divided into three different species, according to the size of the equilibrium distance. Particularly, one of the aspects we are very interested in is the collision avoidance between different species of multi-agents during the flocking process. In order to explore the internal mechanism of this interesting problem, two kinds of simulations are designed to investigate the flocking of multi-agents with different equilibrium distances, which can provide an insight into the properties of the different species of individual flocking.

In brief, the main contributions of this paper are as follows: (1) A distributed flocking model based on the inclusion principle is proposed to simplify the design and analysis of flocking for multi-agent systems. (2) Based on the Lennard-Jones potential function and self-organization process, an improved flocking algorithm is designed to achieve flocking of multi-agents with partial information exchange. (3) Considering the individual heterogeneity in both nature and engineering applications, the flocking problem of multi-agents with different equilibrium distances is also investigated.

An outline of the paper is organized as follows. Section 2 introduces some preliminary knowledge. A distributed flocking model based on the inclusion principle is presented in Section 3. Section 4 proposes an improved flocking algorithm. In Section 5, the stability analysis of multi-agent systems is proven with the help of the Lyapunov stability theorem and LaSalle’s invariance principle. Several simulation results are provided in Section 6. Finally, the conclusions are presented in Section 7.

2. Preliminaries

The flocking algorithm design process proposed in this paper is based on some preliminary knowledge about graph theory, especially, the inclusion principle and its extended pairwise decomposition.

2.1. Graph Theory. For the description of multi-agent systems by graph theory in this paper, we have some definitions as follows [37].

Definition 1. Suppose that a network topology consists of $N$ vertexes. The network topology can be described by an undirected graph $G=(V,E)$, where $V=\{1,2,\ldots,N\}$ is a
vertex set and $E \subseteq \{(i, j): i, j \in v, j \neq i\}$ denotes the edge set
with vertexes of junctions.

**Definition 2.** The adjacency matrix $A_a = [a_{ij}]$ of the undirected
graph $G$ is defined as

$$a_{ij} = \begin{cases} 1, & \text{if } (i, j) \in v, \\ 0, & \text{others}, \end{cases} \quad \text{(1)}$$

where $A_a$ is symmetric, i.e. $A_a^T = A_a$.

**Definition 3.** The degree matrix of undirected graph $G$ is a
diagonal matrix $D = D(A_a)$ with diagonal elements $\sum_{i=1}^{n} a_{ii}$
that are row sums of $A_a$. The Laplacian matrix $L = [l_{ij}]$ is an
$n \times n$ matrix, which is defined as

$$L = D - A_a. \quad \text{(2)}$$

2.2. **Inclusion Principle.** As mentioned in Section 1, the
inclusion principle is a theoretical basis for pair-wise decom-
position and decentralized control of complex inter-
connected systems [32]. In order to control multi-agent
systems coordinately, we introduce some definitions and
theorems as follows [30].

Consider a continuous-time linear time-invariant system
$S$ and its expanded system $\tilde{S}$ described by

$$\begin{align*}
S: & \quad \dot{x} = Ax + Bu, \quad y = Cx, \\
\tilde{S}: & \quad \tilde{x} = \tilde{A}\tilde{x} + \tilde{B}\tilde{u}, \quad \tilde{y} = \tilde{C}\tilde{x},
\end{align*} \quad \text{(3)}$$

where $x \in \mathbb{R}^n$, $u \in \mathbb{R}^m$, and $y \in \mathbb{R}^l$ are the state, input, and
output vectors of the system $S$, respectively. $\tilde{x}(t) \in \mathbb{R}^n$,
$\tilde{u}(t) \in \mathbb{R}^m$, and $\tilde{y}(t) \in \mathbb{R}^l$ are the state, input, and output
vectors of the expanded system $\tilde{S}$, respectively. The matrices
$A, B, C$ and $\tilde{A}, \tilde{B}, \tilde{C}$ have appropriate dimensions. It is
supposed that $n \leq \tilde{n}, m \leq \tilde{m}$, and $l \leq \tilde{l}$.

**Definition 4.** The expanded system $\tilde{S}$ includes the system $S$, namely, $\tilde{S}$.
If there is a quadruplet of full-rank matrices $\{V, U, R, S\}$ satisfying
$UV = I_n$, such that for any initial condition $x_0 \in \mathbb{R}^n$ and any input $u \in \mathbb{R}^m$, when $\tilde{x}_0 = Vx_0$,
$\tilde{u} = Ru$, there is $\tilde{x}(t; t_0, x_0, u) = U\tilde{x}(t; t_0, \tilde{x}_0, \tilde{u})$ and
$y[x(t)] = S\tilde{y}[\tilde{x}(t)]$ for all $t \geq t_0$.

**Theorem 1.** If the system $S$ is a restriction of the expanded
system $\tilde{S}$, namely $S \subset \mathbb{R}^n\tilde{S}$, then there exists a triplet of full-rank matrices $\{V, R, T\}$ such that

$$\begin{align*}
\tilde{A}V = VA, \quad \tilde{B}R = VB, \quad \tilde{C}V = TC, \\
or \quad MAV = 0, MBR = 0, MCV = 0,
\end{align*} \quad \text{(4)}$$

where $V, R, T, M_A, M_B$, and $M_C$ will be illustrated in
Theorem 2.

**Proof.** The detailed proof is carried out in [30].

**Theorem 2.** If the system $S$ is an aggregation of the expanded
system $\tilde{S}$, namely, $S \subset \mathbb{R}^n\tilde{S}$, then there exists a triplet of full-rank matrices $\{U, Q, S\}$ such that

$$U\tilde{A} = AU, U\tilde{B} = BQ, S\tilde{C} = CU,$$

or $U\tilde{M}_A = 0, U\tilde{M}_B\tilde{R} = 0, S\tilde{M}_C\tilde{V} = 0$. \quad \text{(5)}$$

As mentioned above, the relationship between the sys-
tern $S$ and expanded system $\tilde{S}$ is described by

$$\tilde{A} = VAU + M_A, \quad \tilde{B} = VBQ + M_B, \quad \tilde{C} = TCU + M_C,$$

where $V, U, R, Q, T, \text{ and } S$ are full-rank transformation
matrices with dimensions of $n \times n, \tilde{n} \times n, \tilde{m} \times \tilde{m}, \text{ and } \tilde{l} \times \tilde{l}, \tilde{l} \times l, l \times l$, respectively, and satisfy $UV = I_n, \text{QR} = I_n$, and
$ST = I, \text{MA}, \text{MB}, \text{and } MC$ are complementary matrices with
dimensions of $n \times \tilde{n}, \tilde{m} \times \tilde{m}$, and $\tilde{l} \times \tilde{l}$, respectively.

**Proof.** The detailed proof is carried out in [30].

3. **Distributed Flocking Model Based on Inclusion Principle**

During the flocking process of multi-agents, the dynamics of
the multi-agent system should be modeled mathematically
to facilitate further investigation. This does not depend on
where the agent is natural biology (e.g., bees and ants) or a piece of engineering equipment (e.g., robots and UAVs). In this section, we assume that all agents have the same mass and size. Consider a group of $N$ mobile agents (or subsystems) with second-order linear dynamics, which is described by

$$
S_i:\begin{bmatrix}
\dot{p}_i \\
\dot{v}_i
\end{bmatrix} = 
\begin{bmatrix}
0 & 1 \\
0 & 0
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix} + 
\begin{bmatrix}
0 \\
1
\end{bmatrix} u_i,
$$

$$
\begin{bmatrix}
y_{pi} \\
y_{vi}
\end{bmatrix} = 
\begin{bmatrix}
1 & 0 \\
0 & 1
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix}, \quad i = 1, 2, \ldots, N,
$$

where $p_i, v_i \in \mathbb{R}^m$ (e.g., $m = 2$ or 3, when a flocking in a 2D or 3D space) denote the position and velocity vector of the agent $i$, respectively. $u_i \in \mathbb{R}^m$ is the control input of agent $i$.

Suppose that each agent has a limited sensing radius (or interaction range). Let $r > 0$ denote the sensing radius, and then, the neighbor set of the agent $i$ is defined as

$$
N_i = \left\{ j \in \nu : \|p_j - p_i\| < r \right\},
$$

where $\|\cdot\|$ is the Euclidean norm in $\mathbb{R}^m$; each agent can acquire the information of other mobile agents within its neighbor set. Moreover, one dynamic (or static) $\gamma -$ agent represents the group objective (or virtual leader) of multi-agent systems, which drives all agents to track the group objective with the following model:

$$
\begin{align*}
\dot{p}_\gamma &= v_\gamma, \\
\dot{v}_\gamma &= f_\gamma(p_{\gamma}, v_{\gamma}),
\end{align*}
$$

where $p_\gamma, v_\gamma, f_\gamma \in \mathbb{R}^m$ denote the position, velocity, and control input vector of the group objective, respectively. The initial state vector pairs are set to $(p_\gamma(0), v_\gamma(0)) = (p_d, v_d)$. A static $\gamma -$ agent means that its state is fixed, that is, $(p_\gamma, v_\gamma) = (p_d, v_d)$ for all the way. In Section 6, we assume that the group objective always moves at a constant velocity to the fixed direction. Thereby, dynamic equation (12) of the group objective is simplified as $\dot{p}_\gamma = v_d$, $\dot{v}_\gamma = 0$.

Considering that $p_i, v_i$ are measurable, the linear state-space model of the multi-agent system can be described by

$$
S: \dot{x} = Ax + Bu, y = Cx,
$$

where $x = (p_1, v_1, p_2, v_2, \ldots, p_N, v_N)^T$, $u = (u_1, u_2, \ldots, u_N)^T$, and $y = (y_{p1}, y_{v1}, \ldots, y_{pN}, y_{vN})$ are the state, input, and output vectors of the system $S$, respectively. Meanwhile,

$$
A = \begin{bmatrix}
A_{11} & \cdots & A_{1N} \\
\vdots & \ddots & \vdots \\
A_{N1} & \cdots & A_{NN}
\end{bmatrix}, \quad B = \text{blockdiag}[B_{11}, \ldots, B_{NN}],
$$

$$
C = \text{blockdiag}[C_{11}, \ldots, C_{NN}],
$$

in which $A_{ij}$, $B_{ii}$, and $C_{ii}$ are submatrices with dimensions of $n_i \times n_i$, $n_i \times m_i$, and $l_i \times n_i$, respectively, and $i = 1, 2, \ldots, N, j = 1, 2, \ldots, N$. According to equations (10), (11), and (13), pairwise subsystems (or neighbouring agents) can be identified by

$$
S_{ij}:
\begin{bmatrix}
\dot{p}_i \\
\dot{v}_i
\end{bmatrix} = 
\begin{bmatrix}
1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix} + 
\begin{bmatrix}
1 \\
0
\end{bmatrix} u_i,
$$

$$
\begin{bmatrix}
y_{pi} \\
y_{vi}
\end{bmatrix} = 
\begin{bmatrix}
1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix},
$$

where $\|d_{ij}\| = \|p_j - p_i\|$.

To illustrate the applicability of the pair-wise decomposition [30] and the improved flocking algorithm (see Section 4), we consider an interconnected multi-agent system (the number of agents $N \geq 3$) as an example. This system with its equilibrium state is shown in Figure 1(a). The vertices denote the subsystems (or agents), represented by blue spheres. The edges denote the interconnections between the neighbouring subsystems, represented by solid lines. The dashed ovals represent the pair-wise subsystems that can be decomposed, as shown in Figure 1(b). The basic interconnection coefficient $e_{ij}$ is used to represent the subsystems’ connection states, if the subsystem $S_i$ is connected with the subsystem $S_j$, then $e_{ij} = 1$; otherwise, $e_{ij} = 0$. From Figure 1(c), the dashed ovals, respectively, divide the pair-wise subsystem in the system $S$, among which the intersections, namely, intersected subsystems, are the overlapped parts between subsystems, that are also called overlapped subsystems.

Based on the inclusion principle and its extended pair-wise decomposition [32], we consider the coordinated control of each pair-wise subsystem as the basic control of the multi-agent system. Then, the flocking control for the whole system is realized through the parallel coordinated control of these pair-wise subsystems, thereby simplifying the analysis and design of such complex systems. For example, when $N = 3$, the system (shown in Figure 1(b)) is represented by equation (13) with

$$
A = \begin{bmatrix}
A_{11} & A_{12} & A_{13} \\
A_{21} & A_{22} & A_{23} \\
A_{31} & A_{32} & A_{33}
\end{bmatrix},
$$

$$
B = \text{blockdiag}[B_{11}, \ldots, B_{NN}],
$$

$$
C = \text{blockdiag}[C_{11}, \ldots, C_{NN}],
$$

in which $A_{ij}$, $B_{ii}$, and $C_{ii}$ are submatrices with dimensions of $n_i \times n_i$, $n_i \times m_i$, and $l_i \times n_i$, respectively, and $i = 1, 2, \ldots, N, j = 1, 2, \ldots, N$. According to equations (10), (11), and (13), pairwise subsystems (or neighbouring agents) can be identified by

$$
S_{ij}:
\begin{bmatrix}
\dot{p}_i \\
\dot{v}_i
\end{bmatrix} = 
\begin{bmatrix}
0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix} + 
\begin{bmatrix}
1 \\
0
\end{bmatrix} u_i,
$$

$$
\begin{bmatrix}
y_{pi} \\
y_{vi}
\end{bmatrix} = 
\begin{bmatrix}
0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0
\end{bmatrix} 
\begin{bmatrix}
p_i \\
v_i
\end{bmatrix},
$$

where $\|d_{ij}\| = \|p_j - p_i\|$.

The topology can be considered as a composite of the loop structure between the subsystems $S_1$, $S_2$, and $S_3$. From Theorems 1 and 2, we apply the following expansion transformations:
Figure 1: Schematic diagram of stable flocking for multiagent systems. (a) The whole system. (b) The system with \( N = 3 \). (c) The expanded pairwise subsystems.

\[
V = \text{blockdiag}\left( \begin{bmatrix} I_{11}, I_{11} \end{bmatrix}^T, \begin{bmatrix} I_{22}, I_{22} \end{bmatrix}^T, \begin{bmatrix} I_{33}, I_{33} \end{bmatrix}^T \right),
\]

\[
U = 0.5 \text{ blockdiag}\left( \begin{bmatrix} I_{11}, I_{11} \end{bmatrix}, \begin{bmatrix} I_{22}, I_{22} \end{bmatrix}, \begin{bmatrix} I_{33}, I_{33} \end{bmatrix} \right),
\]

\[
M_A = 0.5 \times \begin{bmatrix}
A_{11} & -A_{11} & A_{12} & -A_{12} & -A_{13} & A_{13} \\
-A_{11} & A_{11} & -A_{12} & A_{12} & -A_{13} & A_{13} \\
-A_{21} & A_{21} & A_{22} & -A_{22} & 0 & 0 \\
-A_{21} & A_{21} & -A_{22} & A_{22} & 0 & 0 \\
A_{31} & -A_{31} & 0 & 0 & A_{33} & -A_{33} \\
A_{31} & -A_{31} & 0 & 0 & -A_{33} & A_{33}
\end{bmatrix}
\]

(16)

Hence,

\[
\bar{A} = \begin{bmatrix}
A_{11} & 0 & A_{12} & 0 & 0 & A_{13} \\
0 & A_{11} & A_{12} & 0 & 0 & A_{13} \\
0 & A_{21} & A_{22} & 0 & 0 & 0 \\
0 & A_{21} & 0 & A_{22} & 0 & 0 \\
A_{31} & 0 & 0 & 0 & A_{33} & 0 \\
A_{31} & 0 & 0 & 0 & 0 & A_{33}
\end{bmatrix}
\]

(17)

Then, according to Theorem 3, we use the permutation matrix,

\[
P_A = \prod_{k=1}^{5} P_{A(k+1)} = P_{A12}P_{A23}P_{A34}P_{A45}P_{A56},
\]

resulting in

\[
\bar{A}_p = P_A^{-1} \bar{A} P_A = \begin{bmatrix}
A_{11} & A_{12} & 0 & 0 & A_{13} & 0 \\
A_{21} & A_{22} & 0 & 0 & 0 & 0 \\
0 & 0 & A_{22} & A_{23} & 0 & 0 \\
0 & 0 & A_{32} & A_{33} & 0 & 0 \\
0 & 0 & 0 & 0 & A_{33} & A_{31} \\
0 & A_{12} & 0 & 0 & A_{13} & A_{11}
\end{bmatrix}
\]

where we have the desired pair-wise control structure as shown in Figure 1(c), that is, each pair-wise subsystem satisfies equation (14).

According to equation (11), the equilibrium structure (or desired geometry) of flocking requires that the distance among the neighboring agents is constant, thereby satisfying the following constraints:

\[
\left\| p_j - p_i \right\| = d_e, \quad (20)
\]

where \( i = 1, 2, \ldots, N \) and \( j \in \mathcal{N}_i(t) \) as time \( t \to \infty \). The equilibrium distance \( d_e \) is a positive constant, usually \( d_e < r \). Particularly, the equilibrium distance plays an important role in the desired formation (e.g., a geometric model of flocks).

4. Flocking Control Algorithm

From the above distributed flocking model, an improved flocking algorithm is designed to achieve stable flocking of multi-agents. Since 2006, the typical flocking algorithm, proposed by Olfati-Saber [7], has been extensively applied in fields such as mobile sensor networks and unmanned aircraft systems (UAS), etc. [12–15]. In [7], the control protocol (or input) \( u_i \) for agent \( i \) consists of the following three components:

\[
u_i = u_i^g + u_i^d + u_i^c,
\]

(21)

where \( u_i^g \) and \( u_i^d \) are the gradient-based term and velocity consensus term of the control protocol \( u_i \), respectively. \( u_i^c \) is the navigational feedback to track the group objective, which is described by

\[
u_i^c = c_1 \left( p_i - p_\gamma \right) + c_2 \left( v_i - v_\gamma \right), \quad c_1, c_2 > 0.
\]

(22)

In this section, a novel distributed flocking algorithm is considered based on the Lennard-Jones potential function [35] and the self-organization process [36]. In contrast with the typical flocking algorithm [7], we simplify the interaction
protocols between the neighbouring agents and develop a complex flocking of multi-agents, via a simpler pair-wise action function that only considers repulsion and attraction between the neighbouring agents (or the pair-wise subsystems). More specifically, according to (14) and (20), the control protocol \( u_i \) for the agent \( i \) (or the subsystem \( S_i \)) is given by the following function:

\[
u_i = \sum_{j \in \mathcal{N}_i} f_a\left(\|p_j - p_i\|\right)e_{ij} + \sum_{j \in \mathcal{N}_i} a_{ij}(v_j - v_i) + c_1(p_j - p_i) + c_2(v_j - v_i), \quad i \in \mathcal{V},\]

where \( e_{ij} = \left(\left(p_j - p_i\right)/\left(\sqrt{1 + \epsilon\|p_j - p_i\|^2}\right)\right)\), \( \epsilon \in (0, 1) \) is a value along the line connecting \( p_i \) to \( p_j \), and \( a_{ij} \) is the element of the adjacency matrix \( A_a \) as in Definition 2 and satisfies (1). In addition, compared with the traditional artificial potential function [7, 20–25], a simpler pair-wise action function that only considers repulsion and attraction among the neighbouring agents (or the pair-wise subsystems) is presented to reflect the interaction protocols. In this way, to some extent, the calculation difficulties can be decreased during the flocking process. Thus, the pair-wise action function \( f_a(x) \) is defined as

\[
f_a(x) = \begin{cases} k_1 \left(1 - \frac{x}{d_e}\right), & x \leq d_e, \\ 0, & x > r, \\ k_2 (d_e - x) e^{((x - d_e)^2)/c_2}, & d_e < x \leq r, \end{cases}\]

where \( k_1 \) and \( k_2 \) are defined as the coefficients of the interconnection force, respectively, and especially, the pair-wise action function \( f_{a_1}(x) \) can be drawn as in Figure 2 when \( d_e = 5 \). Figure 2 shows that the action force between the neighbouring agents is an attraction for a long range, but repulsion for a short range. Moreover, the pair-wise potential function \( F_a(x) \) is described by

\[
F_a(x) = \int_{d_e}^{x} f_a(s)ds. \tag{25}
\]

Note that the pair-wise potential function \( F_a(\|p_j - p_i\|) \) is a differential, nonnegative function satisfying the following conditions: (1) \( F_a \) reaches its unique minimum when \( \|p_j - p_i\| = d_e \); (2) \( F_a \rightarrow \infty \) as \( \|p_j - p_i\| \rightarrow 0 \); (3) \( F_a \) remains constant if \( \|p_j - p_i\| > r \).

In the flocking control protocol (23), it is assumed that all agents can obtain the information of the group objective, which is almost impossible and requires considerable energy consumption in practice. For the purpose of reducing the communication energy consumption, in this section, we assume that only a part, but not all, of the agents are informed of the group objective. Consequently, the control protocol (23) is modified as

\[
u_i = \sum_{j \in \mathcal{N}_i} f_a\left(\|p_j - p_i\|\right)e_{ij} + \sum_{j \in \mathcal{N}_i} a_{ij}(v_j - v_i) + h_i\left[c_1(p_j - p_i) + c_2(v_j - v_i)\right], \quad i \in \mathcal{V}, \tag{26}
\]

where the control indicator \( h_i \) is utilized to achieve the flocking of multi-agents with partial information exchange. In other words, if the agent \( i \) is informed of the group objective, \( h_i = 1 \); otherwise, \( h_i = 0 \). For example, in Section 6, we assume that there are \( M \) agents informed agents, which means that \( h_i = 1 \) for \( i = 1, 2, \ldots, M \), but \( h_i = 0 \) for \( i = M + 1, M + 2, \ldots, N \). Note that different from the work of Su et al. [20], we are more focused on selecting an appropriate ratio of the informed agents, which can not only reduce the communication energy consumption, but also enable all the agents (either informed or uninformed) to track the group objective during the flocking process.

5. Stability Analysis

As discussed above, an improved flocking algorithm based on the proposed distributed flocking model is proposed. In this section, we first introduce a theorem to establish a stable flocking of multi-agents and then prove this theorem via the Lyapunov stability theorem and LaSalle's invariance principle.

**Theorem 4.** Consider a group of \( N \) mobile agents (or subsystems) with dynamics (10) applying the flocking control protocol (26). Assume that the initial positions and velocities of all agents are chosen at random with the Gaussian distribution. Then, the following statements hold:

(i) No collisions occur between neighbouring agents

(ii) The velocity of all agents asymptotically become consistent

(iii) Flocking of all agents is formed asymptotically

**Proof.** Define \( \bar{p}_i = p_i - p_p, \bar{v}_i = v_i - v_p \) as the position error vector and velocity error vector, respectively. Consequently, the error dynamic of the agent \( i \) is described by
Moreover, let $p_{ij} = p_i - p_j$ and $\tilde{p}_{ij} = \tilde{p}_i - \tilde{p}_j$; clearly, $\tilde{p}_{ij} = p_{ij}$. Hence, the collective potential function $V_i$ in [7] is modified as
\begin{equation}
\tilde{V}_i(p_{ij}) = \sum_{j \in V(i)} F_a(\|p_{ij}\|) = \sum_{j \in N_i,j \neq i} F_a(r) + \sum_{j \in N_i} F_a(\|p_{ij}\|).
\end{equation}

Similarly, flocking control protocol (26) of the agent $i$ can be rewritten as
\begin{equation}
u_i = \sum_{j \in N_i} f_a(\|\tilde{p}_{ij}\|) c_{ij} + \sum_{j \in N_i} a_{ij}(\tilde{v}_j - \tilde{v}_i) - h_i[c_1 \tilde{p}_i + c_2 \tilde{v}_i],
\end{equation}
where $c_1, c_2 > 0$.

We choose an energy-like Lyapunov function as follows:
\begin{equation}Q(\tilde{p}, \tilde{v}) = \frac{1}{2} \sum_{i=1}^{N} (U_i(\tilde{p}) + \tilde{v}_i^T \tilde{v}_i),\end{equation}
where
\begin{equation}
U_i(\tilde{p}) = \sum_{j=1,j \neq i}^{N} F_a(\|\tilde{p}_{ij}\|) + h_i c_1 \tilde{v}_i^T \tilde{v}_i = \tilde{V}_i(p_{ij}) + h_i c_1 \tilde{v}_i^T \tilde{v}_i, \end{equation}
and $\tilde{p} = [p_1, p_2, \ldots, p_N]^T$, $\tilde{v} = [v_1, v_2, \ldots, v_N]^T$.

Thanks to the symmetry of $F_a(x)$ and the adjacent matrix $A_s$ in Definition 2, it follows that
\begin{equation}\frac{\partial F_a}{\partial p_{ij}} = \frac{\partial F_a}{\partial p_{ij}} = \frac{\partial F_a}{\partial p_{ij}}.\end{equation}

Then,
\begin{equation}\frac{1}{2} \sum_{i=1}^{N} \tilde{U}_i(\tilde{p}) = \sum_{i=1}^{N} (\tilde{V}_i \tilde{v}_i) \tilde{V}_i(p_{ij}) + h_i c_1 \tilde{v}_i^T \tilde{p}_i. \end{equation}

Consequently,
\begin{equation}\dot{Q}(\tilde{p}, \tilde{v}) = \sum_{i=1}^{N} \tilde{U}_i(\tilde{p}) + \sum_{i=1}^{N} \tilde{V}_i \tilde{v}_i = \tilde{V}_i(p_{ij}) + h_i c_1 \tilde{v}_i^T \tilde{v}_i \end{equation}
\begin{equation}\geq \sum_{i=1}^{N} \tilde{V}_i(p_{ij}) \geq Q_0,\end{equation}
where $\otimes$ is the Kronecker product notation, $L(t)$ denotes the Laplacian matrix in Definition 3, $H(t) = \text{diag}(h_1, h_2, \ldots, h_N)$, and $I_n$ is the identity matrix with $n$ dimensions.

Note that the graph of the adjacent matrix $A_s$ is connected, as shown in Figure 1(a). At this moment, $L(t)$ and $H(t)$ are both positive semidefinite matrices, and especially, it is clear that $L(t) + c_2 H(t)$ is positive semidefinite as well. Hence, $Q(\tilde{p}, \tilde{v}) \leq 0$, which indicates that $Q(\tilde{p}, \tilde{v}) = Q(t)$ is a nonincreasing function over time $t$. Therefore, $Q(t) \leq Q_0$ for all $t \geq 0$, and $Q_0$ is the initial value of $Q(t)$. From equations (30) and (31), we conclude that $c_1 \tilde{p}_i^T \tilde{p}_i \leq 2Q_0$ for any agent $i$, which guarantees the flocking of multi-agents. Therefore, part (iii) is proved.

Since $Q(t) > 0$ and $\dot{Q}(t) \leq 0$, we assume that $\Omega = \{ (\tilde{p}^T, \tilde{v}^T) \mid (\tilde{p}^T, \tilde{v}^T) \leq Q_0 \}$ is an invariant set. From LaSalle’s invariance principle, the trajectories of all agents starting from $\Omega$ will converge to the largest set $\Psi = \{ (\tilde{p}^T, \tilde{v}^T) \mid (\tilde{p}^T, \tilde{v}^T) = 0 \}$. Suppose that the number of connected subgraphs is $m(t)$. For any $t > 0$, there is an orthogonal permutation matrix $P(t)$ such that $L(t) + c_2 H(t)$ can be transformed into the following form:
\begin{equation}\tilde{L}(t) + c_2 \tilde{H}(t) = P(t)(L(t) + c_2 H(t)) P^T(t) \end{equation}
\begin{equation}= \text{diag}(L_1(t) + c_2 H_1(t), \ldots, L_k(t) + c_2 H_k(t)), \end{equation}
where $L_k(t)$ and $H_k(t)$ are the Laplacian matrix and the diagonal matrix corresponding to the $k$th connected subgraph, respectively, $k = 1, 2, \ldots, m(t)$. The velocity error vector can also be rewritten as $\tilde{v}^p = (P(t) \otimes I_n) \tilde{v}$, consequently, we have
\begin{equation}\tilde{v}^p [L(t) + c_2 H(t) \otimes I_n] \tilde{v} = (\tilde{v}^p)^T [(\tilde{L}(t) + c_2 \tilde{H}(t) \otimes I_n)] \tilde{v}^p. \end{equation}

As mentioned above, $\tilde{L}(t) + c_2 \tilde{H}(t)$ is positive semidefinite as well. Given (34) and (36), we conclude that $\dot{Q}(t) = 0$ if and only if $(\tilde{v}^p)^T (\tilde{L}(t) \otimes I_n) \tilde{v}^p = 0$ and $(\tilde{v}^p)^T (c_2 \tilde{H}(t) \otimes I_n) \tilde{v}^p = 0$, which is equivalent to $v_1 \equiv \cdots \equiv v_k \equiv \cdots \equiv v_N$. Therefore, part (ii) is proven.

Finally, we prove part (i) by contradiction. Suppose that at least two agents are colliding during the flocking process; then, we acquire
\begin{equation}Q(\tilde{p}, \tilde{v}) = \frac{1}{2} \sum_{i=1}^{N} (U_i(\tilde{p}) + \tilde{v}_i^T \tilde{v}_i) \end{equation}
\begin{equation}= \frac{1}{2} \sum_{i=1}^{N} (\tilde{V}_i(p_{ij}) + h_i c_1 \tilde{v}_i^T \tilde{v}_i) + \frac{1}{2} \sum_{i=1}^{N} \tilde{v}_i^T \tilde{v}_i \geq \frac{1}{2} \sum_{i=1}^{N} \tilde{V}_i(p_{ij}) \geq Q_0, \end{equation}
which contradicts the condition $Q(t) \leq Q_0$, and this hypothesis is not valid. Therefore, part (i) is proven.

6. Simulation Results and Discussion

In this section, examples of numerical simulation are provided to validate the effectiveness of the improved flocking algorithm proposed in Section 4.

Consider a group of 150 agents moving in a two-dimensional space. The initial positions of the agents are randomly distributed in the box $[-50, 50] \times [-50, 50]$, as
Figure 3: The flocking for $N = 150$ agents applying control protocol (26) in 2D. The red triangles denote the positions of all agents; the heading angle of each triangle represents the velocity direction of the corresponding agent, whereas the red $\times$ denotes the position of the group objective. An undirected edge connecting two agents means that they are neighbors of each other. (a) $t = 0$ sec. (b) $t = 5$ sec. (c) $t = 10$ sec. (d) $t = 20$ sec. (e) $t = 30$ sec. (f) $t = 40$ sec.
shown in Figure 3(a). The initial velocities of the agents are randomly chosen from the box $[0, 2] \times [0, 2]$. The initial position and velocity of the group objective are set as $p_c(0) = [0, 0]^T$ and $v_c(0) = [0.5, 0]^T$, respectively, and marked with a red $\times$. The other parameters are given in Table 1, which remain fixed throughout all simulations. It is worth noting that, for convenience, the above control parameters of the proposed algorithm are all derived from the works of Olfati-Saber [7], Su et al. [20], and Lehn [36]. Furthermore, all simulations are implemented on a platform with the following configurations: 2.20GHz CPU, 8.00GB RAM, Windows 10, and MATLAB R2016b.

Figure 3 displays consecutive snapshots of the flocking for $N = 150$ agents applying control protocol (26). To decrease the communication energy consumption, only 30% of the group is randomly selected as the informed agent; hence, there are 45 informed agents. Therefore, the initial state of the group is seen as highly disconnected subgroups, and fragmentation appears in this case, as shown in Figure 3(a). From Figures 3(b) to 3(d), it is clear that more and more agents are connected, but at a certain distance (where no collisions occur). Finally, the group of $N = 150$ agents form a geometric flocking, as shown in Figure 3(e), and maintain the formation thereafter, as shown in Figure 3(f).

The velocity consensus during the flocking process is shown in Figure 4. The plot shows that the velocities in the $x$- and $y$-axes eventually converge to a constant set, which means that the velocities of all agents become consistent asymptotically. Therefore, the velocity consensus is achieved.

The above observations are fully consistent with our predictions of Theorem 4, indicated in Section 5.

In the simulation stated above, we assume that all the agents are of equal mass and size and have the same sensing radius $r$ and equilibrium distance $d_e$. Nevertheless, for instance, due to the difference in the growth background, social status, and economic strength, etc., it is impossible for two people to be completely identical, in reality. Even for a pair of twins, their cognition and ideology will be different. Therefore, we will fully consider the individual heterogeneity in the following simulations, as we are interested in more than two different species of flocking.

Considering the difference in cognitive level among individuals, we divide the group of $N = 150$ agents into three species, whose equilibrium distances $d_e$ are set as 5, 10, and 15, respectively. Correspondingly, their sensing radiuses $r$ are set as 6, 12, and 18, respectively (according to the work of Olfati-Saber [7]). The remaining parameters are the same as given in Table 1.

Figure 5 displays consecutive snapshots of the flocking for three species of agents applying the control protocol (26). The initial positions of all agents are randomly distributed in the box $[-70, 70] \times [-70, 70]$. The initial velocities of all agents are randomly chosen from the box $[-1, 1] \times [-1, 1]$. The initial position and velocity of the group objective are set

<table>
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<th>Table 1: Parameters of the improved flocking algorithm.</th>
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<td>Parameters</td>
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<tr>
<td>The equilibrium distance</td>
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<td>The sensing radius</td>
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<td>Repulsion coefficient of function (24)</td>
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<td>Attraction coefficient of function (24)</td>
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<td>Coefficients of control protocol (26)</td>
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![Figure 4](image-url)
Figure 5: The flocking for three species of agents applying control protocol (26) in 2D. The triangles denote the positions of all agents; the heading angle of each triangle represents the velocity direction of the corresponding agent, whereas the red $\times$ denotes the position of the group objective. The red, green, and blue triangles represent agents whose equilibrium distance $d_e$ is 5, 10, and 15, respectively. The numbers of red, green, and blue triangles are 110, 30 and 10, respectively. An undirected edge connecting two agents means that they are neighbors of each other. (a) $t = 0$ sec. (b) $t = 4$ sec. (c) $t = 18$ sec. (d) $t = 30$ sec. (e) $t = 36$ sec. (f) $t = 50$ sec.
Figure 6: The flocking for three species of agents applying control protocol (26) in 2D. The triangles denote the positions of all agents; the heading angle of each triangle represents the velocity direction of the corresponding agent, whereas the red x denotes the position of the group objective. The red, green, and blue triangles represent the agent whose equilibrium distance $d_e$ is 5, 10, and 15, respectively, and the number of the three species is equal to 50. An undirected edge connecting two agents means that they are neighbors of each other. (a) $t = 0$ sec. (b) $t = 5$ sec. (c) $t = 18$ sec. (d) $t = 25$ sec. (e) $t = 40$ sec. (f) $t = 50$ sec.
as $p_y(0) = [-10, 0]^T$ and $v_y(0) = [0.5, 0]^T$, respectively. Similarly, there are 45 informed agents. The red, green, and blue triangles represent agents with equilibrium distances $d_e$ of 5, 10, and 15, respectively. The numbers of red, green, and blue triangles are 110, 30, and 10, respectively. Hence, according to the initial state of the group shown in Figure 5(a), the lengths of links between the neighboring agents are not equal. From Figure 5(b) to 5(e), it is clear that the group of $N = 150$ agents eventually form an irregular flocking and maintain the formation thereafter, as can be seen in Figure 5(f). Likewise, no collisions are seen to occur during the flocking process. In this simulation, we find that the green and blue triangles asymptotically move to the periphery and eventually become distributed around the edge of the group, when the multiagent system is stable.

Figure 5(a) shows the initial state of the group, where the agents are positioned with various distances from the center of the group. The green and blue triangles are positioned closer to the center, while the red triangles are positioned further away. The numbers of agents in each category are indicated by the different colors and sizes of the triangles.

In Figure 5(b), the group starts to form an irregular flocking pattern. The distances between the agents are gradually adjusted to maintain a stable formation. In Figure 5(c), the flocking pattern becomes more defined, with the agents moving closer to the edge of the group. In Figure 5(d), the flocking pattern is further refined, with the agents aligning their positions to maintain a cohesive flocking formation. In Figure 5(e), the flocking pattern stabilizes, with the agents maintaining their positions and distances to form a stable flocking formation.

In Figure 5(f), the flocking pattern is shown to have stabilized, with the agents maintaining their positions and distances to form a stable flocking formation. The flocking pattern is shown to have stabilized, with the agents maintaining their positions and distances to form a stable flocking formation.

It is also worth noting that there are very few agents with large equilibrium distances that are surrounded by agents with small equilibrium distances. At this point, its links appear denser because its sensing radius $r$ is larger than that of the neighboring agents.

To compare with the typical flocking algorithm [7] and the rigid flocking algorithm based on the IWKCF [25], a simple model is used to calculate the communication energy consumption [38], which is usually applied to wireless microsensor networks. In [38], the radio hardware energy dissipation is related to the distance $d$ between the transmitter and receiver. In other words, for a relatively short distance, the energy dissipation is inversely proportional to $d^2$, whereas $d^4$ for a relatively long distance. Hence, the communication energy consumption for transmitting a 1-bit message with a distance $d$ is defined as

$$E_{\text{Tx}}(l, d) = E_{\text{Tx-electr}}(l) + E_{\text{Tx-amp}}(l, d)$$

$$= \begin{cases} 
   lE_{\text{electr}} + lE_{\text{amp}}d^2, & d < d_c, \\
   lE_{\text{electr}} + lE_{\text{amp}}d^4, & d \geq d_c,
\end{cases}$$

(38)
where $E_{\text{elec}}$ denotes the electronic energy, which depends on factors such as the digital coding, filtering, and spreading of the signal, etc. $\epsilon_{fs}$ and $\epsilon_{mp}$ are related to the sensitivity of the receiver and the admissible bit-error rate, respectively. $d_c$ represents the cross-over distance. For the theoretical experiments in this work, these parameters are set as $E_{\text{elec}} = 50 \text{nJ/bit}$, $\epsilon_{fs} = 10 \text{ pJ/bit/m}^2$, $d_c = 100 \text{ m}$, and $\epsilon_{mp} = 0.0013 \text{ pJ/bit/m}^2$, which are the same as the experiments conducted in [25, 38].

Figure 8 shows the communication energy consumption compared with the proposed flocking algorithm, the algorithm proposed by Olfati-Saber [7] and Luo et al. [25]. The plot reveals that the communication energy consumption in this paper is significantly lower than that obtained in previous works, which has mainly benefited from the inclusion principle and improved flocking algorithm.

7. Conclusion

In this paper, the flocking problem of multi-agents with partial information exchange was investigated based on the inclusion principle. For the purpose of reducing the communication energy consumption, an improved flocking algorithm has been proposed to achieve a stable collective behavior of the multi-agent system. The stability of the multiagent system was analyzed using Theorem 4 and its relevant proof. In addition, considering the individual heterogeneity in both nature and engineering applications, we then investigated the flocking problem of multi-agents with different sensing radii and equilibrium distances. Finally, numerical simulations verified the effectiveness of the proposed results. Compared with the above simulation results, we found that there exists an interesting self-organizing classification during the flocking process. It is believed that this phenomenon can provide an insight into the properties of the different species of individual flocking. However, the main limitation of this paper is that the proposed flocking model is only applicable to linear systems. Therefore, future works include, but are not limited to, the following research interests: (1) considering the flocking problem of multi-agents with nonlinear dynamics, (2) applying the proposed flocking algorithm to the formation control of swarm robots.

Data Availability

The data used to support the findings of this study are available from the corresponding author upon request.

Conflicts of Interest

The authors declare that there are no conflicts of interest regarding the publication of this paper.

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